CSE 332 Winter 2024
Lecture 13: Hashing and Sorting

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Collision Resolution: Linear Probing

• When there’s a collision, use the next open space in the table
Linear Probing: Insert Procedure

- To insert $k, v$
  - Calculate $i = h(k) \% arrsize$
  - If $table[i]$ is occupied then try $(i + 1) \% arrsize$
  - If that is occupied try $(i + 2) \% arrsize$
  - If that is occupied try $(i + 3) \% arrsize$
  - ...
Linear Probing: Find

• \( i = h(k) \% \text{arrsize} \)
  • If \( i \) has the key or it’s empty, then we’re done
  • Otherwise:
    • Check \((i + 1) \% \text{arrsize}\) if it’s there, done else
    • Check \((i + 2) \% \text{arrsize}\) if it’s there, done else
    • Check \((i + 3) \% \text{arrsize}\)
    • ...
    • Until we hit an empty cell
Linear Probing: Find

• To find key \( k \)
  • Calculate \( i = h(k) \mod \text{arrsize} \)
  • If \( \text{table}[i] \) is occupied and does not contain \( k \) then look at \( (i + 1) \mod \text{arrsize} \)
  • If that is occupied and does not contain \( k \) then look at \( (i + 2) \mod \text{arrsize} \)
  • If that is occupied and does not contain \( k \) then look at \( (i + 3) \mod \text{arrsize} \)
  • Repeat until you either find \( k \) or else you reach an empty cell in the table
Linear Probing: Delete

• Problem: don’t want to leave an empty space when deleting
• Option 1: when we delete, move the “last thing” with a matching hash to that location
• Option 2: “tombstone” deletion. When deleting something, leave a special marker to indicate something used to be there
Linear Probing: Delete

• Option 1: Find the last thing with a matching hash, move that into the spot you deleted from

• Option 2: Called “tombstone” deletion. Leave a special object that indicates an object was deleted from there
  • The tombstone does not act as an open space when finding (so keep looking after its reached)
  • When inserting you can replace a tombstone with a new item
Downsides of Linear Probing

• What happens when $\lambda$ approaches 1?
  • Longer and longer clusters of items
  • Runnings times get longer and longer
Quadratic Probing: Insert Procedure

- To insert \( k, v \)
  - Calculate \( i = h(k) \mod arrsize \)
  - If \( table[i] \) is occupied then try \( (i + 1^2) \mod arrsize \)
  - If that is occupied try \( (i + 2^2) \mod arrsize \)
  - If that is occupied try \( (i + 3^2) \mod arrsize \)
  - If that is occupied try \( (i + 4^2) \mod arrsize \)
  - ...
Quadratic Probing: Example

• Insert:
  • 76
  • 40
  • 48
  • 5
  • 55
  • 47
Using Quadratic Probing

- If you probe `tablesize` times, you start repeating the same indices
- If `tablesize` is prime and $\lambda < \frac{1}{2}$ then you’re guaranteed to find an open spot in at most $\frac{tablesize}{2}$ probes

- Helps with the clustering problem of linear probing, but does not help if many things hash to the same value
Double Hashing: Insert Procedure

• Given $h$ and $g$ are both good hash functions

• To insert $k, v$
  • Calculate $i = h(k) \mod size$
  • If $table[i]$ is occupied then try $(i + g(k)) \mod size$
  • If that is occupied try $(i + 2 \cdot g(k)) \mod size$
  • If that is occupied try $(i + 3 \cdot g(k)) \mod size$
  • If that is occupied try $(i + 4 \cdot g(k)) \mod size$
  • ...
Rehashing

• If your load factor $\lambda$ gets too large, copy everything over to a larger hash table
  • To do this: make a new array with a new hash function (maybe just a new modulus)
  • Re-insert all items into the new hash table with the new hash function
  • New hash table should be “roughly” double the size (but probably still want it to be prime)

• General Guideline:
  • Separate Chaining: rehash when $\lambda = 2$
  • Open Addressing: rehash when $\lambda = \frac{1}{2}$
Sorting

- Rearrangement of items into some defined sequence
  - Usually: reordering a list from smallest to largest according to some metric

- Why sort things?
More Formal Definition

• Input:
  • An array $A$ of items
  • A comparison function for these items
    • Given two items $x$ and $y$, we can determine whether $x < y$, $x > y$, or $x = y$

• Output:
  • A permutation of $A$ such that if $i \leq j$ then $A[i] \leq A[j]$
  • Permutation: a sequence of the same items but perhaps in a different order
Sorting “Landscape”

• There is no singular best algorithm for sorting
• Some are faster, some are slower
• Some use more memory, some use less
• Some are super extra fast if your data matches particular assumptions
• Some have other special properties that make them valuable
• No sorting algorithm can have only all the “best” attributes
“Moving Day” Sorting Algorithm
Selection Sort

• Idea: Find the next smallest element, swap it into the next index in the array

```
1 2 3 4 5 6 10 8 7 9 12 11
```

Already In Position

```
1 2 3 4 5 6 7 8 10 9 12 11
```

Already In Position
Selection Sort

- Swap the thing at index 0 with the smallest thing in the array
- Swap the thing at index 1 with the smallest thing after index 0
- ...
- Swap the thing at index $i$ with the smallest thing after index $i - 1$

```java
for (i=0; i<a.length; i++){
    smallest = i;
    for (j=i; j<a.length; j++){
        if (a[j]<a[smallest]){ smallest=j;}
    }
    temp = a[i];
    a[i] = a[smallest];
    a[smallest] = a[i];
}
```

Running Time:
- Worst Case: $\Theta(n^2)$
- Best Case: $\Theta(n^2)$

<table>
<thead>
<tr>
<th>10</th>
<th>77</th>
<th>5</th>
<th>15</th>
<th>2</th>
<th>22</th>
<th>64</th>
<th>41</th>
<th>18</th>
<th>19</th>
<th>30</th>
<th>21</th>
<th>3</th>
<th>24</th>
<th>23</th>
<th>33</th>
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<td>6</td>
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<td>11</td>
<td>12</td>
<td>13</td>
<td>14</td>
<td>15</td>
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</table>
Insertion Sort

- Idea: Maintain a sorted list prefix, extend that prefix by “inserting” the next element.
Insertion Sort

- If the items at index 0 and 1 are out of order, swap them
- Keep swapping the item at index 2 with the thing to its left as long as the left thing is larger
- ...
- Keep swapping the item at index $i$ with the thing to its left as long as the left thing is larger

for (i=1; i<a.length; i++){
    prev = i-1;
    while(a[i] < a[prev] && prev > -1){
        temp = a[i];
        a[i] = a[prev];
        a[prev] = a[i];
        i--;
        prev--;
    }
}

Running Time:
- Worst Case: $\Theta(\cdot)$
- Best Case: $\Theta(\cdot)$

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Aside: Bubble Sort – we won’t cover it

"the bubble sort seems to have nothing to recommend it, except a catchy name and the fact that it leads to some interesting theoretical problems” –Donald Knuth, The Art of Computer Programming
Heap Sort

- **Idea**: Build a maxHeap, repeatedly delete the max element from the heap to build sorted list Right-to-Left

Max Heap Property: Each node is larger than its children
Heap Sort

- Remove the Max element (i.e. the root) from the Heap: replace with last element, call percolateDown(root)

Percolate Down(node): if node satisfies heap property, done. Else swap with largest child and repeat on that subtree
Heap Sort

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Heap Sort

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```
9 8 6 3 7 5 2 4 1
```

**Percolate Down(node):** if node satisfies heap property, done. Else swap with largest child and repeat on that subtree
Heap Sort

- Remove the Max element (i.e. the root) from the Heap: replace with last element, call percolateDown(root)

Percolate Down(node): if node satisfies heap property, done. Else swap with largest child and repeat on that subtree
Heap Sort

• Build a heap
• Call deleteMax
• Put that at the end of the array

myHeap = buildHeap(a);
for (int i = a.length-1; i>=0; i--){
    item = myHeap.deleteMax();
    a[i] = item;
}

Running Time:
  Worst Case: $\Theta(\cdot)$
  Best Case: $\Theta(\cdot)$
“In Place” Sorting Algorithm

- A sorting algorithm which requires no extra data structures
- Idea: It sorts items just by swapping things in the same array given
- Definition: it only uses $\Theta(1)$ extra space

- Selection sort: In Place!
- Insertion sort: In Place!
- Heap sort: Not In Place!
  - But we can fix that!
In Place Heap Sort

- **Idea**: When “removing” an element from the heap, swap it with the last item of the heap then “pretend” the heap is one item shorter
Heap Sort

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Heap Sort

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In Place Heap Sort

• Build a heap using the same array (Floyd’s build heap algorithm works)
• Call deleteMax
• Put that at the end of the array

```java
buildHeap(a);
for (int i = a.length-1; i>=0; i--){
    temp = a[i]
    a[i] = a[0];
    a[0] = temp;
    percolateDown(0);
}
```

Running Time:
- Worst Case: $\Theta(\cdot)$
- Best Case: $\Theta(\cdot)$
Floyd’s buildHeap method

• Working towards the root, one row at a time, percolate down

```java
buildHeap()
{
    for(int i = size; i>0; i--){
        percolateDown(i);
    }
}
```