CSE 332 Spring 2023 Final Exam

Name: _____

Email address (UWNetID): _____

Instructions:

- The allotted time is 1 hour and 50 minutes.
- Please do not turn the page until the staff says to do so.
- This is a closed-book and closed-notes exam. You are NOT permitted to access electronic devices including calculators.
- Read the directions carefully, especially for problems that require you to show work or provide an explanation.
- When provided, write your answers in the box or on the line provided.
- Unless otherwise noted, every time we ask for an O, Ω, or Θ bound, it must be simplified and tight.
- For answers that involve bubbling in a \bigcirc or \bigsqcup , make sure to fill in the shape completely.
- If you run out of room on a page, indicate where the answer continues. Try to avoid writing on the very edges of the pages: we scan your exams and edges often get cropped off.
- A formula sheet has been included at the end of the exam.

Advice

- If you feel like you're stuck on a problem, you may want to skip it and come back at the end if you have time.
- Look at the question titles on the cover page to see if you want to start somewhere other than problem 1.
- Relax and take a few deep breaths. You've got this! :-).

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Total: 97 points

2

Q1: Short-answer questions (10 pts)

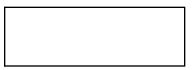
- For questions asking you about runtime, give a simplified, tight Big-O bound. This means that, for example, O(5n² + 7n+ 3) (not simplified) or O(2^{n!}) (not tight enough) are unlikely to get points. Unless otherwise specified, all logs are base 2.
- For questions with a mathematical answer, you may leave you answer as an unsimplified formula (e.g. 7 · 10³).

We will only grade what is in the provided answer box.

a. True or False: The worst-case runtime of delete in a Binary Search Tree containing N elements is $\Omega(\log N)$.

		False
b. What is the maximum number of nodes in a binary max	heap with	height h?
	<i>.</i> .	
c. i) For a Binary Min Heap, in Floyd's build heap algorithr		k one) . Opercolate elements down
ii) For a Binary Min Heap, in Floyd's build heap algorith	m, you go	from (pick one):
\bigcirc the bottom (where the leaf nodes are) up to the top of top of the top of top	he heap	
O the top (where the root node is) down to the bottom of t	he heap	
d. Give the worst-case runtime of an increaseKey operation on	n a Binary N	lin Heap containing N values.
	O()

e. What is the minimum number of nodes in an AVL tree of height 3?



Q2: More Short-answer questions (10 pts)

(Same instructions as for Q1)

a.What is the worst-case runtime for Kruskal's algorithm as described in lecture. (Give your answer in terms of V, E and/or d)

b.What is the worst-case runtime of finding all **incoming edges of a vertex** using an adjacency list representation of a graph. Use the adjacency list as described in lecture - with a linked list node for each edge, NOT what you used for P3. (Give your answer in terms of V, E and/or d)

c. What is the worst-case runtime of a find operation on a hash table containing N elements, using separate chaining, where each bucket is a sorted linked list. The load factor of this table is log N.

d. Assuming the implementations described in lecture, give the name of a **sequential** sort that is **not** in-place but is stable.

e. If 1/4 of your program must be run sequentially (and the remaining 75% can be parallelized) what is the maximum speedup you would expect to get with 5 processors?

3







4 Q3: Hashing (10 pts)

For question a), insert 82, 4, 57, 13, 24, 17, 32 into the tables below. For each table, TableSize = 10, and you should use the primary hash function h(k) = k%10. If an item cannot be inserted into the table, please indicate this and continue inserting the remaining values.

a) [3 pts] **Double Hashing Hashtable**. You should use the primary hash function: h(k) = k % 10 and a secondary hash function: g(k) = 2 + (k % 3).

If any values cannot be inserted, write them here:_____

_	
0	
1	
2	
3 4	
4	
5	
6	
7	
8	
9	

b) [1 pt] What is the load factor for the table in part a)?

c) [5 pts] You are designing a hash table to store your album collection. Each album will be represented by an Album object:

```
public class Album {
    public static final String UUID = "08634082047204792740274a";
    private String artist;
    private String artist;
    private List<Song> songs;

    public Album (String t, String a, List<Song> s) {
        title = t;
        artist = a;
        songs = s;
    }
    public int hashCode() {
        ...
    }
}
```

We define two Albums as being equal if they have the same title, artist, and songs. Just like in the real world, no two albums in a hashmap can have the same title or same set of songs.

For each of the hashCode() implementations below, fill in the appropriate bubble to indicate whether the implementation is valid or invalid.

- A hashcode is "**valid**" if the methods of a hash table implementation (e.g. HashSet) always return the correct results.
- A hashcode is "**invalid**" if the code either doesn't compile, or some methods of a hash table implementation (e.g. HashSet) would not work as expected.

In the code below, you should assume that the calls to hashCode() on objects other than Albums are implemented correctly and efficiently.

hash	Code implementation	valid	invalid
i)	<pre>return title.hashCode() + artist.hashCode();</pre>	\bigcirc	0
ii)	return UUID;	\bigcirc	0
iii)	<pre>return UUID.hashCode();</pre>	0	\bigcirc
iv)	<pre>Random newRand = new Random(); return newRand.nextInt();</pre>	\bigcirc	\bigcirc
♥)	<pre>int result = 0; for (Song s: songs) { result += s.hashCode(); } return result;</pre>	0	\bigcirc

d) [1 pt] Which of the **valid** hashCode implementations in c) would you expect to have the **most** collisions (pick ONE):

 \bigcirc iv)

v)

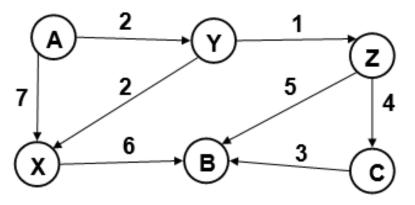
◯ iii)

ii)

i)

⁶ Q4: Graphs (19 pts)

Use the following graph for the problems on this page:



a) [2 pts] List a valid **topological ordering** of the nodes in the graph above (if there are no valid orderings, state why not).

_____, _____,

___, __

Yes

No

b) [4 pts] Step through Dijkstra's Algorithm to calculate the **single source shortest path from A** to every other vertex. Break ties by choosing the lexicographically smallest letter first; ex. if B and C were tied, you would explore B first. *Note that the next question asks you to recall what order vertices were declared known*. Make sure the <u>final</u> distance and predecessor are clear in the table below.

_,

Vertex	Known	Distance	Predecessor
А			
В			
С			
x			
Y			
Z			

c) [1 pt] In what order would Dijkstra's algorithm mark each node as known?

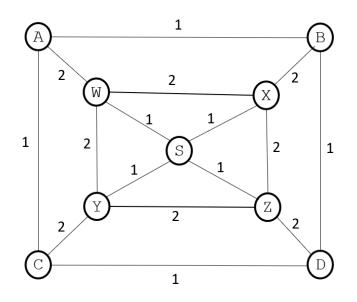
_, _

d) [1 pt] List the **shortest path** from A to B. (Give the actual path **NOT** the cost.)

e) [2 pts] Is this graph complete? Explain your answer in 1-2 sentences for any credit.

_, _

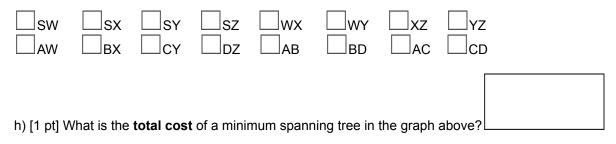
Use the following graph for the problems on this page:



f) [2 pts] If we run Kruskal's algorithm, which of the following edges could be the last edge added to the Minimum Spanning Tree? **Bubble in the box for all that apply.** Assume any ties are broken randomly.

Sw	□sx	SY	Sz	□wx	WY	□xz	Ωyz
AW	Швх	Сү	Dz	АВ	BD		ССС

g) [2 points] If we run Prim's algorithm from vertex S, which of the following edges could be the last edge added to the Minimum Spanning Tree? **Bubble in the box for all that apply.** Assume any ties are broken randomly.



i) [2 pts] ASSUMING the edges above are **unweighted**, give a valid **breadth first search** of this graph, <u>starting at vertex S</u>, using the algorithm described in lecture. When adding elements to the data structure, you should break ties by choosing the lexicographically smallest letter first; ex. if A and B were tied, you would add A to the data structure first. You only need to show the final breadth first search.

j) [2 pts] ASSUMING the edges above are **unweighted**, give a valid **depth first search** of this graph, <u>starting</u> <u>at vertex S</u>, using the non-recursive algorithm described in lecture. When adding elements to the data **structure**, you should break ties by choosing the lexicographically smallest letter first; ex. if A and B were tied, you would add A to the data structure first. You only need to show the final depth first search.

_, _

_, ___

_, ____

_, _

Q5: ForkJoin (14 pts)

In Java using the ForkJoin Framework, write code to solve the following problem:

- Input: A root Node of a valid (i.e. you may assume correct fields) binary tree of unique integers.
- Output: True iff the binary tree is a valid Binary Search Tree (BST). False otherwise.

Examples:

Inputs	5 / \ 3 8 /\ /\ 1 4 7 9	8 / \ 4 10 /\ \ 2 6 9	8 / \ 4 9 /\ 2 10
Output	True	False	False
Why	All numbers are in correct ordering	10 occurs before 9	10 is greater than the root node

Notes:

- Use the given CUTOFF field to employ a sequential cutoff. You may use the given sequential method that determines whether a binary tree is a BST (true) or not (false).
 - Make sure your code has $O(\log n)$ span and O(n) work, where n is the number of nodes.
 - <u>To achieve this span, you may assume the given binary tree is balanced (similar to an AVL tree).</u>
- You may NOT use any global data structures or synchronization primitives (locks).
- Do **NOT** worry about off-by-one errors.

Fill in the <u>underlines</u> in the function validateBST below.

```
import java.util.concurrent.ForkJoinPool;
import java.util.concurrent.RecursiveTask;
import java.util.concurrent.RecursiveAction;
public class Main {
    public static final ForkJoinPool fjPool = new ForkJoinPool();
   private static class Node { // You should use this class
        int key;
        int height;
       Node left, right;
    public static Boolean validateBST(Node root) {
        return fjPool.invoke(new Task(____
                                                                             _));
    }
    /**
     * @param root root of a valid binary tree
     * @param min maximum value root.key can be
     * @param max minimum value root.key can be
     * @return true if the root is the root of a valid BST (where
                     min <= root.key <= max) and false otherwise.</pre>
     * Note: Every node in the tree must also also have a key between min and
            max and be a valid BST.
     * Runs in O(n) time, where n is the number of nodes in the Binary Tree
     **/
    private static boolean sequential (Node root, int min, int max) {
        // some good code here that follows the spec :)
    // Your class goes here (write it on the next page)
```

Write your class here:

L

pub	public static final int CUTOFF = 3; // Fields go here	_ {
	public Task()) {
	}	
	<pre>public compute() {</pre>	

¹⁰ Q6: Concurrency (10 pts)

The UWBookstore class manages the bookstore's stock and balance by selling books and restocking books. Multiple threads (students) might be accessing the same UWBookstore object. Assume Map objects ARE THREAD-SAFE, have enough space, and operations on them will not throw an exception.

```
1
   public class UWBookstore {
2
       private Map<String, Integer> books;
3
       private Map<String, Object> bookLocks;
4
       private int balance;
5
       // The bookstore will ONLY carry the books in bookList and you can assume that
6
7
       // methods in the class will only be called on these books.
8
       public UWBookstore(List<String> bookList) {
           this.books = new HashMap<>();
9
           this.bookLocks = new HashMap<>();
10
11
           for (String book : bookList) {
12
                this.books.put(book, 0);
13
                this.bookLocks.put(book, new Object());
14
            }
15
            this.balance = 100;
        }
16
17
        public void sell(String book, int price) {
18
19
20
            synchronized(bookLocks.get(book)) { // assume book exists
21
22
                if (books.get(book) > 0) {
23
                    books.put(book, books.get(book) - 1); // remove the book from stock
24
25
26
                    balance += price; // The store gets $ for each book it sells.
27
28
                }
29
30
            }
31
32
        }
33
34
        public void restock(String book, int quantity, int costEach) {
35
36
            int totalCost = quantity * costEach;
37
38
            synchronized(bookLocks.get(book)) { // assume book exists
39
40
                if (balance >= totalCost) {
41
42
                    books.put(book, books.get(book) + quantity); // add books to stock
43
44
                    balance -= totalCost; // The store must pay for these books!
45
46
47
                }
48
49
            }
50
        }
51
52 }
```

```
a) [3 pts]Does the UWBookstore class above have (bubble in all that apply):
                          potential for deadlock
                                                        a data race
                                                                              none of these
   a race condition
If there are any problems, describe each problem selected in 1-2 sentences.
b) [3 pts] You decide to add one more method to the UWBookstore class:
58 public void giveawayBundle(String book1, String book2) {
        // Removes one copy of book1 and book2 from stock
59
60
        String firstBook = book1;
61
        String secondBook = book2;
62
63
        if (book1.compareTo(book2) > 0) {
64
65
             firstBook = book2;
             secondBook = book1;
66
67
        }
68
69
70
        synchronized(bookLocks.get(firstBook)) { // assume firstBook exists
71
             synchronized(bookLocks.get(secondBook)) { // assume secondBook exists
72
73
74
                 if (books.get(book1) > 0 && books.get(book2) > 0) {
75
                      books.put(book1, books.get(book1) - 1);
76
                      books.put(book2, books.get(book2) - 1);
77
78
79
                 }
80
81
             }
82
83
        }
84
85
    }
Does adding this method to the UWBookstore class cause any new (bubble in all that apply):
                                                                              none of these
   a race condition
                          potential for deadlock
                                                        a data race
If there are any new problems, describe each problem selected in 1-2 sentences.
```

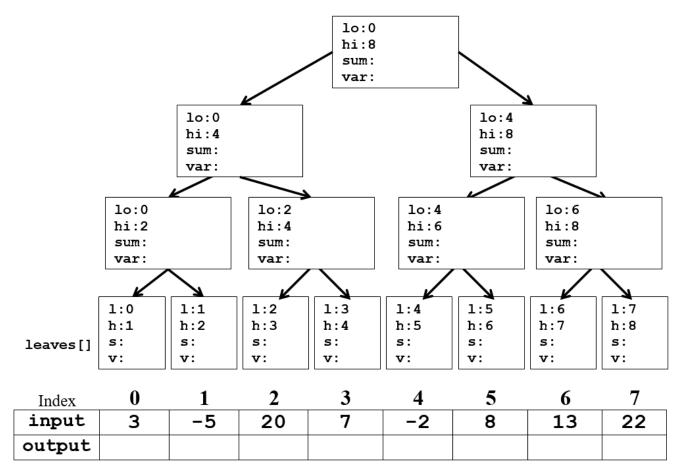
c) [4 pts] Modify the <u>code above in part b) and on the previous page</u> (draw arrows if needed) to *allow the most concurrent access* and to avoid all of the potential problems listed above. For full credit you must allow the most concurrent access possible without introducing any errors or extra locks. <u>You should use</u> synchronized appropriately to provide locking. Given the following array as input, perform a parallel **suffix** algorithm to fill the output array with the **sum of values contained in all of the cells to the** <u>right</u> (including the value contained in that cell) in the input array. The first pass of the algorithm is similar to the first pass of the parallel prefix code you have seen before. Do not use a sequential cutoff. For example, for

input: {3, 14, -1, 4, 5, 5, 6, 1},

output should be: {37, 34, 20, 21, 17, 12, 7, 1}.

Note: This question continues on the next page.

a) [5 pts] Fill in the values for sum, var, and the output array in the picture below. Note that problems b-e, on the next page, ask you to give the formulas you used in your calculation.



Give formulas for the following values where **p** is a reference to a non-leaf tree node and **leaves[i]** refers to the leaf node in the tree visible just above the corresponding location in the **input** and **output** arrays in the picture on the previous page.

b) [1 pt] Give pseudocode for how you assigned a value to leaves[i].sum

c) [1 pt] Give code for assigning p.left.var.

p.left.var =

d) [1 pt] Give code for assigning p.right.var.

p.right.var =

e) [1 pt] How is output[i] computed? Give <u>exact code</u> assuming <u>leaves[i]</u> refers to the leaf node in the tree visible just above the corresponding location in the <u>input</u> and output arrays in the picture above.

output[i] =

Q8: Sorting (9 pts)

(a-d) [4 pts] You are given an initial array: [22, 10, 14, 37, 14, 4, 3].

For each of the arrays shown below, indicate whether they could represent <u>one or more</u> of the given sorting algorithms at any point (after an iteration has completed) during the sorting algorithm's execution. If none of the sorting algorithms could have the given state, select **None**.

Note: For Radix sort, if you have following items in buckets 0-9: 0:[a], 1:[b, c], 2:[d], 3:[], 4:[], 5:[e, f], 6:[], 7:[g], 8:[], 9:[] we would write the resulting array as: [a, b, c, d, e, f, g]
a) [3, 4, 14, 37, 14, 10, 22]
b) [10, 22, 3, 14, 14, 37, 4] Insertion sort Selection sort Radix sort None
c) [10, 22, 3, 14, 14, 4, 37]
d) [10, 14, 14, 22, 37,4, 3] Insertion sort Selection sort Radix sort None
e) [1 pt] True or False: Heap sort is a stable sort. O True O False Explain why or why not in 1-2 sentences.

f) [2 pts] Give the <u>recurrence</u> for PARALLEL Merge Sort (parallel sort & sequential merge) - <u>best</u> <u>case</u>

T(n) =

g) [2 pts] Give the *recurrence* for SEQUENTIAL Quicksort - best case.

T(n) =

Q9: P/NP (6 pts)

- a) [1 pt] "NP" stands for _____
- b) [2 pt] Draw a diagram demonstrating how (we are pretty sure) the sets P, NP, and NP-Complete, overlap / don't overlap with each other.

For the following problems, bubble in <u>ALL</u> the sets each problem belongs to:

c) [1 pt] Given a weighted, undirected graph, determine if a minimum spanning tree of cost *k* exists.

NP-complete NP P none of these

d) [1 pt] Given a weighted, directed graph that has an edge between each vertex, determine whether there is a path of total cost *k* that begins and ends at the same vertex and goes through each **vertex** exactly once.

NP-complete	ПР	none of these

e) [1 pt] Given an undirected graph, determine whether there is a path that begins and ends at the same vertex and goes through each **edge** exactly once.

NP-complete	P	none of these
NP-complete	LΗP	lone of thes

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