

# Memory and Arrays

CSE 333 Spring 2018

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# Administrivia

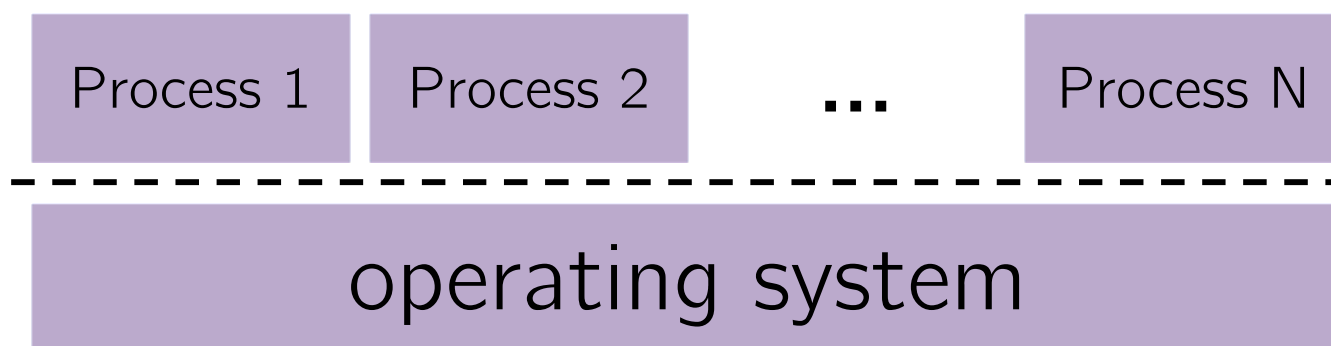
- ❖ Pre-Course Survey & Mini-Bio due tomorrow night
- ❖ Exercise 0 was due this morning
  - Solutions will be posted today after 4 pm
- ❖ Exercise 1 out today and due Friday morning
- ❖ Homework 0 released today
  - Logistics and infrastructure for projects
  - Demos and setup in sections this week – **bring laptop!**
    - Slightly updated CSE VM this quarter – run **\$ sudo yum update** if older version already installed

# Lecture Outline

- ❖ **C's Memory Model** (refresher)
- ❖ Pointers (refresher)
- ❖ Arrays

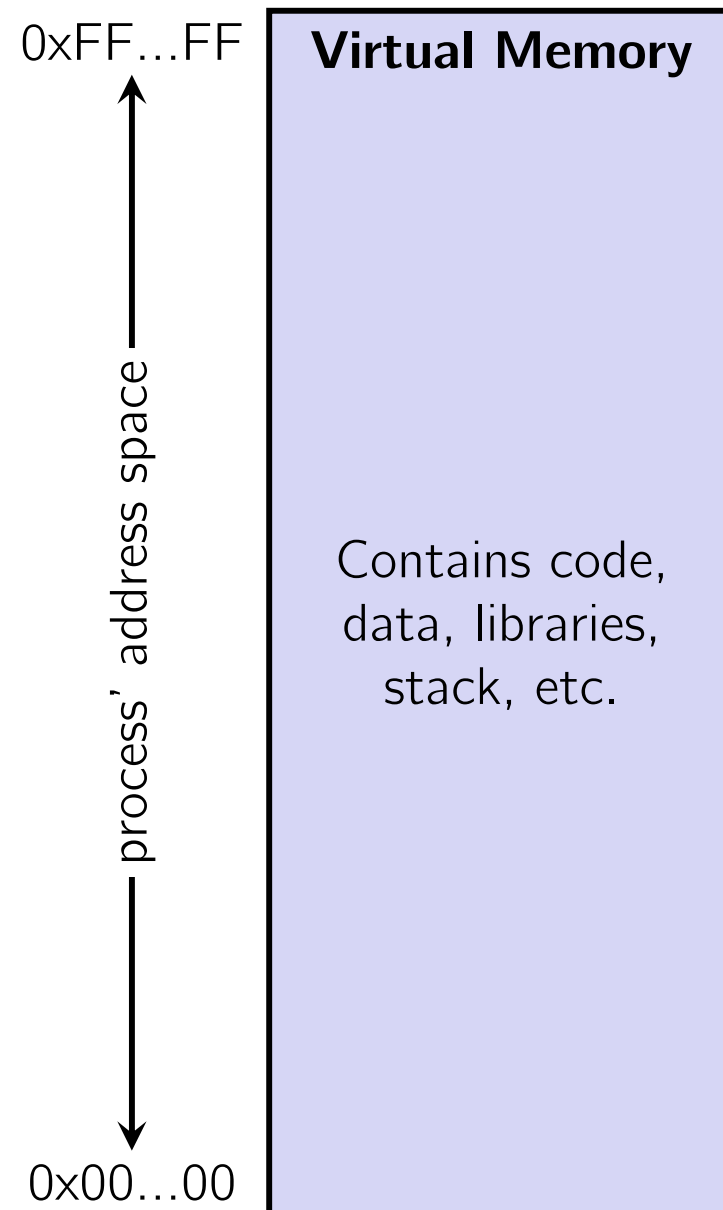
# OS and Processes

- ❖ The OS lets you run multiple applications at once
  - An application runs within an OS “process”
  - The OS timeslices each CPU between runnable processes
    - This happens *very quickly*:  $\sim 100$  times per second



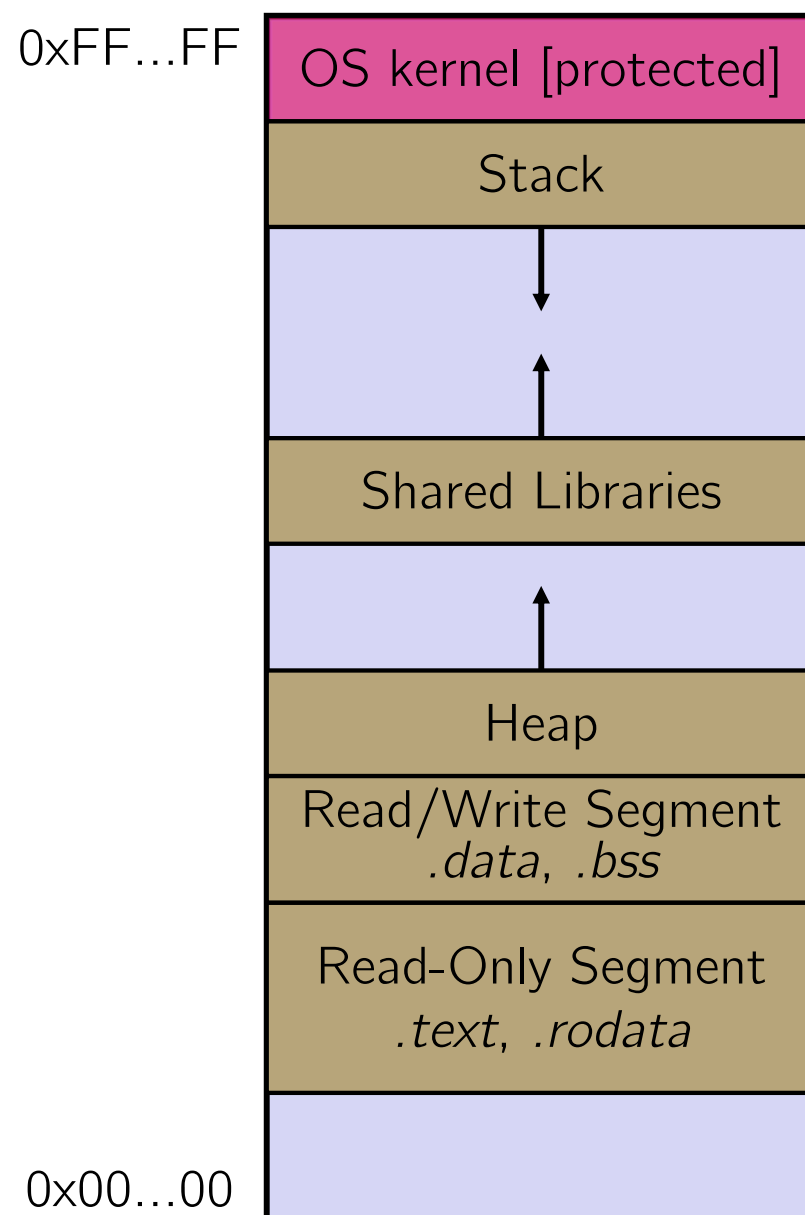
# Processes and Virtual Memory

- ❖ The OS gives each process the illusion of its own private memory
  - Called the process' **address space**
  - Contains the process' virtual memory, visible only to it (via translation)
  - $2^{64}$  bytes on a 64-bit machine



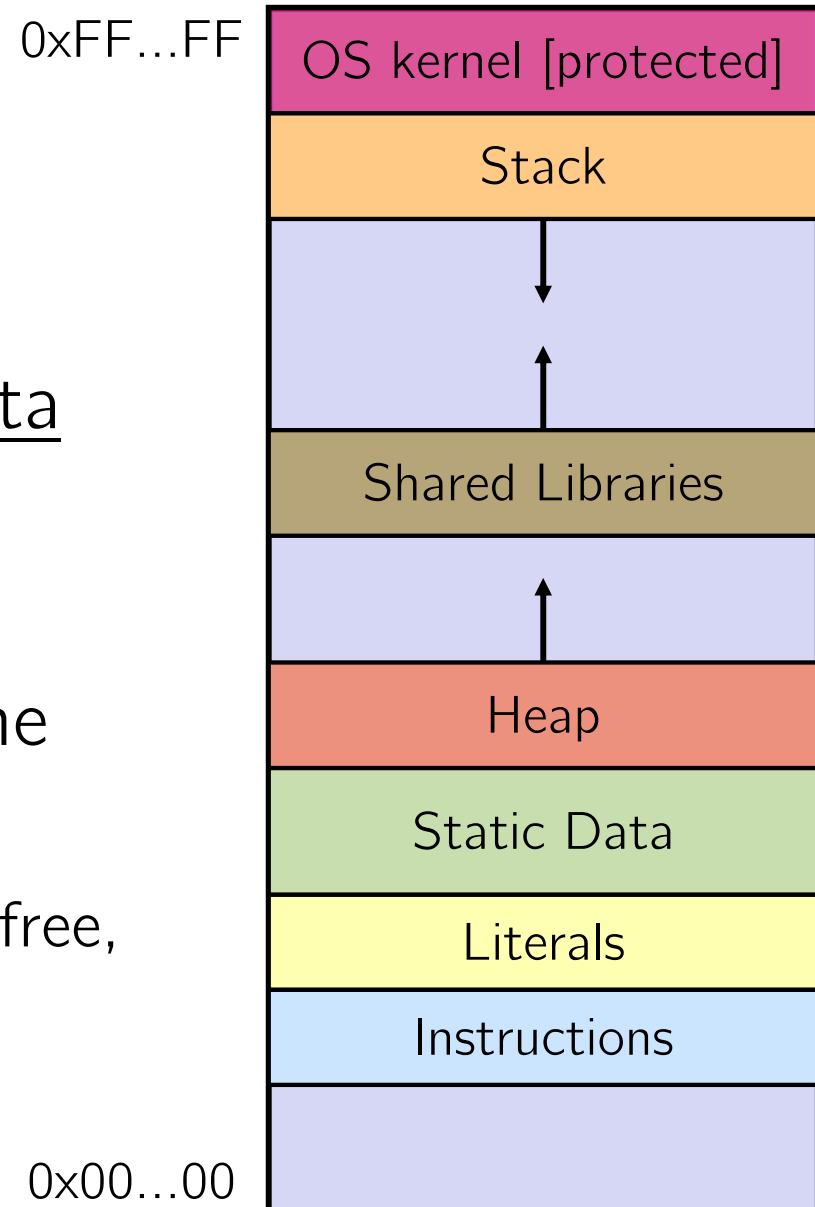
# Loading

- ❖ When the OS loads a program it:
  - 1) Creates an address space
  - 2) Inspects the executable file to see what's in it
  - 3) (Lazily) copies regions of the file into the right place in the address space
  - 4) Does any final linking, relocation, or other needed preparation



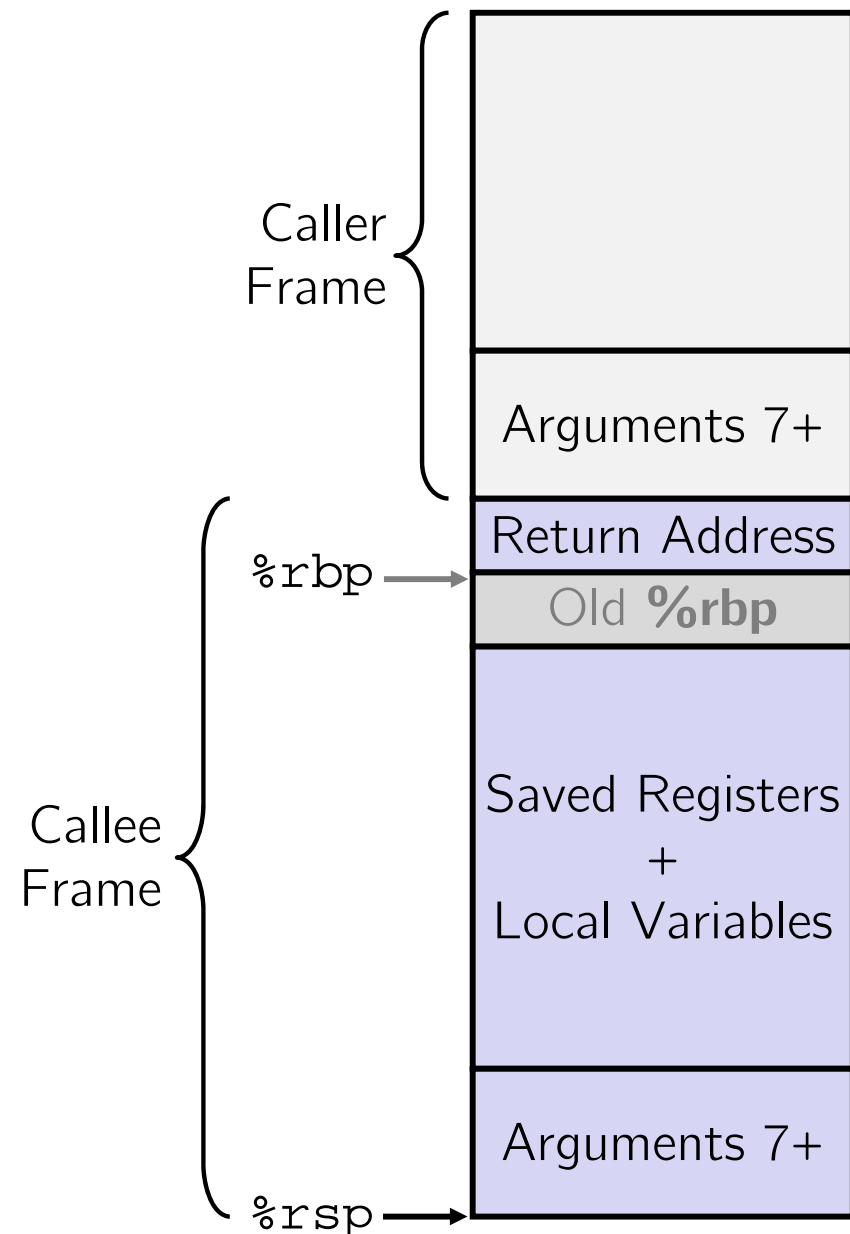
# Memory Management

- ❖ *Local* variables on the Stack
  - Allocated and freed via calling conventions (`push`, `pop`, `mov`)
- ❖ *Global* and *static* variables in Data
  - Allocated/freed when the process starts/exits
- ❖ *Dynamically-allocated* data on the Heap
  - `malloc()` to request; `free()` to free, otherwise **memory leak**



# Review: The Stack

- ❖ Used to store data associated with function calls
  - Compiler-inserted code manages stack frames for you
- ❖ Stack frame (x86-64) includes:
  - Address to return to
  - Saved registers
    - Based on calling conventions
  - Local variables
  - Argument build
    - Only if > 6 used





# Stack in Action

Note: arrow points to *next* instruction to be executed (like in gdb).

stack.c

```

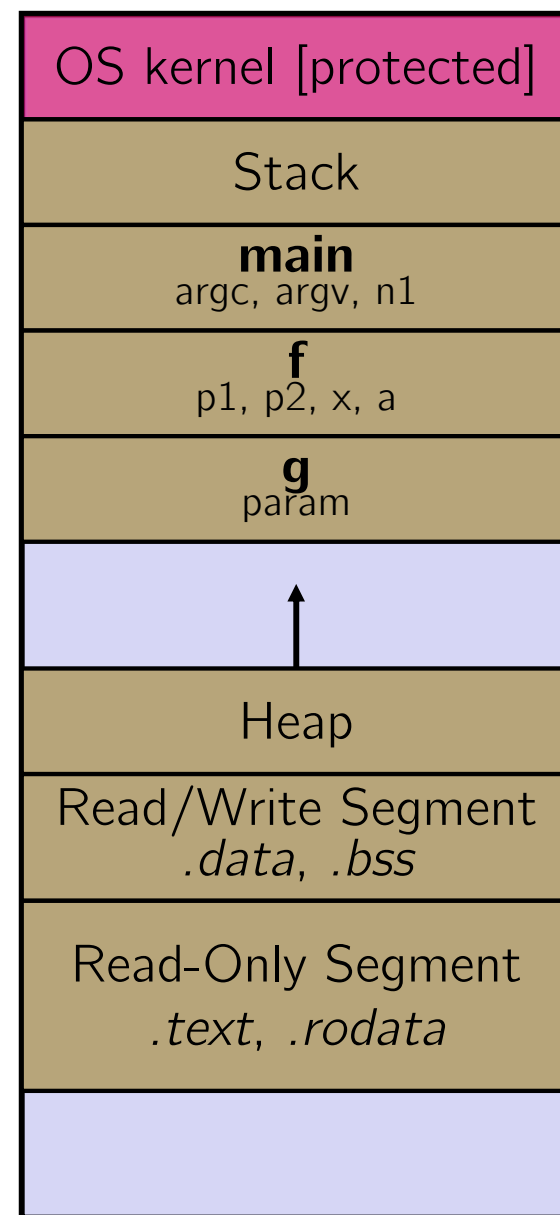
#include <stdint.h>

int f(int, int);
int g(int);

int main(int argc, char** argv) {
    int n1 = f(3, -5);
    n1 = g(n1);
}

int f(int p1, int p2) {
    int x;
    int a[3];
    ...
    x = g(a[2]);
    return x;
}

int g(int param) {
    return param * 2;
}
    
```



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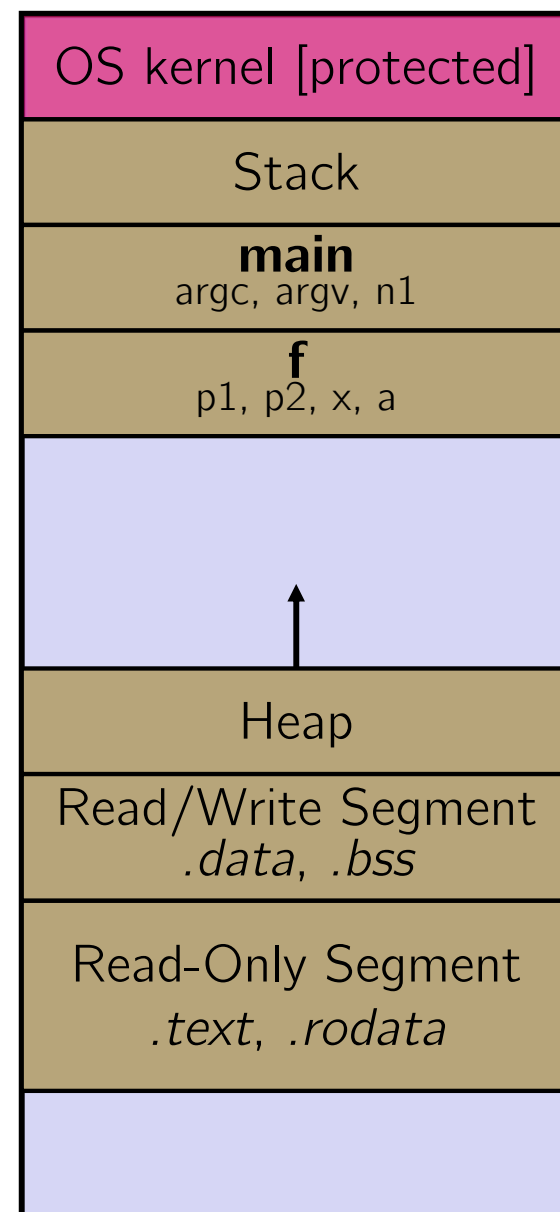
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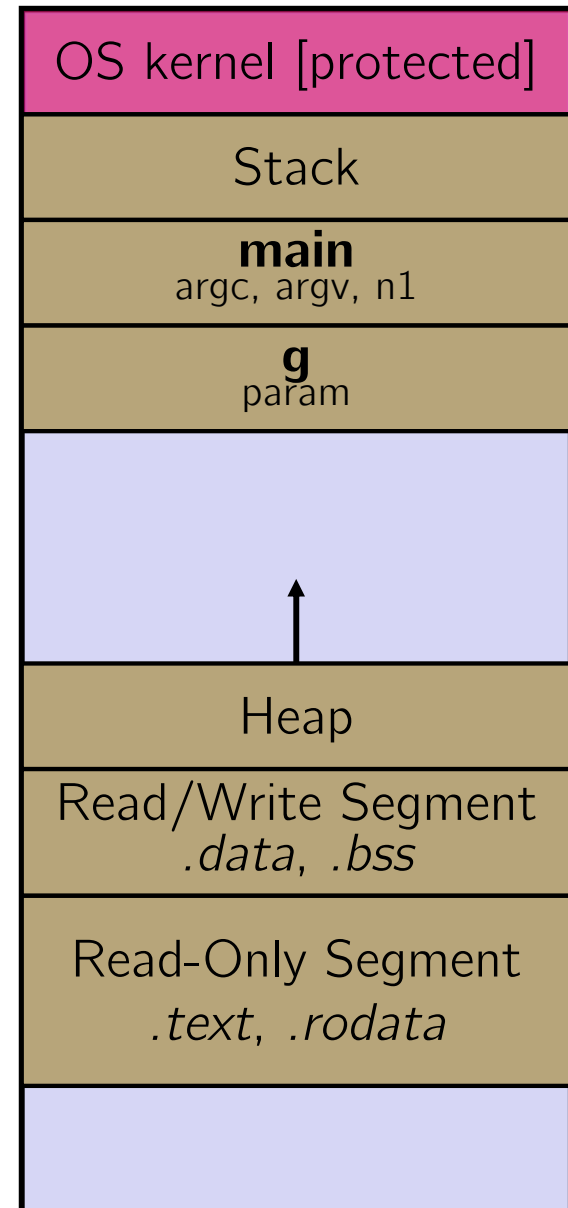
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int main(int argc, char** argv) {
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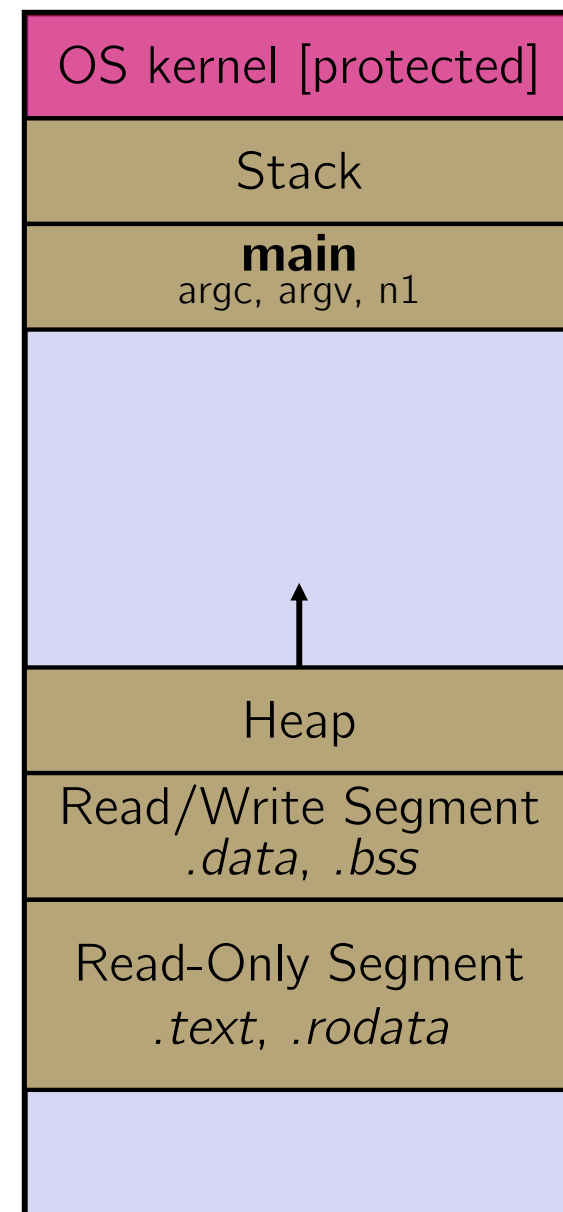
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int f(int, int);
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int main(int argc, char** argv) {
    int n1 = f(3, -5);
    n1 = g(n1);
}

int f(int p1, int p2) {
    int x;
    int a[3];
    ...
    x = g(a[2]);
    return x;
}

int g(int param) {
    return param * 2;
}
```



# Lecture Outline

- ❖ C's Memory Model (refresher)
- ❖ **Pointers** (refresher)
- ❖ Arrays

# Pointers

- ❖ Variables that store addresses
  - It points to somewhere in the process' virtual address space
  - `&foo` produces the virtual address of `foo`
- ❖ Generic definition: `type* name;` or `type *name;`
  - Recommended to not define multiple pointers on same line:  
`int *p1, p2;` not the same as `int *p1, *p2;`
  - Instead, use:  
`int *p1;`  
`int *p2;`
- ❖ *Dereference* a pointer using the unary `*` operator
  - Access the memory referred to by a pointer

# Pointer Example

pointy.c

```
#include <stdio.h>
#include <stdint.h>

int main(int argc, char** argv) {
    int x = 351;
    int* p;      // p is a pointer to a int

    p = &x;     // p now contains the addr of x
    printf("&x is %p\n", &x);
    printf(" p is %p\n", p);
    printf(" x is %d\n", x);

    *p = 333;   // change value of x
    printf(" x is %d\n", x);

    return 0;
}
```

# Something Curious

- ❖ What happens if we run `pointy.c` several times?

```
bash$ gcc -Wall -std=c11 -o pointy pointy.c
```

Run 1:

```
bash$ ./pointy
&x is 0x7ffff9e28524
p is 0x7ffff9e28524
x is 351
x is 333
```

Run 2:

```
bash$ ./pointy
&x is 0x7fffe847be34
p is 0x7fffe847be34
x is 351
x is 333
```

Run 3:

```
bash$ ./pointy
&x is 0x7fffe7b14644
p is 0x7fffe7b14644
x is 351
x is 333
```

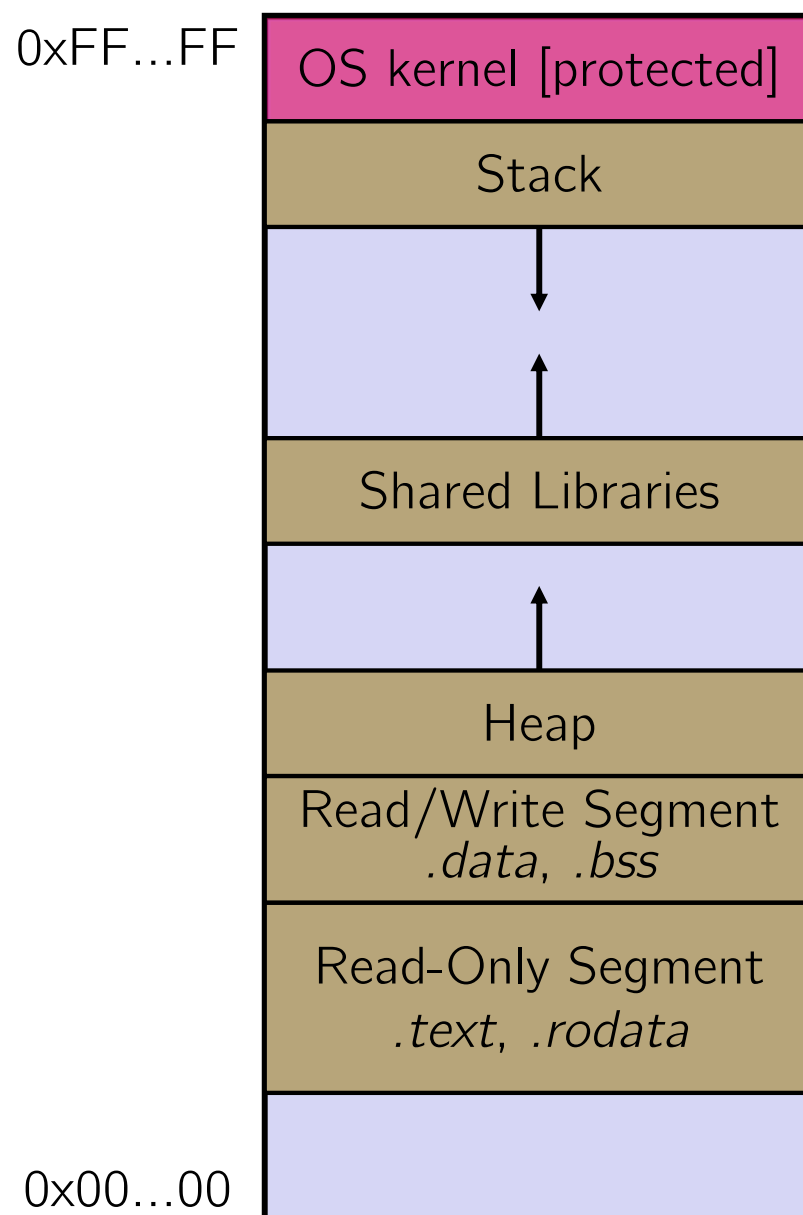
Run 4:

```
bash$ ./pointy
&x is 0x7fffff0dfe54
p is 0x7fffff0dfe54
x is 351
x is 333
```



# Address Space Layout Randomization

- ❖ Linux uses *address space layout randomization* (ASLR) for added security
  - Randomizes:
    - Base of stack
    - Shared library (`mmap`) location
  - Makes Stack-based buffer overflow attacks tougher
  - Makes debugging tougher
  - Can be disabled (`gdb` does this by default); Google if curious



# Lecture Outline

- ❖ C's Memory Model (refresher)
- ❖ Pointers (refresher)
- ❖ **Arrays**

# Arrays

- ❖ Definition: `type name[size]`
  - Allocates `size * sizeof(type)` bytes of *contiguous* memory
  - Normal usage is a compile-time constant for `size` (e.g. `int scores[175];`)
  - Initially, array values are “garbage”
- ❖ Size of an array
  - Not stored anywhere – array does not know its own size!
    - `sizeof(array)` only works in variable scope of array definition
  - Recent versions of C allow for variable-length arrays
    - Uncommon and can be considered bad practice [*we won't use*]

```
int n = 175;  
int scores[n]; // OK in C99
```

# Challenge Question

- ❖ The code snippets both use a variable-length array. What will happen when we compile with C99?
  - Vote at <http://PollEv.com/justinh>

```
int m = 175;
int scores[m];

void foo(int n) {
    ...
}
```

```
int m = 175;

void foo(int n) {
    int scores[n];
    ...
}
```

- A. **Compiler Error**
- B. **Compiler Error**
- C. **No Error**
- D. **No Error**
- E. **We're lost...**

- Compiler Error**
- No Error**
- Compiler Error**
- No Error**

# Using Arrays

- ❖ Initialization: `type name[size] = {val0, ..., valN};`
  - `{ }` initialization can *only* be used at time of definition
  - If no `size` supplied, infers from length of array initializer
- ❖ Array name used as identifier for “collection of data”
  - `name[index]` specifies an element of the array and can be used as an assignment target or as a value in an expression
  - Array name (by itself) produces the address of the start of the array
    - Cannot be assigned to / changed

```
int primes[6] = {2, 3, 5, 6, 11, 13};  
primes[3] = 7;  
primes[100] = 0; // memory smash!
```

# Multi-dimensional Arrays

## ❖ Generic 2D format:

```
type name[rows][cols] = {{values}, ..., {values}};
```

- Still allocates a single, contiguous chunk of memory
- C is *row-major*

```
// a 2-row, 3-column array of doubles
double grid[2][3];

// a 3-row, 5-column array of ints
int matrix[3][5] = {
    {0, 1, 2, 3, 4},
    {0, 2, 4, 6, 8},
    {1, 3, 5, 7, 9}
};
```

# Parameters: reference vs. value

- ❖ There are two fundamental parameter-passing schemes in programming languages
- ❖ **Call-by-value**
  - Parameter is a local variable initialized when the function is called and gets a copy of the calling argument; manipulating the parameter only changes copy, *not* the calling argument
  - **C, Java**, C++ primitives
- ❖ **Call-by-reference**
  - Parameter is an alias for the supplied argument; manipulating the parameter manipulates the calling argument
  - C++ arrays and references (we'll see more later)

# Arrays as Parameters

- ❖ It's tricky to use arrays as parameters
  - What happens when you use an array name as an argument?
  - Arrays do not know their own size

```
int sumAll(int a[]); // prototype

int main(int argc, char** argv) {
    int numbers[] = {9, 8, 1, 9, 5};
    int sum = sumAll(numbers);
    return 0;
}

int sumAll(int a[]) {
    int i, sum = 0;
    for (i = 0; i < ...???)
}
```



# Solution 1: Declare Array Size

```
int sumAll(int a[5]); // prototype

int main(int argc, char** argv) {
    int numbers[] = {9, 8, 1, 9, 5};
    int sum = sumAll(numbers);
    printf("sum is: %d\n", sum);
    return 0;
}

int sumAll(int a[5]) {
    int i, sum = 0;
    for (i = 0; i < 5; i++) {
        sum += a[i];
    }
    return sum;
}
```

- ❖ Problem: loss of generality/flexibility!

# Solution 2: Pass Size as Parameter

```
int sumAll(int a[], int size); // prototype

int main(int argc, char** argv) {
    int numbers[] = {9, 8, 1, 9, 5};
    int sum = sumAll(numbers, 5);
    printf("sum is: %d\n", sum);
    return 0;
}

int sumAll(int a[], int size) {
    int i, sum = 0;
    for (i = 0; i < size; i++) {
        sum += a[i];
    }
    return sum;
}
```

arraysum.c

# Returning an Array

- ❖ Local variables, including arrays, are allocated on the Stack
  - They “disappear” when a function returns!
  - Can't safely return local arrays from functions
    - Can't return an array as a return value – why not?

```
int* copyArray(int src[], int size) {  
    int i, dst[size];    // OK in C99  
  
    for (i = 0; i < size; i++) {  
        dst[i] = src[i];  
    }  
  
    return dst;    // no compiler error, but wrong!  
}
```

buggy\_copyarray.c

# Solution: Output Parameter

- ❖ Create the “returned” array in the caller
  - Pass it as an **output parameter** to `copyarray()`
    - A pointer parameter that allows the callee to leave values for the caller to use
  - Works because arrays are “passed” as pointers
    - “Feels” like call-by-reference, *but it's not*

```
void copyArray(int src[], int dst[], int size) {  
    int i;  
  
    for (i = 0; i < size; i++) {  
        dst[i] = src[i];  
    }  
}
```

copyarray.c

# Output Parameters

- ❖ Output parameters are common in library functions
  - `long int strtol(char* str, char** endptr, int base);`
  - `int sscanf(char* str, char* format, ...);`

```
int    num, i;
char*  pEnd, str1 = "333 rocks";
char   str2[10];

// converts "333 rocks" into long -- pEnd is conversion end
num = (int) strtol(str1, &pEnd, 10);

// reads string into arguments based on format string
num = sscanf("3 blind mice", "%d %s", &i, str2);
```

outparam.c

# Extra Exercises

- ❖ Some lectures contain “Extra Exercise” slides
  - Extra practice for you to do on your own without the pressure of being graded
  - You may use libraries and helper functions as needed
    - Early ones may require reviewing 351 material or looking at documentation for things we haven't reviewed in 333 yet
  - Always good to provide test cases in `main( )`
  
- ❖ Solutions for these exercises will be posted on the course website (as `extra#.c` or `extra#.cc`)
  - You will get the most benefit from implementing your own solution before looking at the provided one

# Extra Exercise #1

- ❖ Write a function that:
  - Accepts an array of 32-bit unsigned integers and a length
  - Reverses the elements of the array in place
  - Returns nothing (`void`)

# Extra Exercise #2

- ❖ Write a function that:
  - Accepts a string as a parameter
  - Returns:
    - The first white-space separated word in the string as a newly-allocated string
    - AND the size of that word