

# Lecture 14 - More Graphs

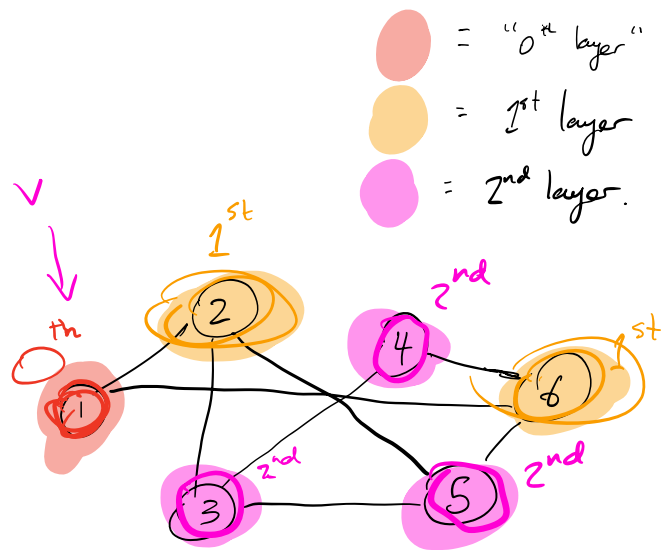
## Announcements:

- Sorting Assignments out - start early!
- Midterm grades out probably tomorrow
- Mid-quarter feedback form out tomorrow  
get +2 midterm points for filling out!

## Last time: Breadth-First-Search

Idea: pick source node v

- 1) process all nodes 1 away from v  
all neighbors
- 2) process all nodes 2 away from v  
all neighbors' neighbors
- 3) " " 3 " "
- ...



## BFS Pseudocode

bfs( $G, v$ ):

found = new empty queue

found.enqueue(v)

mark v as "visited"

while found is not empty:

curr = found.dequeue() ←

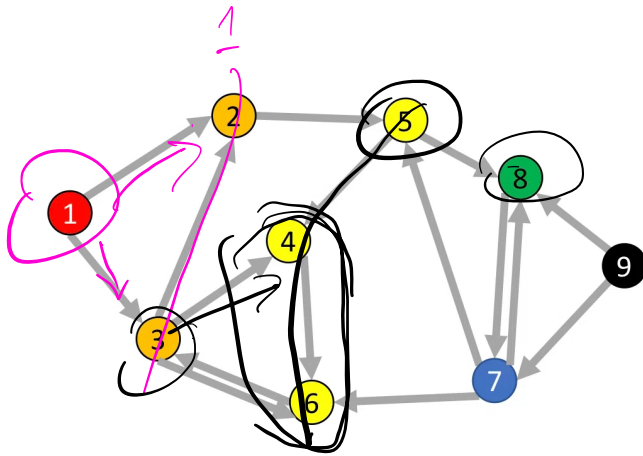
for u in neighbors(curr):

if u not visited already:

found.enqueue(u)

mark u as "visited"

Ex



Queue:

~~2~~ 3 5 4 6

| node | visited? | depth |
|------|----------|-------|
| 1    | True     | 0     |
| 2    | True     | 1     |
| 3    | T        | 1     |
| 4    | T        | 2     |
| 5    | T        | 2     |
| 6    | T        | 2     |
| 7    |          |       |
| 8    | T        | 3     |
| 9    |          |       |

## Find Distance Pseudocode

findDistance( $G, s, t$ )

found = new empty queue

found.enqueue( $s$ )

→ layer/depth( $s$ ) = 0

mark  $s$  as "visited"

while found is not empty:

curr = found.dequeue

→ layer = layer of curr

for  $u$  in neighbors(curr):

if  $u$  not visited:

found.enqueue( $u$ )

mark  $u$  as "visited"

depth( $u$ ) = layer(curr) + 1

→ [ return depth( $t$ ) ]

Idea: when vertex is seen first, record its depth!

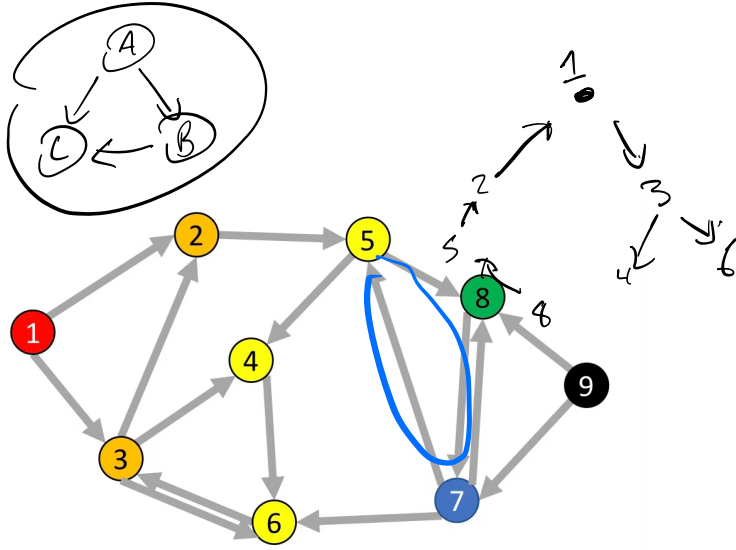
## Shortest Path:

We can use a very similar idea to find shortest path from  $s$  to  $t$ !

idea: instead of (or in addition to) recording depth of  $v$ , record  $v$ 's "previous": the

find  $u$  that adds  $v$  to the queue.

Ex



Queue:

~~2~~ ~~3~~ ~~5~~ 4 6 8

| node | visited? | previous     |
|------|----------|--------------|
| 1    | True     | <del>0</del> |
| 2    | T        | <del>1</del> |
| 3    | T        | 1            |
| 4    | T        | 3            |
| 5    | T        | <del>2</del> |
| 6    | T        | 3            |
| 7    |          |              |
| 8    |          | 5            |
| 9    |          |              |

1-2-3-4-5-6-7-8

Longest Path?

This is NP-hard (probably not solvable quickly)

Traversal #2: Depth-first Search

Idea: explore each path as deep as possible:

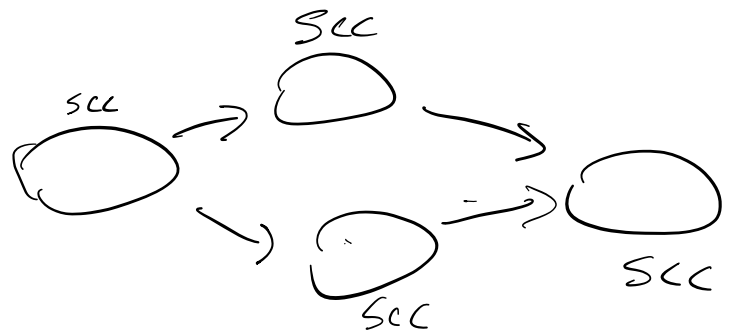
- 1) visit start node  $s$
- 2) visit a neighbor of  $s$ , then all nodes reachable from that neighbor, ...
- 3) repeat with next neighbor



What can we learn from this?

finding cycles!

topological sort



## DFS non-recursive Pseudocode

$dfs(G, s)$ :

found = new empty stack

found. push (s)

mark s as "visited"

while found is not empty:

curr = found.pop

for u in neighbors(curr):

if u not visited:

found. push (u)

mark u as "visited"

runtime?

$$\Theta(|V| + |E|)$$

## Recursive implementation?

$dfs(G, curr)$ :

mark curr as "visited" ←

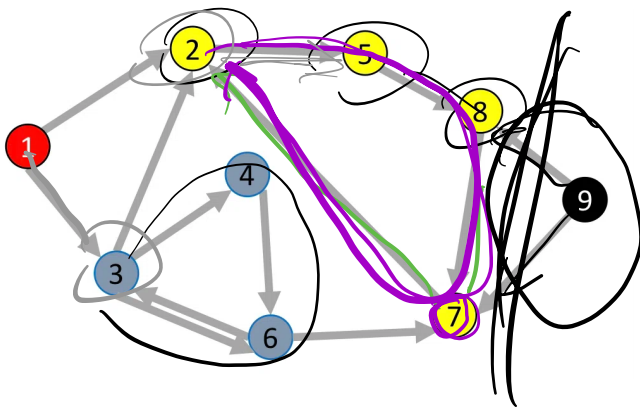
for u in neighbors(curr):

if not visited(u):

$dfs(G, u)$

mark curr as "done" ←

Ex 1



call stack :

7 8 5 2 1

| node | visited? | done?          | other |
|------|----------|----------------|-------|
| 1    | True     | F              |       |
| 2    | T        | <del>E</del> T |       |
| 3    | T        | F              |       |
| 4    |          |                |       |
| 5    | T        | <del>E</del> T |       |
| 6    |          |                |       |
| 7    | T        | <del>E</del> T |       |
| 8    | T        | <del>E</del> T |       |
| 9    |          |                |       |

## Characterizing Edges in DFS

Tree Edges

cross edges

$$t_{done}(b) < t_{visit}(a)$$

forward edges

$$t_{visit}(a) < t_{visit}(b) < t_{done}(b) < t_{done}(a)$$

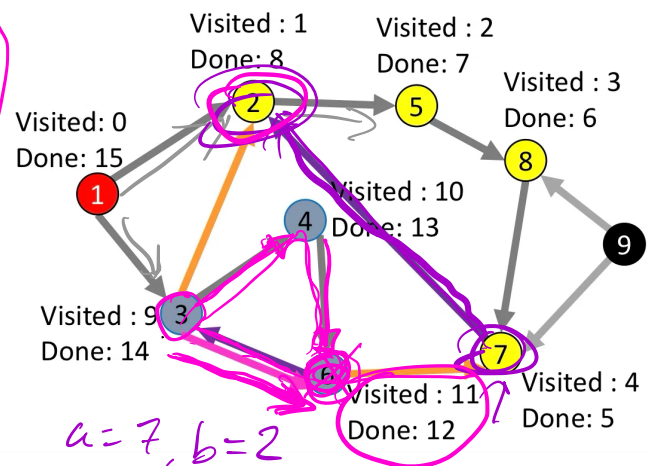
"pointing to an ancestor"

back edges

$$t_{visit}(b) < t_{visit}(a) \rightarrow t_{done}(a) < t_{done}(b)$$

current nodes neighbor is "in progress" / reachable

⇒ CYCLE!



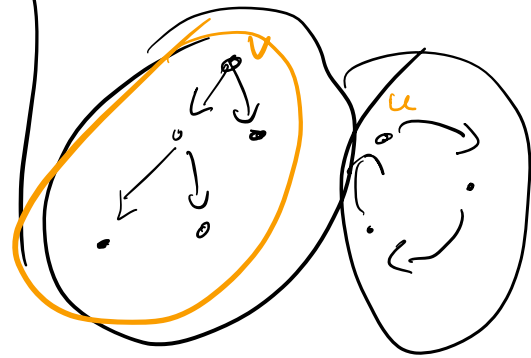
## Cycle Detection

idea: look for a back edge!

```
cycleDetect (G, curr):  
    mark curr as "visited"  
    cycleFound = False  
    for u in neighbors(curr):  
        if u visited but not done:  
            cycleFound = True  
        if u not visited and not cycleFound:  
            cycleFound = cycleDetect(G, u)  
    mark curr as "done"  
    return cycleFound
```

hasCycle(G)

```
for each node v:  
    if v not done:  
        if hasCycle(G, v):  
            return True  
return False
```



## Single Source Shortest Path

Task: given start node  $s$  in a weighted graph with non-negative edge weights, find the shortest path from  $s$  to all other vertices (or a specific node)

## Dijkstra's Algorithm!

idea: when a vertex is closest "not done" node to start, we have found its shortest path!

## Dijkstra's Pseudocode

① extract node *curr* from priority queue

this makes *curr* "done"

② for each not *done* neighbor of *curr*:

- update its distance if we have found a better path.

Seen = added to PQ

done = removed from PQ

## BFS Pseudocode

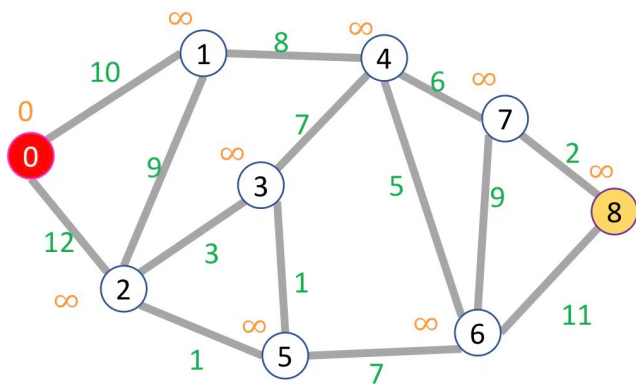
① extract node *curr*

② for each not visited neighbor of *curr*:

add it to queue

depth of neighbor = depth(*curr*) + 1

Ex 1



PQ:

node seen? done? distance

|   |      |  |          |
|---|------|--|----------|
| 0 | True |  | 0        |
| 1 |      |  | $\infty$ |
| 2 |      |  | $\infty$ |
| 3 |      |  | $\infty$ |
| 4 |      |  | $\infty$ |
| 5 |      |  | $\infty$ |
| 6 |      |  | $\infty$ |
| 7 |      |  | $\infty$ |
| 8 |      |  | $\infty$ |

