Priority Queue ADT

Min Priority Queue ADT

state

Set of comparable values

- Ordered based on "priority" behavior

insert(value) – add a new element to the collection.
removeMin() – returns the element with the smallest priority, removes it from the collection.
peekMin() – find, but do not remove the element with the smallest priority.

Uses:

- Operating System
- Well-designed printers
- Some Compression Schemes (google Huffman Codes)
- Sorting
- · Graph algorithms

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Even More Operations

BuildHeap(elements $e_1, ..., e_n$) – Given n elements, create a heap containing exactly those n elements.

Try 1: Just call insert n times.

Worst case running time?

n calls, each worst case $\Theta(\log n)$. So it's $\Theta(n \log n)$ right?

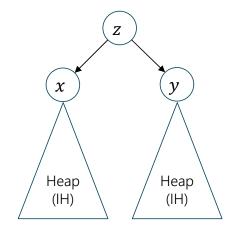
That proof isn't valid.

There are at least two distinct problems (bugs or gaps that need much more explanation), can you find them?

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Let's Prove It!

Well, let's sketch the proof of it.



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Amortization

AMORTIZED

It costs \$1800/month (which we pay once)

So the cost per day is $\frac{1800}{30} = 60$.

Good answer if the question is "what does my daily pay need to be to afford housing?"

UNAMORTIZED

On the first it costs \$1800.

Every other day of the month it costs \$0

Good answer if the question is "how much do I need to keep in my bank account so it doesn't get overdrawn?"

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