CSE 332 Summer 2024
Lecture 12: Sorting

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http://www.cs.uw.edu/332
Selection Sort

• **Idea**: Find the next smallest element, swap it into the next index in the array
Selection Sort

• Swap the thing at index 0 with the smallest thing in the array
• Swap the thing at index 1 with the smallest thing after index 0
• …
• Swap the thing at index $i$ with the smallest thing after index $i - 1$

for (i=0; i<a.length; i++){  
    smallest = i;  
    for (j=i; j<a.length; j++){  
        if (a[j]<a[smallest]){ smallest=j;}
    }  
    temp = a[i];  
a[i] = a[smallest];  
a[smallest] = a[i];
}

Running Time:
Worst Case: $\Theta(\cdot)$
Best Case: $\Theta(\cdot)$

<table>
<thead>
<tr>
<th>10</th>
<th>77</th>
<th>5</th>
<th>15</th>
<th>2</th>
<th>22</th>
<th>64</th>
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<th>18</th>
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<tbody>
<tr>
<td>0</td>
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<td>11</td>
<td>12</td>
<td>13</td>
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</table>
Insertion Sort

- **Idea:** Maintain a sorted list prefix, extend that prefix by “inserting” the next element.
Insertion Sort

• If the items at index 0 and 1 are out of order, swap them
• Keep swapping the item at index 2 with the thing to its left as long as the left thing is larger
• ...
• Keep swapping the item at index i with the thing to its left as long as the left thing is larger

for (i=1; i<a.length; i++){
    prev = i-1;
    while(a[i] < a[prev] && prev > -1){
        temp = a[i];
        a[i] = a[prev];
        a[prev] = temp;
        i--;
        prev--;
    }
}

Running Time:

Worst Case: \( \Theta(n^2) \)
Best Case: \( \Theta(n) \)
Aside: Bubble Sort – we won’t cover it

"the bubble sort seems to have nothing to recommend it, except a catchy name and the fact that it leads to some interesting theoretical problems” –Donald Knuth, The Art of Computer Programming
Heap Sort

- **Idea**: Build a maxHeap, repeatedly delete the max element from the heap to build sorted list Right-to-Left

<table>
<thead>
<tr>
<th>10</th>
<th>9</th>
<th>6</th>
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<th>2</th>
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Max Heap Property: Each node is larger than its children
Heap Sort

- Remove the Max element (i.e. the root) from the Heap: replace with last element, call percolateDown(root)

Percolate Down(node): if node satisfies heap property, done. Else swap with largest child and repeat on that subtree.
Heap Sort

- Remove the Max element (i.e. the root) from the Heap: replace with last element, call percolateDown(root)

Percolate Down(node): if node satisfies heap property, done. Else swap with largest child and repeat on that subtree.
Heap Sort

• **Remove the Max element (i.e. the root) from the Heap:** replace with last element, call percolateDown(root)

```
9 8 6 3 7 5 2 4 1
```

**Percolate Down(node):** if node satisfies heap property, done. Else swap with largest child and repeat on that subtree
Heap Sort

- Remove the Max element (i.e. the root) from the Heap: replace with last element, call percolateDown(root)

**Percolate Down(node):** if node satisfies heap property, done. Else swap with largest child and repeat on that subtree
Heap Sort

• Build a heap
• Call deleteMax
• Put that at the end of the array

myHeap = buildHeap(a);
for (int i = a.length-1; i>=0; i--){
    item = myHeap.deleteMax();
    a[i] = item;
}

Running Time:
  Worst Case: \(\Theta(\cdot)\)
  Best Case: \(\Theta(\cdot)\)
“In Place” Sorting Algorithm

• A sorting algorithm which requires no extra data structures
• Idea: It sorts items just by swapping things in the same array given
• Definition: it only uses $\Theta(1)$ extra space

• Selection sort: In Place!
• Insertion sort: In Place!
• Heap sort: Not In Place!
  • But we can fix that!
In Place Heap Sort

- **Idea:** When “removing” an element from the heap, swap it with the last item of the heap then “pretend” the heap is one item shorter
Heap Sort

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In Place Heap Sort

- Build a heap using the same array (Floyd’s build heap algorithm works)
- Call deleteMax
- Put that at the end of the array

```java
buildHeap(a);
for (int i = a.length-1; i>=0; i--){
    temp=a[i]
    a[i] = a[0];
    a[0] = temp;
    percolateDown(0);
}
```

Running Time:
- Worst Case: $\Theta(\cdot)$
- Best Case: $\Theta(\cdot)$
Floyd’s buildHeap method

• Working towards the root, one row at a time, percolate down

buildHeap()
{
    for(int i = size; i>0; i--){
        percolateDown(i);
    }
}
Merge Sort

- **Base Case:**
  - If the list is of length 1 or 0, it’s already sorted, so just return it

- **Divide:**
  - Split the list into two “sublists” of (roughly) equal length

- **Conquer:**
  - Sort both lists recursively

- **Combine:**
  - **Merge** sorted sublists into one sorted list
Merge Sort In Action!

Sort between indices $low$ and $high$

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Base Case: if $low == high$ then that range is already sorted!

Divide and Conquer: Otherwise call mergesort on ranges $\left( low, \frac{low+high}{2} \right)$ and $\left( \frac{low+high}{2} + 1, high \right)$

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After Recursion:

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Create a new array to merge into, and 3 pointers/indices:

- **L_next**: the smallest “unmerged” thing on the left
- **R_next**: the smallest “unmerged” thing on the right
- **M_next**: where the next smallest thing goes in the merged array

One-by-one: put the smallest of L_next and R_next into M_next, then advance both M_next and whichever of L/R was used.
Merge Sort Pseudocode

void mergesort(myArray)
{
    ms_helper(myArray, 0, myArray.length());
}

void mshelper(myArray, low, high)
{
    if (low == high){return;}  // Base Case
    mid = (low+high)/2;
    ms_helper(low, mid);
    ms_helper(mid+1, high);
    merge(myArray, low, mid, high);
}
void merge(myArray, low, mid, high) {
    merged = new int[high-low+1]; // or whatever type is in myArray
    l_next = low;
    r_next = high;
    m_next = 0;
    while (l_next <= mid && r_next <= high) {
        if (myArray[l_next] <= myArray[r_next]) {
            merged[m_next++] = myArray[l_next++];
        }
        else{
            merged[m_next++] = myArray[r_next++];
        }
    }
    while (l_next <= mid){ merged[m_next++] = myArray[l_next++]; }
    while (r_next <= high){  merged[m_next++] = myArray[r_next++]; }
    for(i=0; i<=merged.length; i++) { myArray[i+low] = merged[i];}
}
Analyzing Merge Sort

1. Identify time required to Divide and Combine
2. Identify all subproblems and their sizes
3. Use recurrence relation to express recursive running time
4. Solve and express running time asymptotically

• **Divide**: 0 comparisons
• **Conquer**: recursively sort two lists of size $\frac{n}{2}$
• **Combine**: $n$ comparisons
• **Recurrence:**
  
  \[ T(n) = 0 + T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + n \]
  
  \[ T(n) = 2T\left(\frac{n}{2}\right) + n \]
\[ T(n) = 2T\left(\frac{n}{2}\right) + n \]

Red box represents a problem instance
Blue value represents time spent at that level of recursion

\[ \Rightarrow n \text{ comparisons / level} \]

\[ \log_2 n \text{ levels of recursion} \]

\[ T(n) = \sum_{i=1}^{\log_2 n} n = n \log_2 n \]
Properties of Merge Sort

• Worst Case Running time:
  • $\Theta(n \log n)$

• In-Place?
  • No!

• Adaptive?
  • No!

• Stable?
  • Yes!
    • As long as in a tie you always pick $l_{\text{next}}$
Quicksort

• Like Mergesort:
  • Divide and conquer
  • $O(n \log n)$ run time (kind of…)

• Unlike Mergesort:
  • Divide step is the “hard” part
  • *Typically* faster than Mergesort
Quicksort

Idea: pick a *pivot* element, recursively sort two sublists around that element

- **Divide**: select *pivot* element $p$, $\text{Partition}(p)$
- **Conquer**: recursively sort left and right sublists
- **Combine**: Nothing!
Partition (Divide step)

Given: a list, a pivot $p$

Start: unordered list

| 8 | 5 | 7 | 3 | 12 | 10 | 1 | 2 | 4 | 9 | 6 | 11 |

Goal: All elements $< p$ on left, all $> p$ on right

| 5 | 7 | 3 | 1 | 2 | 4 | 6 | 8 | 12 | 10 | 9 | 11 |
Partition, Procedure

If Begin value < \( p \), move Begin right
Else swap Begin value with End value, move End Left
Done when Begin = End
Partition, Procedure

If Begin value < \( p \), move Begin right
Else swap Begin value with End value, move End Left
Done when Begin = End
Partition, Procedure

If \texttt{Begin} value $< p$, move \texttt{Begin} right
Else swap \texttt{Begin} value with \texttt{End} value, move \texttt{End} Left
Done when \texttt{Begin} = \texttt{End}

Case 1: meet at element $< p$
Swap $p$ with pointer position (2 in this case)
Partition, Procedure

If Begin value < \( p \), move Begin right
Else swap Begin value with End value, move End Left
Done when Begin = End

Case 2: meet at element > \( p \)

Swap \( p \) with value to the left (2 in this case)
Partition Summary

1. Put $p$ at beginning of list
2. Put a pointer (Begin) just after $p$, and a pointer (End) at the end of the list
3. While Begin < End:
   1. If Begin value $< p$, move Begin right
   2. Else swap Begin value with End value, move End Left
4. If pointers meet at element $< p$: Swap $p$ with pointer position
5. Else If pointers meet at element $> p$: Swap $p$ with value to the left

Run time? $O(n)$
Conquer

2 5 7 3 6 4 1

All elements < \( p \)

Exactly where it belongs!

8

All elements > \( p \)

10 9 11 12

Recursively sort Left and Right sublists
Quicksort Run Time (Best)

If the **pivot** is always the median:

| 2 | 5 | 1 | 3 | 6 | 4 | 7 | 8 | 10 | 9 | 11 | 12 |

Then we divide in half each time

\[
T(n) = 2T\left(\frac{n}{2}\right) + n
\]

\[
T(n) = O(n \log n)
\]
Quicksort Run Time (Worst)

If the pivot is always at the extreme:

\[
T(n) = T(n - 1) + n
\]

Then we shorten by 1 each time

\[
T(n) = O(n^2)
\]
Quicksort Run Time (Worst)

\[ T(n) = T(n - 1) + n \]

\[ T(n) = 1 + 2 + 3 + \cdots + n \]

\[ T(n) = \frac{n(n + 1)}{2} \]

\[ T(n) = O(n^2) \]
Quicksort on a (nearly) Sorted List

First element always yields unbalanced pivot

So we shorten by 1 each time

$$T(n) = T(n-1) + n$$

$$T(n) = O(n^2)$$
Good Pivot

• What makes a good Pivot?
  • Roughly even split between left and right
  • Ideally: median

• There are ways to find the median in linear time, but it’s complicated and slow and you’re better off using mergesort

• In Practice:
  • Pick a random value as a pivot
  • Pick the middle of 3 random values as the pivot
Properties of Quick Sort

• Worst Case Running time:
  • $\Theta(n^2)$
  • But $\Theta(n \log n)$ average! And typically faster than mergesort!

• In-Place?
  • ....Debatable

• Adaptive?
  • No!

• Stable?
  • No!
Improving Running time

• Recall our definition of the sorting problem:
  • Input:
    • An array $A$ of items
    • A comparison function for these items
      • Given two items $x$ and $y$, we can determine whether $x < y$, $x > y$, or $x = y$
  • Output:
    • A permutation of $A$ such that if $i \leq j$ then $A[i] \leq A[j]$
• Under this definition, it is impossible to write an algorithm faster than $n \log n$ asymptotically.

• Observation:
  • Sometimes there might be ways to determine the position of values without comparisons!
“Linear Time” Sorting Algorithms

• Useable when you are able to make additional assumptions about the contents of your list (beyond the ability to compare)
  • Examples:
    • The list contains only positive integers less than $k$
    • The number of distinct values in the list is much smaller than the length of the list

• The running time expression will always have a term other than the list’s length to account for this assumption
  • Examples:
    • Running time might be $\Theta(k \cdot n)$ where $k$ is the range/count of values
BucketSort

- Assumes the array contains integers between 0 and $k - 1$ (or some other small range)

- Idea:
  - Use each value as an index into an array of size $k$
  - Add the item into the “bucket” at that index (e.g. linked list)
  - Get sorted array by “appending” all the buckets
BucketSort Running Time

• Create array of $k$ buckets
  • Either $\Theta(k)$ or $\Theta(1)$ depending on some things...
• Insert all $n$ things into buckets
  • $\Theta(n)$
• Empty buckets into an array
  • $\Theta(n + k)$
• Overall:
  • $\Theta(n + k)$
• When is this better than mergesort?
Properties of BucketSort

• In-Place?
  • No

• Adaptive?
  • No

• Stable?
  • Yes!
RadixSort

• Radix: The base of a number system
  • We’ll use base 10, most implementations will use larger bases

• Idea:
  • BucketSort by each digit, one at a time, from least significant to most significant

Place each element into a “bucket” according to its 1’s place
RadixSort

• Radix: The base of a number system
  • We’ll use base 10, most implementations will use larger bases

• Idea:
  • BucketSort by each digit, one at a time, from least significant to most significant

<table>
<thead>
<tr>
<th>800</th>
<th>801</th>
<th>401</th>
<th>101</th>
<th>901</th>
<th>121</th>
<th>512</th>
<th>103</th>
<th>323</th>
<th>823</th>
<th>113</th>
<th>255</th>
<th>555</th>
<th>245</th>
<th>018</th>
<th>999</th>
</tr>
</thead>
</table>

Place each element into a “bucket” according to its 10’s place
RadixSort

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• Idea:
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Place each element into a “bucket” according to its 100’s place
RadixSort

• Radix: The base of a number system
  • We’ll use base 10, most implementations will use larger bases

• Idea:
  • BucketSort by each digit, one at a time, from least significant to most significant

```
 018 101 103 113 121 245 255 323 401 512 555 800 801 823 901 999
```

Convert back into an array
RadixSort Running Time

• Suppose largest value is $m$
• Choose a radix (base of representation) $b$
• BucketSort all $n$ things using $b$ buckets
  • $\Theta(n + k)$
• Repeat once per each digit
  • $\log_b m$ iterations
• Overall:
  • $\Theta(n \log_b m + b \log_b m)$
• In practice, you can select the value of $b$ to optimize running time
• When is this better than mergesort?