

# 23su CSE332 Final 1

Full Name:

Email Address (UW NetID):

@uw.edu

## Instructions:

- The allotted time is 1 hour.
- Do not turn the page until the staff says to do so.
- Read the directions carefully, especially for problems that require you to show work or provide an explanation.
- This is a closed-book and closed-notes exam.
- You are NOT permitted to access electronic devices including calculators.
- You must put your final answer inside the box.
  - If you run out of space, indicate where the answer continues.
  - Try to avoid writing on the very edges of the pages as we scan the exams.
- Unless otherwise noted, any bounds must be the **worst-case**, **simplified** and **tight**.
- Unless otherwise noted, logs are base 2.
- Unless otherwise noted, all material is assumed as in lecture.
- For answers that involve bubbling in a ☐ or ☐, fill in the shape completely.
- A formula sheet has been included at the end of the exam.

## Advice:

- If you feel like you're stuck on a problem, you may want to skip it and come back at the end if you have time.
- Look at the question titles on the cover page to see if you want to start somewhere other than problem 1.
- **Relax and take a few deep breaths. You've got this! :-).**

<b>Q1: Short Answer Questions (17 pts)</b> .....	<b>2</b>
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## Q1: Short Answer Questions (17 pts)

- $O$ ,  $\Omega$ , or  $\Theta$  bound must be **simplified** and **tight**. This means that, for example,  $7n + 3 \in O(7n + 3)$  (not simplified) or  $n \in O(2^{n!})$  (not tight enough) are unlikely to get points.
- Unless otherwise specified, all logs are base 2.
- For questions with a mathematical answer, you may leave your answer as an unsimplified formula (e.g.,  $7 \cdot 103$ ).
- Points are all or nothing (i.e., we will only grade what is inside the box).

a) (2 pts) True or False: Even when a `QuadraticProbingHashTable` has a prime number `TableSize`, it's still possible for cycling to occur.

☐ True
 ☐ False

b) (2 pts) Give the worst-case runtime of a delete operation on a `ChainingHashTable` containing  $N$  elements, where each bucket is an unsorted linked list. The load factor of this table is  $\log M$ .

 $O(\text{[ ]})$ 

c) (2 pts) Give the worst-case runtime of inserting  $N$  elements in an empty `ChainingHashTable` where each bucket is an unsorted linked list.

 $O(\text{[ ]})$ 

d) (1 pt) As in lecture, give a stable, **not** in-place,  $n \log n$  sorting algorithm.

e) (2 pts) True or False: A `Queue` (instead of a `Stack`) can be used to implement DFS.

☐ True
 ☐ False

- f) (2 pts) If 30% of your program must be run sequentially (and the rest can be parallelized), what is the maximum speedup you would expect to get with infinite processors?

times

Bonus Question (0 pts): Who is credited with this law?

- g) (1 pt) Give the worst-case span of parallel merge sort (as in lecture, using parallel sort and merge).

$O(\text{  )$

- h) (2 pts) True or False: Using a reentrant lock for read-only operations will eliminate all possibilities of potential data race situations.

☐ True
☐ False

- i) (2 pts) True or False: The topological sort algorithm can be run on a **cyclic** graph without giving an incorrect topological ordering.

☐ True
☐ False

- j) (1 pt) True or False: For Kruskal's algorithm, does processing the edges in ascending order of weight produce the correct results?

☐ True
☐ False

## Q2: Hashing (12 pts)

### a) (6 pts) Double Hashing

Consider the following hash table of `TableSize = 10` using double hashing with the following hash functions:

$$h(key) = key \% 10$$

$$g(key) = 7 - (key \% 7)$$

- i) (1 pt) Insert the numbers 23, 45, 13, 29, 37, 9, 19 into the table. It is possible that some elements cannot be inserted.

0	
1	
2	
3	
4	
5	
6	
7	
8	
9	

- ii) (1 pt) If any elements cannot be inserted, write them here:

- iii) (1 pt) What is the load factor of the above hash table? Write your answer as a precise fraction without simplifying as per the definition.

## b) (6 pts) Quadratic Probing

Consider the following hash table of `TableSize = 10` using quadratic probing and the hash function  $h(key) = key \% 10$

- i) (2 pts) Using lazy deletion, delete the numbers 91, 3, 31, 66 from the hash table on the left (then put the new hash table after deletion on the right):

0	11
1	91
2	12
3	3
4	23
5	31
6	66
7	
8	
9	

0	
1	
2	
3	
4	
5	
6	
7	
8	
9	

- ii) (2 pts) Suppose that we make the following change to lazy deletion within a table using quadratic probing. When deleting an element, if the next slot in the probe sequence is empty, then we can mark the slot as 'empty' instead of 'deleted'.

True or False: Future search operations will work.

☐ True ☐ False

- iii) (2 pts) Consider Device X, a small computer with limited memory. As the designer of Device X, you are faced with the challenge of incorporating a hash table to manage key-value pairs efficiently. Device X must handle numerous key-value pairs reliably, but high-speed performance is not a critical requirement.

Given this scenario, choose the best hash table type.

☐ Separate ☐ Linear ☐ Quadratic

### Q3: Sorting (12 pts)

You are given an initial array:  $[50, 42_1, 42_2, 37, 34]$ .

For each of the arrays shown below, indicate whether they could represent **one or more** of the given sorting algorithms **at any point** (after an iteration has completed) during the sorting algorithm's execution. If none of the sorting algorithms could have the given state, select **None**.

Assume all algorithms work exactly as in lecture.

Note: For Radix sort, if you have following items in buckets 0-9:

0:[a], 1:[b, c], 2:[d], 3:[], 4:[], 5:[e, f], 6:[], 7:[g], 8:[], 9:[]  
we would write the resulting array as: [a, b, c, d, e, f, g].

a) (3 pts)  $[34, 42_1, 42_2, 37, 50]$

<input type="radio"/> Insertion Sort	<input type="radio"/> Selection sort	<input type="radio"/> Radix sort	<input type="radio"/> None
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b) (3 pts)  $[34, 37, 42_1, 42_2, 50]$

<input type="radio"/> Insertion Sort	<input type="radio"/> Selection sort	<input type="radio"/> Radix sort	<input type="radio"/> None
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c) (3 pts)  $[50, 42_1, 42_2, 34, 37]$

<input type="radio"/> Insertion Sort	<input type="radio"/> Selection sort	<input type="radio"/> Radix sort	<input type="radio"/> None
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d) (3 pts)  $[42_1, 42_2, 50, 37, 34]$

<input type="radio"/> Insertion Sort	<input type="radio"/> Selection sort	<input type="radio"/> Radix sort	<input type="radio"/> None
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### Q4: Graphs (18 pts)

- a) (2 pts) Given a **Dense, Directed** Graph, which graph implementation would be preferable for computing the reverse graph (i.e. a graph where all edges are reversed)?

- b) (2 pts) Give a simplified, tight bound for the runtime of computing the reverse graph. Express your answer in terms of  $V$  (number of vertices) only.

$O(\text{ })$

- c) (2 pts) Give an exact number for the maximum number of edges in an **Undirected, Acyclic** graph. Express your answer in terms of  $V$  (number of vertices) only.

- d) (2 pts) A WinstonHeap has these worst-case runtimes:

insert -  $O(1)$   
 deleteMin -  $O(1)$   
 decreaseKey -  $O(1)$

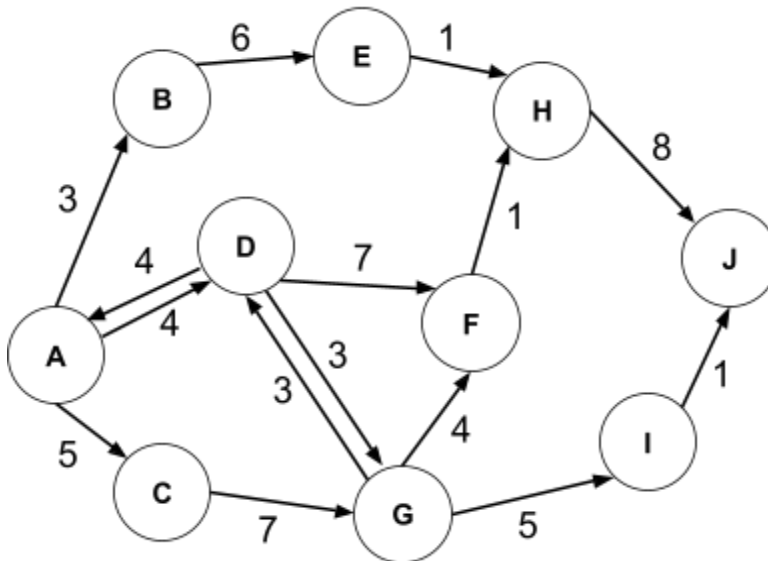
Give the worst-case runtime for Dijkstra's algorithm using a WinstonHeap. Express your answer in terms of  $V$  (number of vertices) only.

$O(\text{ })$

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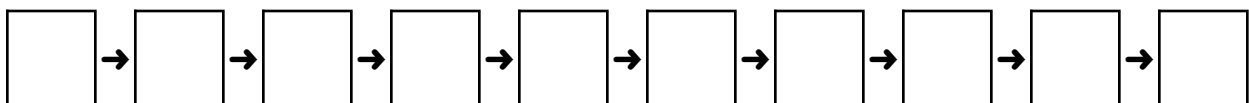
e) Dijkstra's



- i) With the graph above, step through Dijkstra's Algorithm to calculate the **single source shortest path from A** to every other vertex. Break ties by choosing the lexicographically (alphabetically) smallest letter first (i.e., if B and C were tied, you would explore B first). Note that later on, you will need to recall what order vertices were declared *visited*. For full credit, you should only show the final values (use scratch paper or erase old values).

Node	Visited?	Cost	Predecessor
A			
B			
C			
D			
E			
F			
G			
H			
I			
J			

- ii) In what order would Dijkstra's algorithm mark each node as *visited*?



iii) List the shortest path from A to J (Give the actual path, not the cost).



## CSE 332: Data Structures and Parallelism

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### Useful Math Identities

#### Summations

1.  $\sum_{i=0}^{\infty} x^i = \frac{1}{1-x}$  for  $|x| < 1$
2.  $\sum_{i=0}^{n-1} 1 = \sum_{i=1}^n 1 = n$
3.  $\sum_{i=0}^n i = 0 + \sum_{i=1}^n i = \frac{n(n+1)}{2}$
4.  $\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{n^3}{3} + \frac{n^2}{2} + \frac{n}{6}$
5.  $\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4}$
6.  $\sum_{i=0}^{n-1} x^i = \frac{1-x^n}{1-x}$
7.  $\sum_{i=0}^{n-1} \frac{1}{2^i} = 2 - \frac{1}{2^{n-1}}$

#### Logs

1.  $x^{\log_x n} = n$
2.  $a^{\log_b c} = c^{\log_b a}$
3.  $\log_b a = \frac{\log_d a}{\log_d b}$

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