



# CSE 332: Data Structures & Parallelism

## Lecture 22: Minimum Spanning Trees

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# Minimum Spanning Trees

Given an undirected graph  $G=(V,E)$ , find a graph  $G'=(V, E')$  such that:

- $E'$  is a subset of  $E$
- $|E'| = |V| - 1$
- $G'$  is connected

**$G'$  is a minimum spanning tree.**

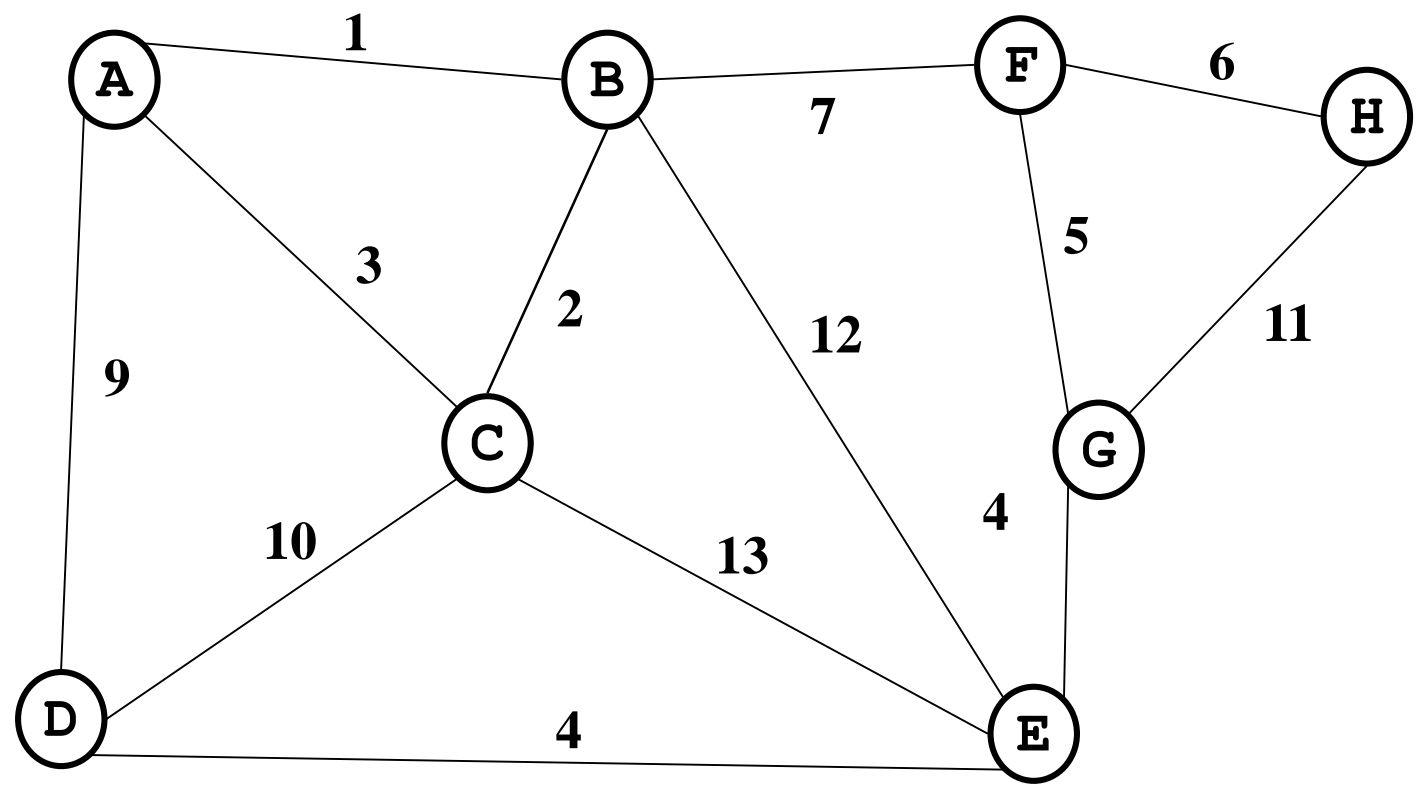
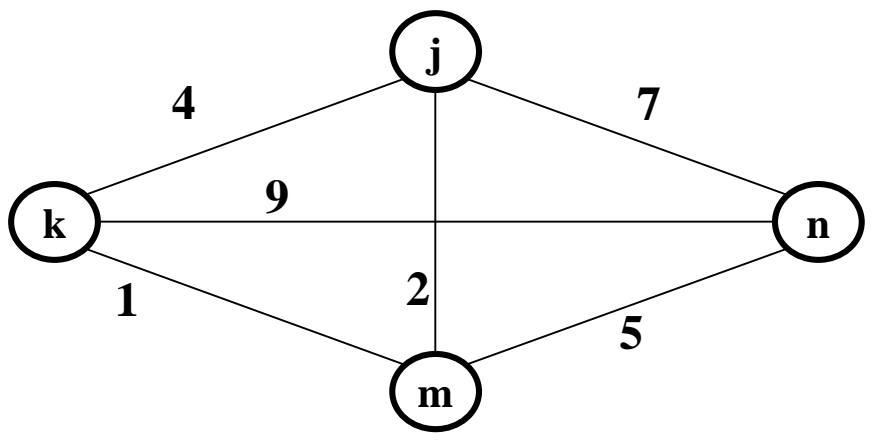
- $\sum_{(u,v) \in E'} c_{uv}$  is minimal

## Applications:

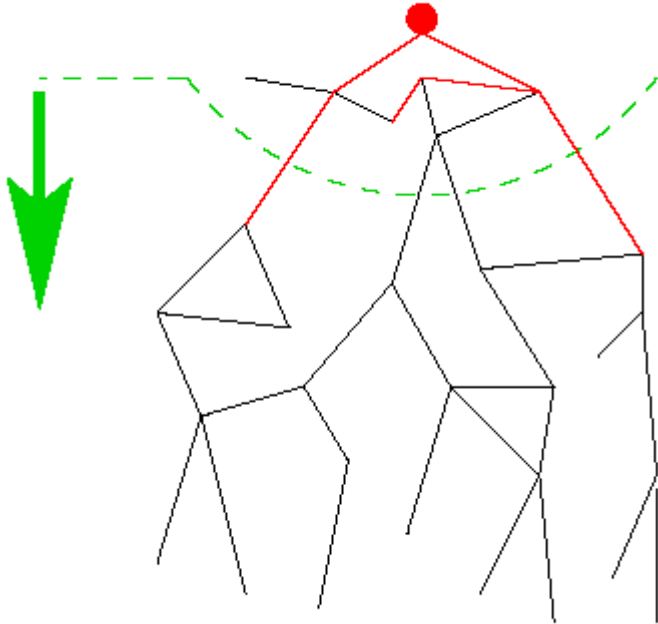
- Example: Electrical wiring for a house or clock wires on a chip
- Example: A road network if you cared about asphalt cost rather than travel time

# Student Activity

*Find the MST*

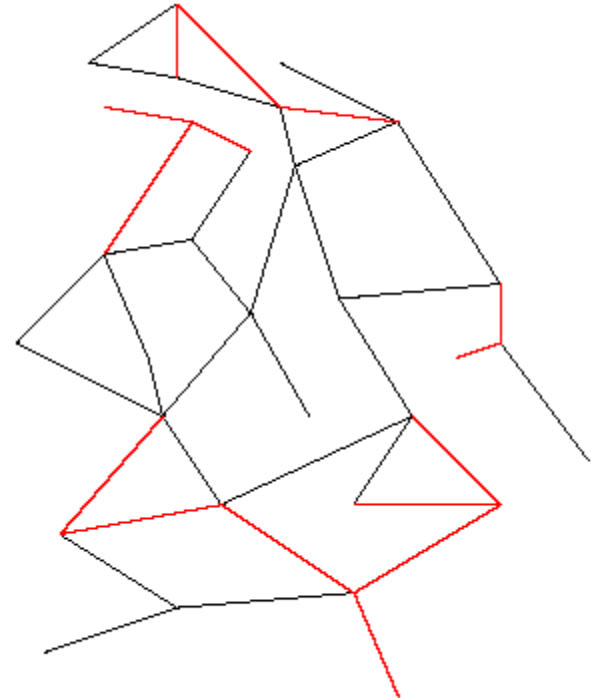


# Two Different Approaches



**Prim's Algorithm**

**Almost identical to Dijkstra's**



**Kruskal's Algorithm**

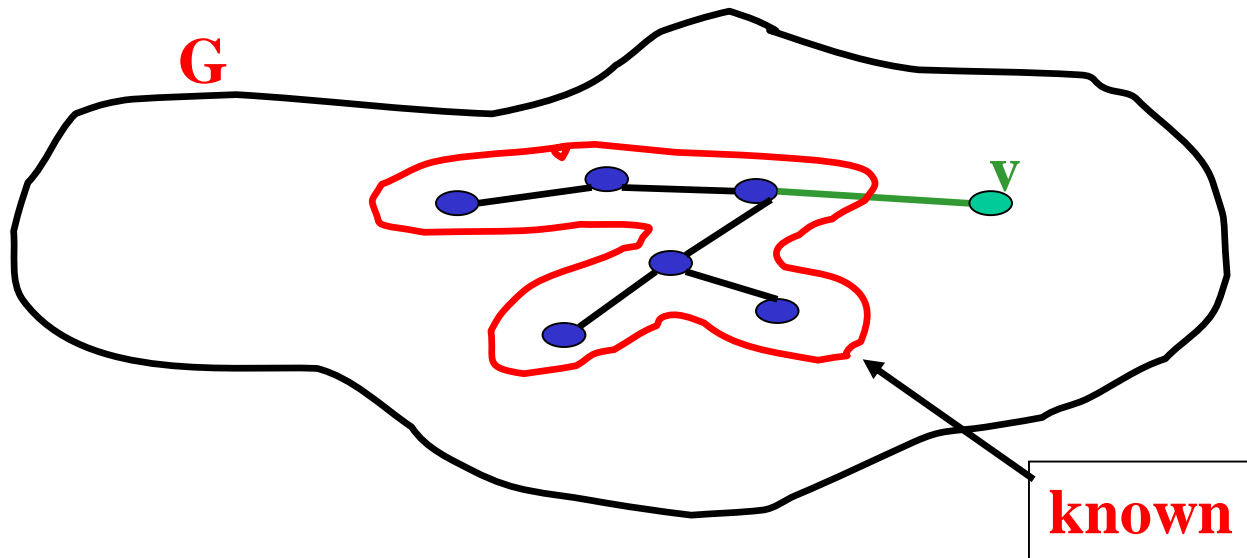
**Completely different!**

# Prim's algorithm

**Idea:** Grow a tree by picking a vertex from the unknown set that has the smallest cost. Here cost = cost of the edge that connects that vertex to the known set. *Pick the vertex with the smallest cost that connects “known” to “unknown.”*

**A node-based greedy algorithm**

**Builds MST by greedily adding nodes**



# *Prim's Algorithm vs. Dijkstra's*

Recall:

**Dijkstra** picked the unknown vertex with smallest cost where  
cost = *distance to the source*.

**Prim's** pick the unknown vertex with smallest cost where  
cost = *distance from this vertex to the known set* (in other words,  
the cost of the smallest edge connecting this vertex to the known  
set)

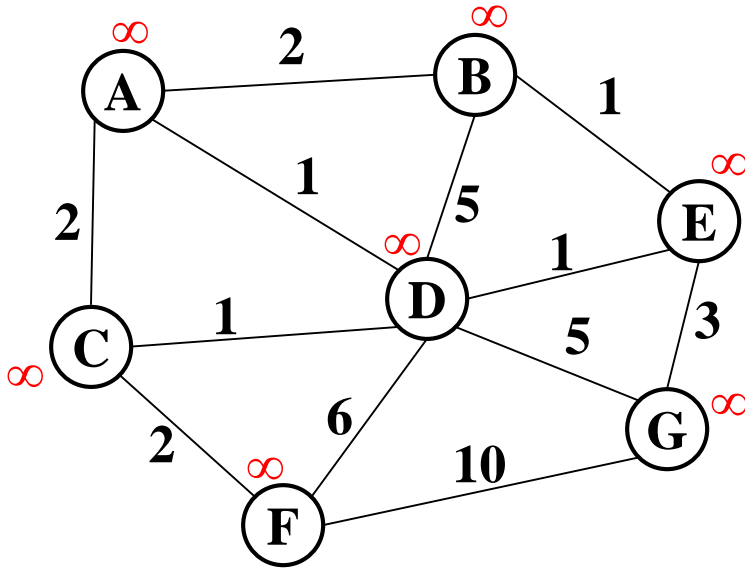
- Otherwise identical
- Compare to slides in Dijkstra lecture!

# *Prim's Algorithm for MST*

1. For each node  $v$ , set  $v.cost = \infty$  and  $v.known = false$
2. Choose any node  $v$ . (this is like your “start” vertex in Dijkstra)
  - a) Mark  $v$  as known
  - b) For each edge  $(v, u)$  with weight  $w$ :  
set  $u.cost = w$  and  $u.prev = v$
3. While there are unknown nodes in the graph
  - a) Select the unknown node  $v$  with lowest **cost**
  - b) Mark  $v$  as known and add  $(v, v.prev)$  to output (the MST)
  - c) For each edge  $(v, u)$  with weight  $w$ ,

```
        if(w < u.cost) {
            u.cost = w;
            u.prev = v;
        }
```

# Example: Find MST using Prim's

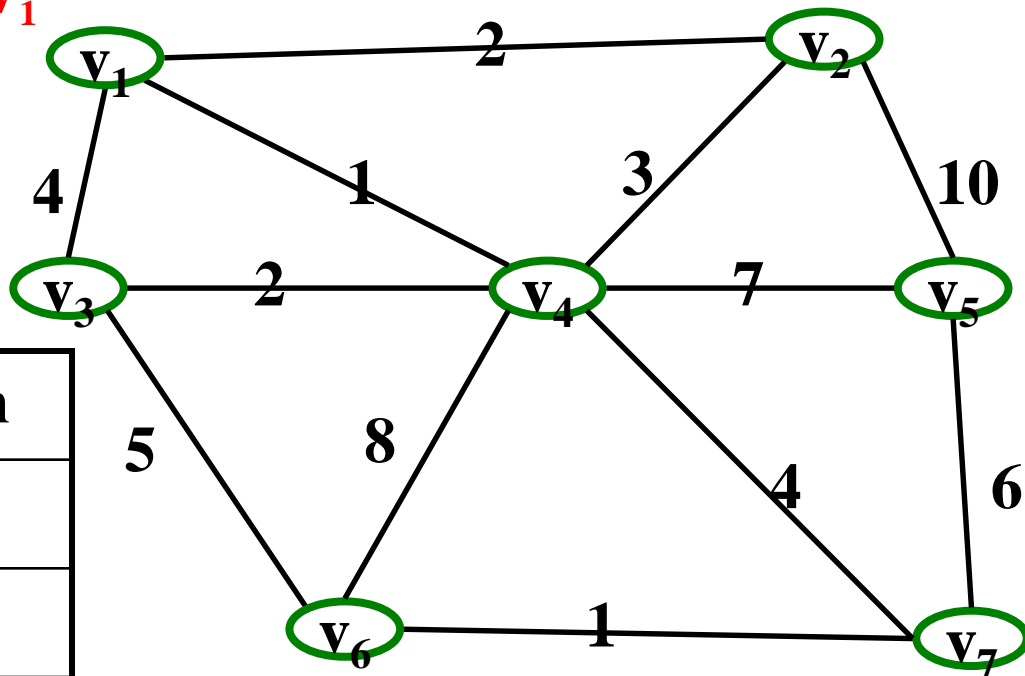


Order added to known set:

vertex	known?	cost	prev
A			
B			
C			
D			
E			
F			
G			

Start with  $V_1$

Find MST using Prim's



V	Kwn	Distance	path
v1			
v2			
v3			
v4			
v5			
v6			
v7			

Order Declared Known:

$V_1$

Total Cost:

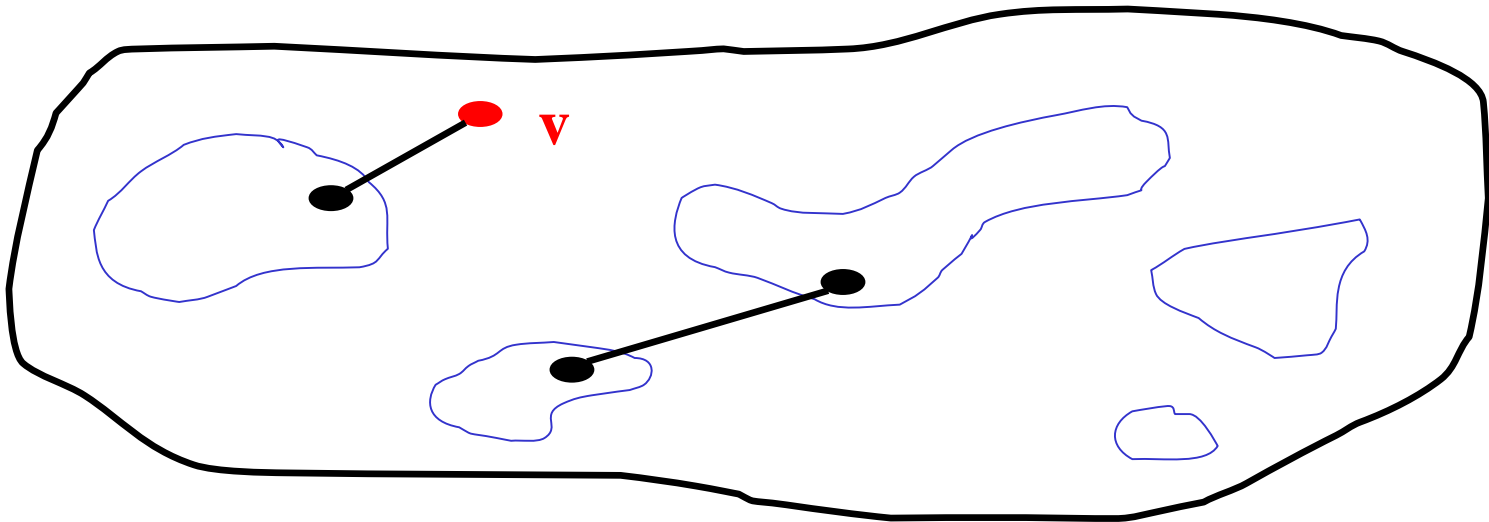
# *Prim's Analysis*

- Correctness ??
  - A bit tricky
  - Intuitively similar to Dijkstra
  - Might return to this time permitting (unlikely)
- Run-time
  - Same as Dijkstra
  - $O(|E| \log |V|)$  using a priority queue

# Kruskal's MST Algorithm

**Idea:** Grow a **forest** out of edges that do not create a cycle. Pick an edge with the smallest weight.

$G=(V,E)$



# *Kruskal's Algorithm for MST*

## **An edge-based greedy algorithm**

**Builds MST by greedily adding edges**

1. Initialize with
  - empty MST
  - all vertices marked unconnected
  - all edges unmarked
2. While all vertices are not connected
  - a. Pick the lowest cost edge  $(u, v)$  and mark it
  - b. If  $u$  and  $v$  are not already connected, add  $(u, v)$  to the MST and mark  $u$  and  $v$  as connected to each other

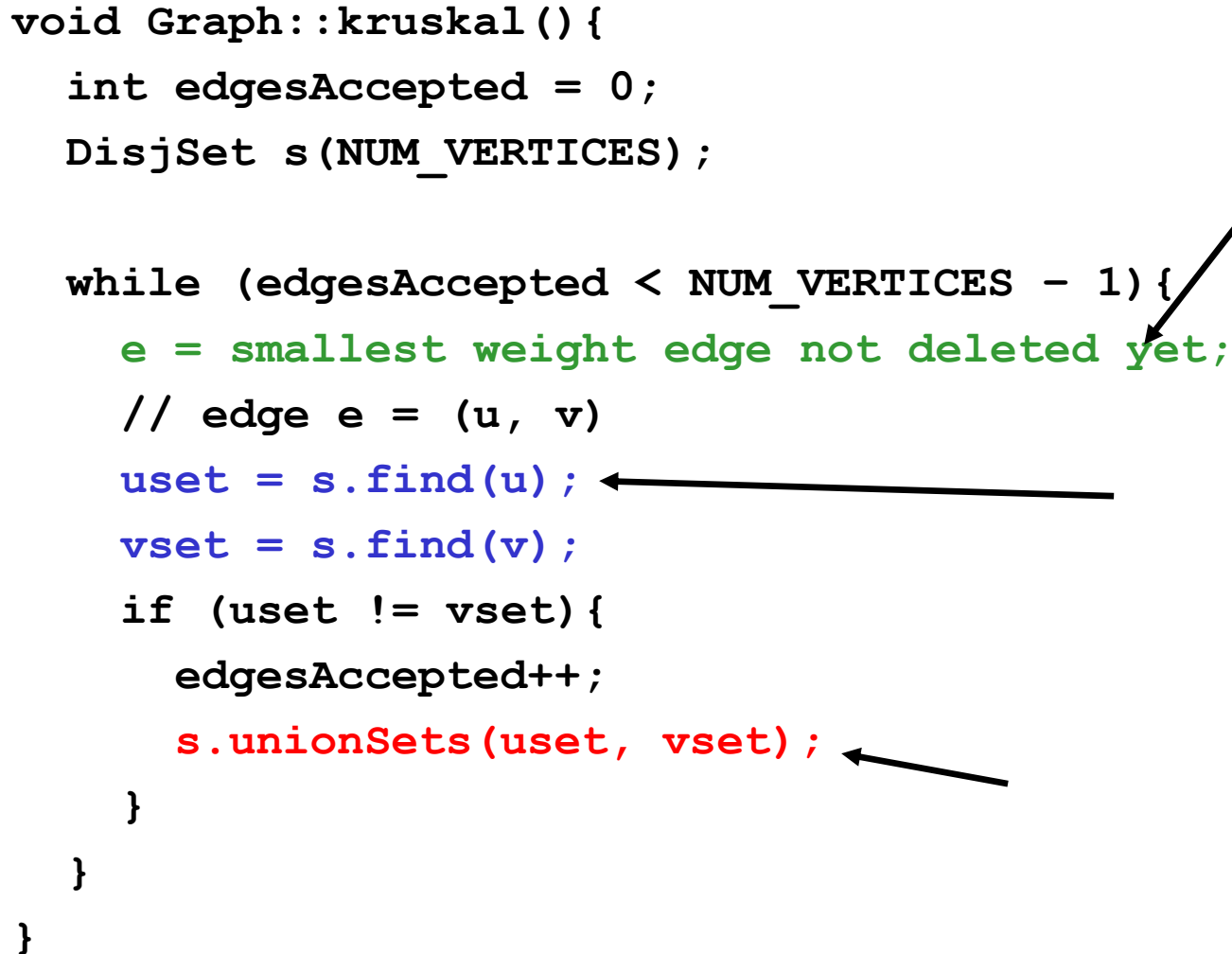
## *Aside: Union-Find aka Disjoint Set ADT*

- **Union(x,y)** – take the union of two sets named x and y
  - Given sets: {3,5,7} , {4,2,8}, {9}, {1,6}
  - **Union(5,1)**  
Result: {3,5,7,1,6}, {4,2,8}, {9},
  - To perform the union operation, we replace sets x and y by  $(x \cup y)$
- **Find(x)** – return the name of the set containing x.
  - Given sets: {3,5,7,1,6}, {4,2,8}, {9},
  - **Find(1)** returns 5
  - **Find(4)** returns 8
- We can do Union in constant time.
- We can get Find to be ***amortized*** constant time (worst case  $O(\log n)$  for an individual Find operation).

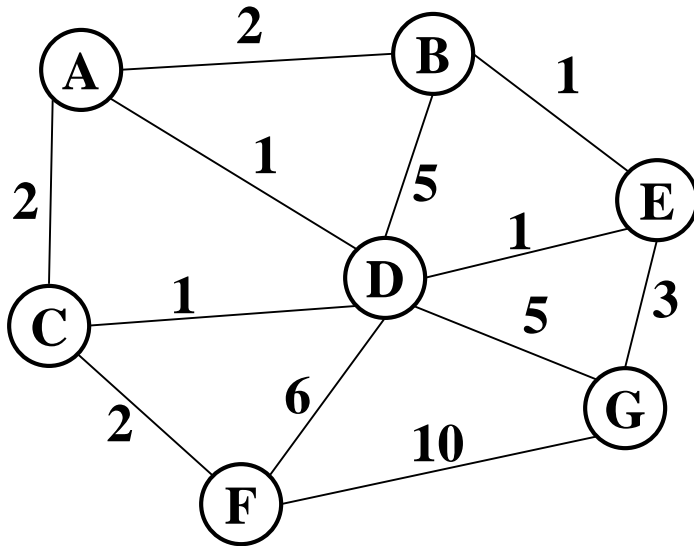
# Kruskal's pseudo code

```
void Graph::kruskal() {
    int edgesAccepted = 0;
    DisjSet s(NUM_VERTICES);

    while (edgesAccepted < NUM_VERTICES - 1) {
        e = smallest weight edge not deleted yet;
        // edge e = (u, v)
        uset = s.find(u);
        vset = s.find(v);
        if (uset != vset) {
            edgesAccepted++;
            s.unionSets(uset, vset);
        }
    }
}
```

The diagram shows three arrows pointing to specific lines of code in the Kruskal's algorithm pseudo-code. The first arrow points to the line 'e = smallest weight edge not deleted yet;', which is highlighted in green. The second arrow points to the line 'uset = s.find(u);', which is highlighted in blue. The third arrow points to the line 's.unionSets(uset, vset);', which is highlighted in red.

# Example: Find MST using Kruskal's



Edges in sorted order:

1: (A,D), (C,D), (B,E), (D,E)

2: (A,B), (C,F), (A,C)

3: (E,G)

5: (D,G), (B,D)

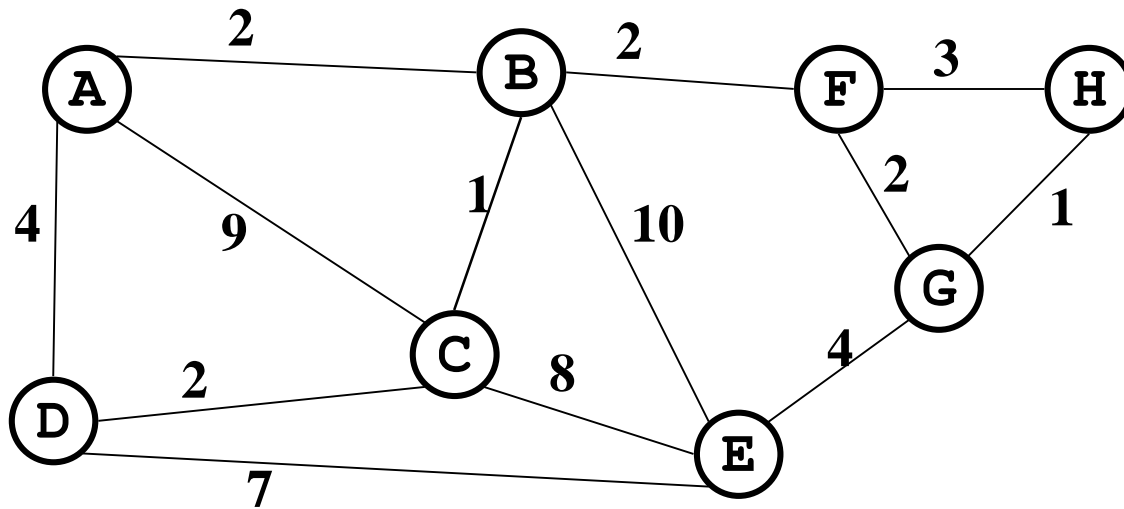
6: (D,F)

10: (F,G)

Output:

Note: At each step, the union/find sets are the trees in the forest

*Find MST using Kruskal's*



**Total Cost:**

- **Now find the MST using Prim's method.**
- **Under what conditions will these methods give the same result?**

# Correctness

Kruskal's algorithm is clever, simple, and efficient

- But does it generate a minimum spanning tree?
- How can we prove it?

First: it generates a spanning tree

- Intuition: Graph started connected and we added every edge that did not create a cycle
- Proof by contradiction: Suppose  $u$  and  $v$  are disconnected in Kruskal's result. Then there's a path from  $u$  to  $v$  in the initial graph with an edge we could add without creating a cycle. But Kruskal would have added that edge. Contradiction.

Second: There is no spanning tree with lower total cost...

# *The inductive proof set-up*

Let  $\mathbf{F}$  (stands for “forest”) be the set of edges Kruskal has added at some point during its execution.

Claim:  $\mathbf{F}$  is a subset of *one or more* MSTs for the graph  
(Therefore, once  $|\mathbf{F}|=|\mathbf{V}|-1$ , we have an MST.)

Proof: By induction on  $|\mathbf{F}|$

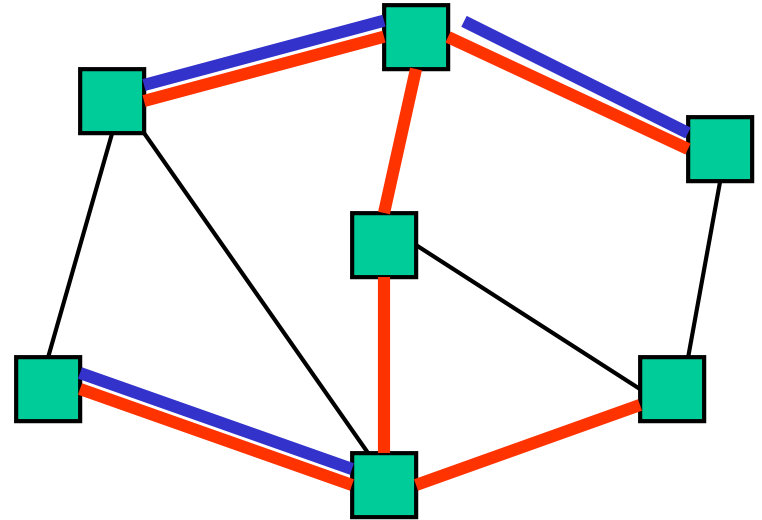
Base case:  $|\mathbf{F}|=0$ : The empty set is a subset of all MSTs

Inductive case:  $|\mathbf{F}|=k+1$ : By induction, before adding the  $(k+1)^{\text{th}}$  edge (call it  $\mathbf{e}$ ), there was some MST  $\mathbf{T}$  such that  $\mathbf{F}-\{\mathbf{e}\} \subseteq \mathbf{T} \dots$

# Staying a subset of **some** MST

Claim: **F** is a subset of *one or more* MSTs for the graph

So far: **F**-{**e**}  $\subseteq$  **T**:



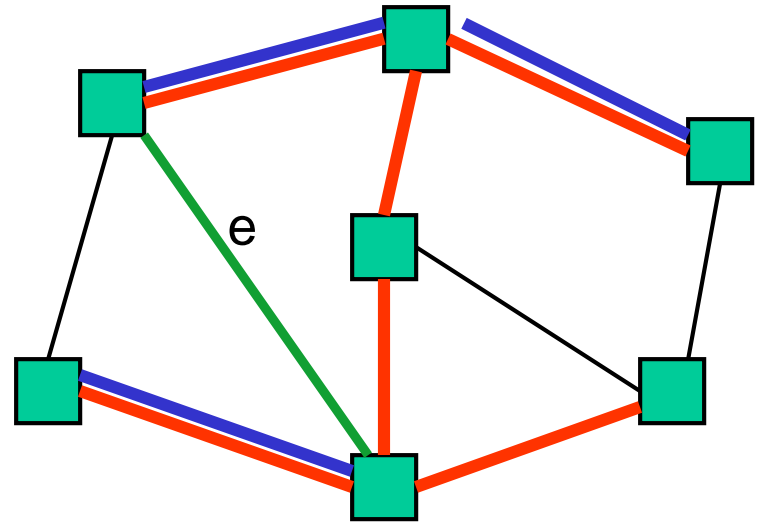
Two disjoint cases:

- If **{e}**  $\subseteq$  **T**: Then **F**  $\subseteq$  **T** and we're done
- Else **e** forms a cycle with some simple path (call it **p**) in **T**
  - Must be since **T** is a spanning tree

# Staying a subset of **some** MST

Claim: **F** is a subset of *one or more* MSTs for the graph

So far:  $\mathbf{F} - \{\mathbf{e}\} \subseteq \mathbf{T}$  and  
 $\mathbf{e}$  forms a cycle with  $\mathbf{p} \subseteq \mathbf{T}$



- There must be an edge **e2** on **p** such that **e2** is not in **F**
  - Else Kruskal would not have added **e**
- Claim: **e2.weight == e.weight**

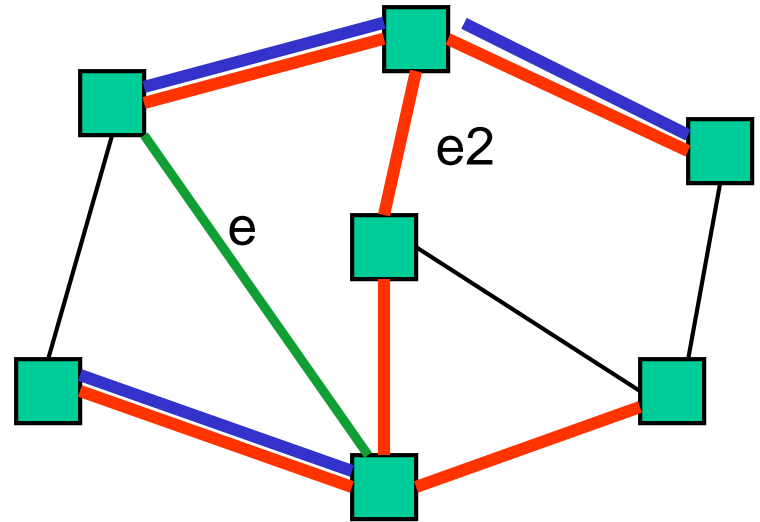
# Staying a subset of **some** MST

Claim: **F** is a subset of *one or more* MSTs for the graph

So far: **F** - {**e**}  $\subseteq$  **T**

**e** forms a cycle with **p**  $\subseteq$  **T**

**e2** on **p** is not in **F**



- Claim: **e2.weight** == **e.weight**
  - If **e2.weight** > **e.weight**, then **T** is not an MST because **T** - {**e2**} + {**e**} is a spanning tree with lower cost: contradiction
  - If **e2.weight** < **e.weight**, then Kruskal would have already considered **e2**. It would have added it since **T** has no cycles and **F** - {**e**}  $\subseteq$  **T**. But **e2** is not in **F**: contradiction

# Staying a subset of **some** MST

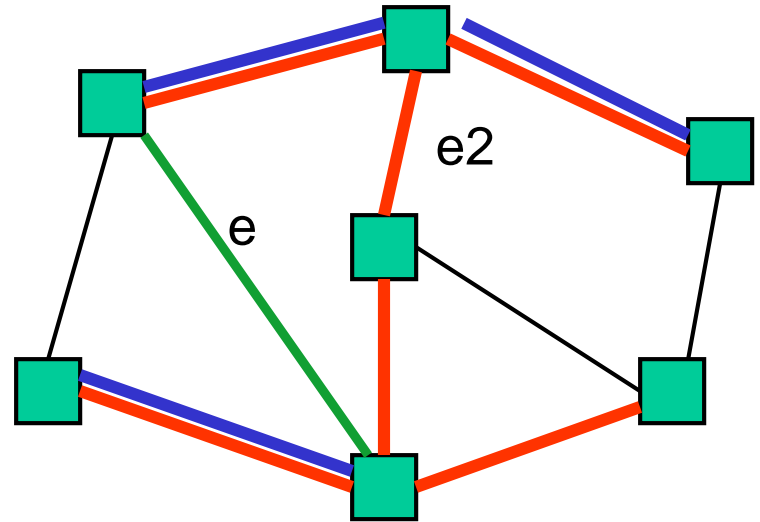
Claim: **F** is a subset of *one or more* MSTs for the graph

So far: **F** - {**e**}  $\subseteq$  **T**

**e** forms a cycle with **p**  $\subseteq$  **T**

**e2** on **p** is not in **F**

**e2.weight** == **e.weight**



- Claim: **T** - {**e2**} + {**e**} is an MST
  - It's a spanning tree because **p** - {**e2**} + {**e**} connects the same nodes as **p**
  - It's minimal because its cost equals cost of **T**, an MST
- Since **F**  $\subseteq$  **T** - {**e2**} + {**e**}, **F** is a subset of one or more MSTs

Done.