

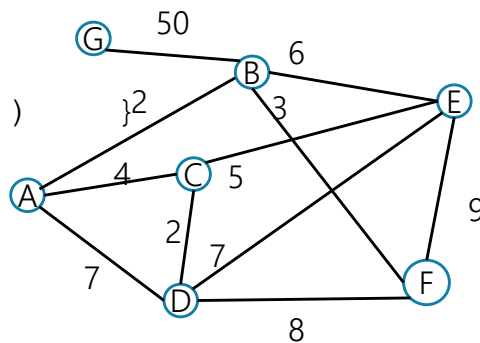
## Try it Out

KruskalMST(Graph G)

```

  initialize each vertex to be a connected component
  sort the edges by weight
  foreach(edge (u, v) in sorted order){
    if(find(u) != find(v)){
      add (u,v) to the MST
      union(find(u), find(v))
    }
  }

```



Operation	Worst-case Amortized
MakeSet()	$\Theta(1)$
Union()	$O(\log^* n)$
Find()	$O(\log^* n)$

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## P vs. NP

**P (stands for "Polynomial")**

The set of all decision problems that have an algorithm that runs in time  $O(n^k)$  for some constant  $k$ .

**NP (stands for "nondeterministic polynomial")**

The set of all decision problems such that if the answer is YES, there is a proof of that which can be verified in polynomial time.

Claim:  $P \subseteq NP$  (do you see why?)

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# EXP

EXP (stands for "Exponential")

The set of all decision problems that have an algorithm that runs in time  $O(2^{n^k})$  for some constant  $k$ .

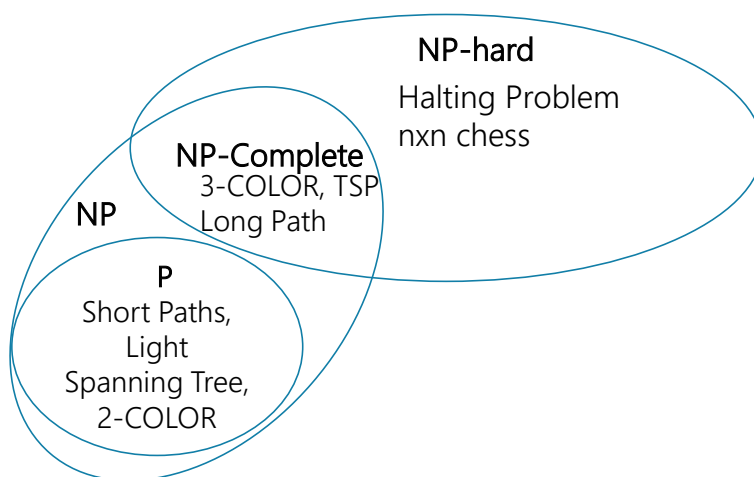
3-COLOR is in EXP (we just saw why on the last slide)

So is

Claim:  $NP \subseteq EXP$  (do you see why?)

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## What The World Looks Like (We Think)



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