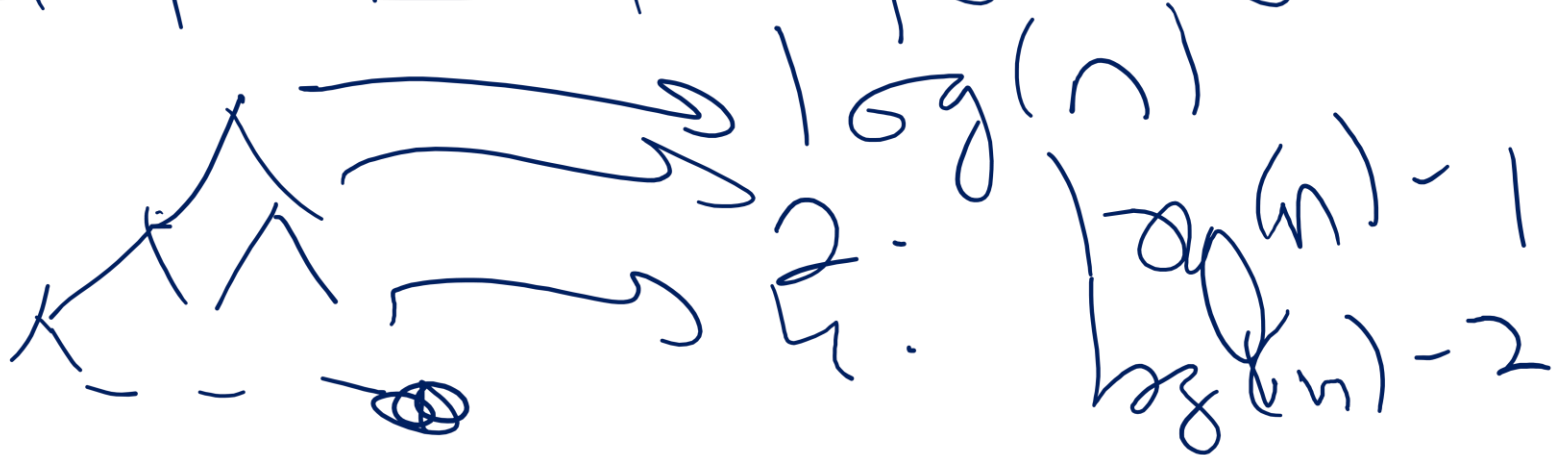
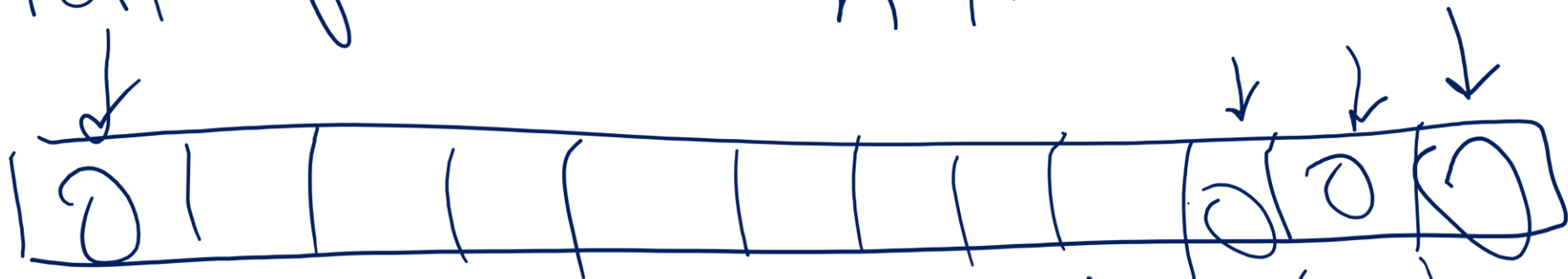


Heap's
BuildHeap

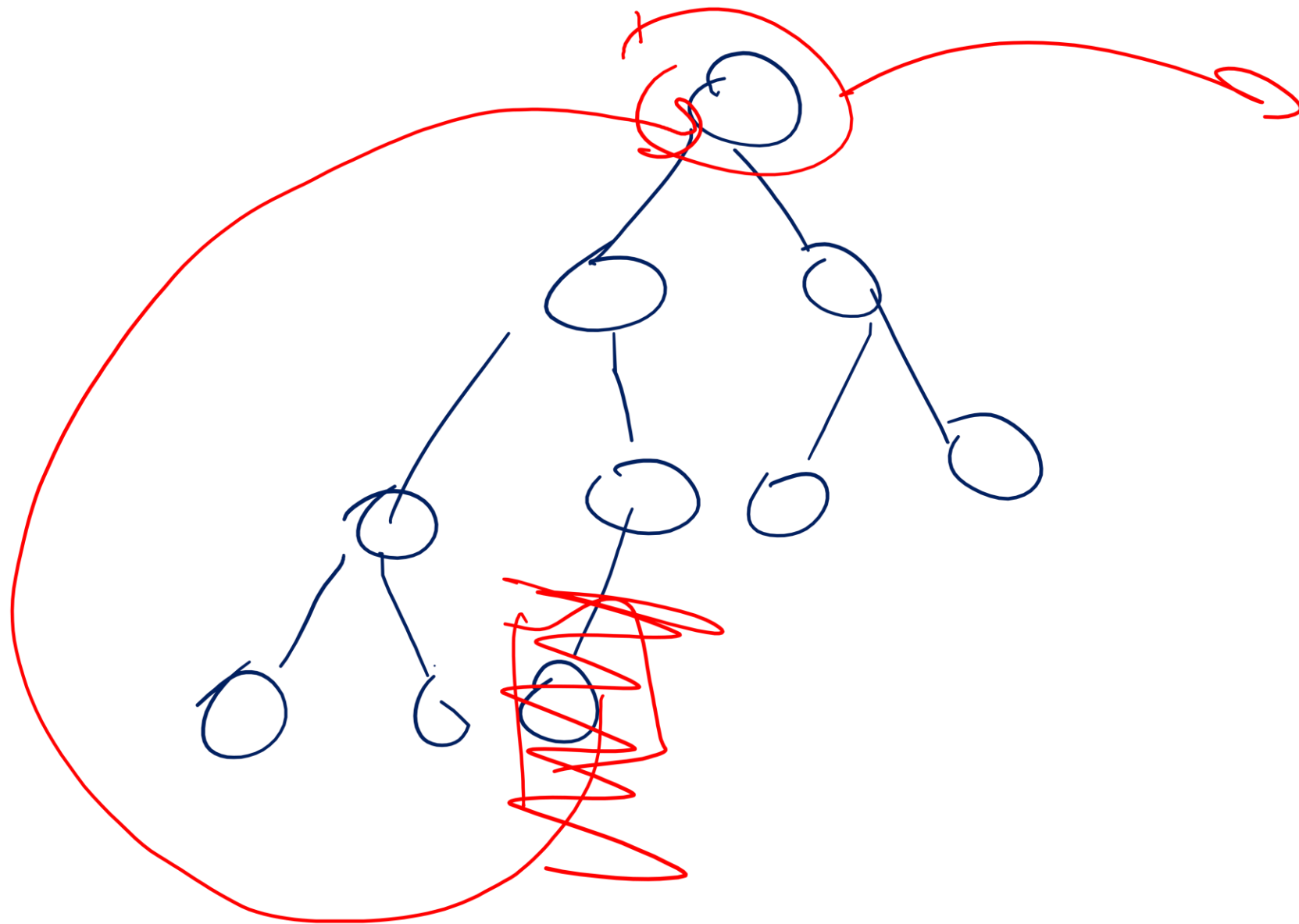
why is it

$O(n)$

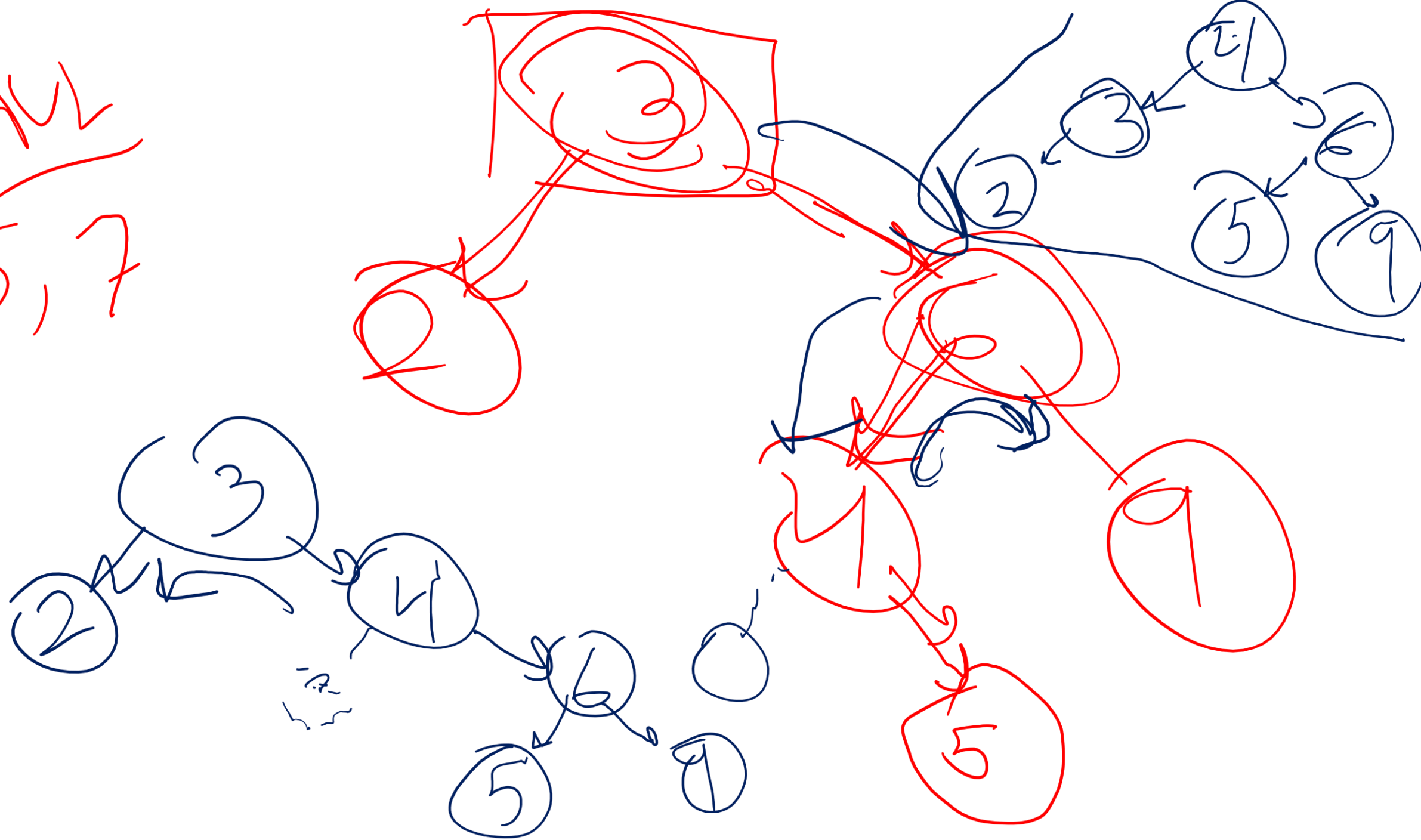
n items



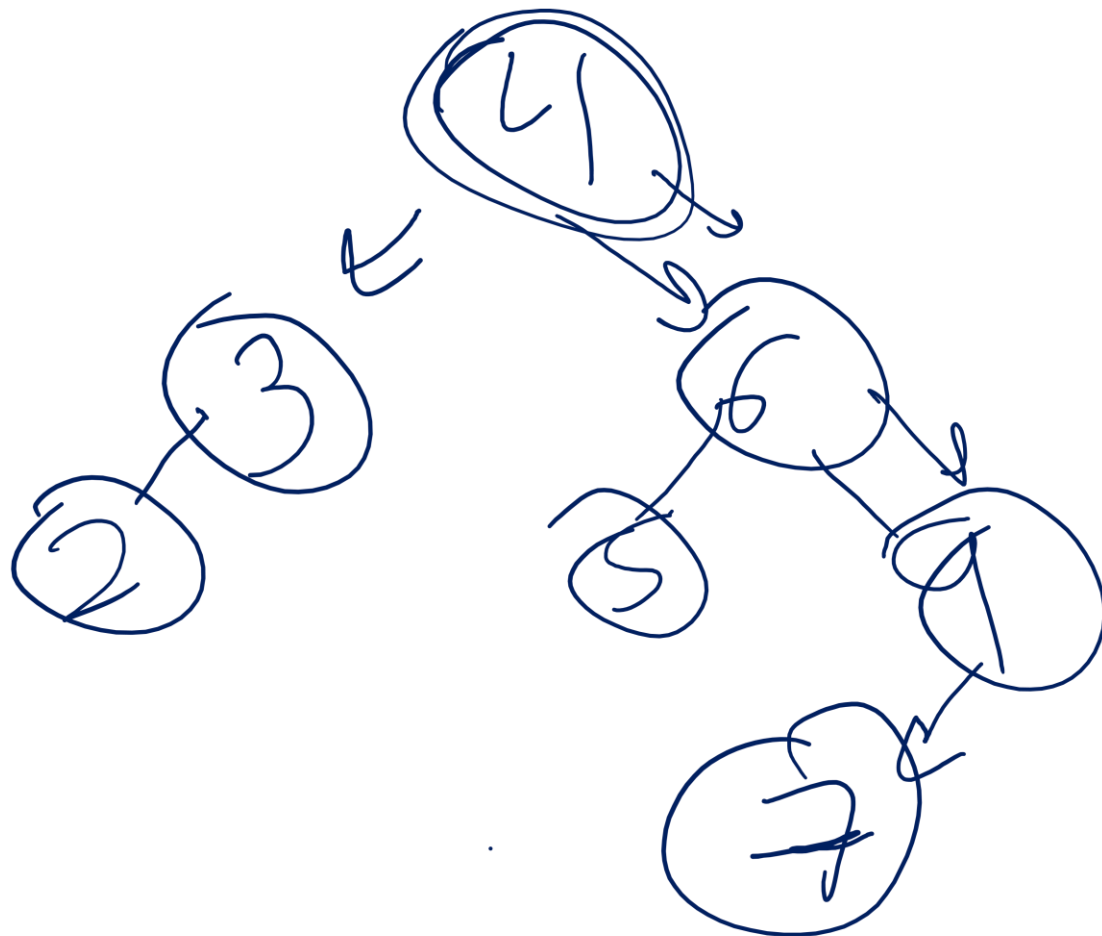
$\sim Cn$



AVL
5, 7



7



No handouts/slides.

Dictionaries $\leadsto$ Insert, find, delete

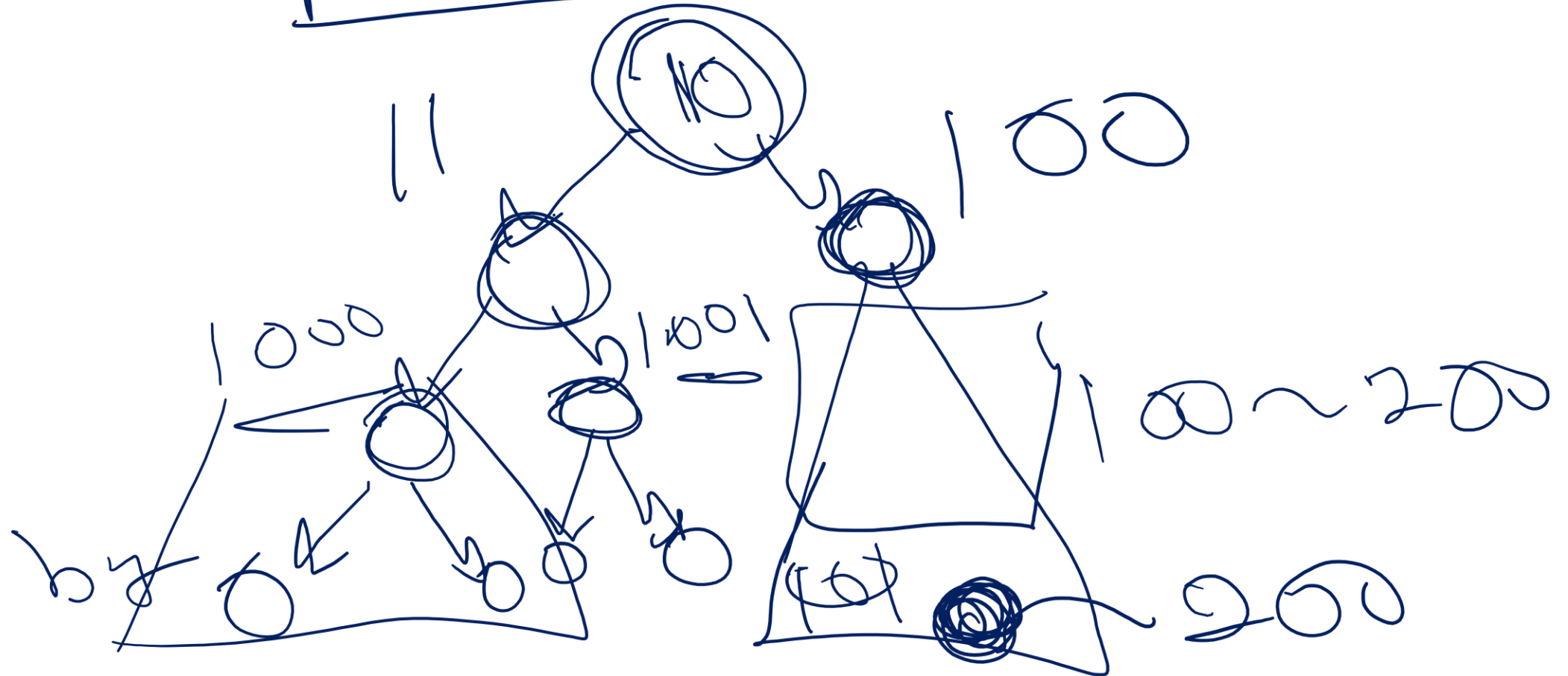
AVL trees
Hash tables

$$2^{\log_2(n)+1} = 2 \cdot 2^{\log_2(n)}$$

$$= 2 \cdot n$$

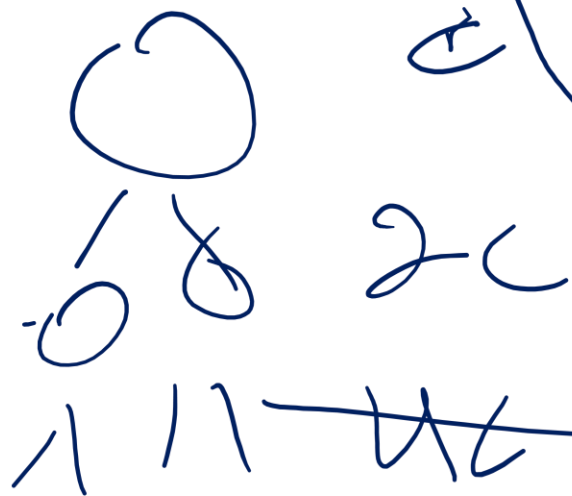
$\left\{ \begin{array}{l} \text{for } (i=0; i < n; i++) \\ \quad \text{for } (j=0; j < i; j++) \end{array} \right\}$

Remove Min heap best case?
→ repeated priorities



$$T(n) = \underbrace{2T\left(\frac{n}{2}\right) + 16}_{\text{recursion}} \quad \text{6/w}$$

~~50~~ 50 if $n < 10$



1.6C 2.1C

$$\underline{2^{\log_2(n)}} = \underline{n} \cdot C$$

for ($i=0; i < n; i = i * 2$)

for ($j=0; j < i; j++$)

$$\underbrace{1 + 2 + 4 + 8 + 16 + \dots + n}_{\log_2(n)} = \sum_{i=0}^{\log_2(n)-1} 2^i = \frac{2^{\log_2(n)+1} - 1}{2-1} = \underline{\underline{2^{n-1}}}$$

→ BSTs

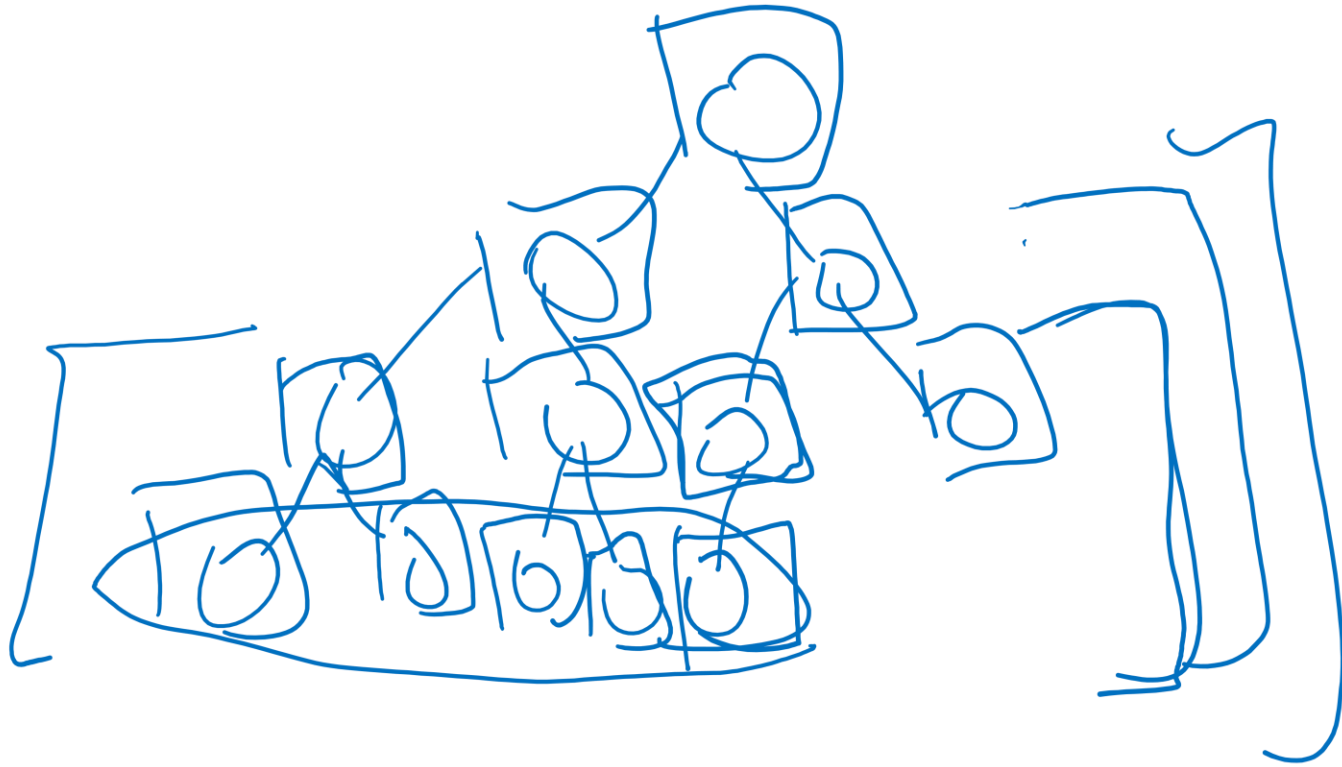
- Balance condition

$$\min \underline{N(h)} = \begin{cases} 1 & h=0 \\ 2 & h=1 \\ N(h-1) + N(h-2) + 1 & \text{o.w.} \end{cases}$$

$$\max \underline{P(h)} = \underline{2^{h+1} - 1}$$

build heap

input array int
output heap w/ those same priorities



$\sim C \cdot h$