CSE 332: Data Structures and Parallelism

Section 3: Recurrences and Closed Forms

0. Not to Tree

Consider the function $f(n)$. Find a recurrence modeling the worst-case runtime of this function and then find a Big-Oh bound for this recurrence.

```
1 f(n) {
2     if (n <= 0) {
3         return 1;
4     }
5     return 2 * f(n - 1) + 1;
6 }
```

a) Find a recurrence $T(n)$ modeling the *worst-case runtime complexity* of $f(n)$

$T(n) = c_0$, if $n \leq 0$

$T(n) = T(n-1) + c_1$, otherwise

$$T(n) = \begin{cases} c_0 & \texttt{if } n <= 0 \\ T(n-1) + c_1 & \texttt{otherwise} \end{cases}$$

b) Find a closed form for $T(n)$

Unrolling the recurrence, we get $T(n) = c_1 + c_1 + \cdots + c_1 + c_0 = c_1 n + c_0$

1. To Tree

Consider the function $h(n)$. Find a recurrence modeling the worst-case runtime of this function and then find a Big-Oh bound for this recurrence.

```
1 h(n) {
2     if (n <= 1) {
3         return 1
4     } else {
5         return h(n/2) + n + 2*h(n/2)
6     }
7 }
```

a) Find a recurrence $T(n)$ modeling the *worst-case runtime complexity* of $h(n)$

$T(n) = c_0$, if $n \leq 1$

$T(n) = 2T\left(\frac{n}{2}\right) + c_1$, otherwise

$$T(n) = \begin{cases} c_0 & \text{if } n \leq 1 \\ 2T\left(\frac{n}{2}\right) + c_1 & \text{otherwise} \end{cases}$$

b) Find a closed form for $T(n)$

The recursion tree has height $\lg(n)$, each non-leaf level i has work $c_1 2^i$, and the leaf level has work $c_0 2^{\lg(n)}$. Putting this together, we have:

$$\left(\sum_{i=0}^{\lg(n)-1} c_1 2^i\right) + c_0 2^{\lg(n)} = c_1\left(\sum_{i=0}^{\lg(n)-1} 2^i\right) + c_0 n$$

$$= c_1 \frac{1-2^{\lg(n)}}{1-2} + c_0 n$$

$$= c_1\left(2^{\lg(n)} - 1\right) + c_0 n$$

$$= c_1(n-1) + c_0 n$$

$$= (c_0 + c_1)n - c_1$$

2. To Tree or Not to Tree

Consider the function $f(n)$. Find a recurrence modeling the worst-case runtime of this function and then find a Big-Oh bound for this recurrence.

```
1  f(n) {
2      if (n <= 1) {
3          return 0
4      }
5      int result = f(n/2)
6      for (int i = 0; i < n; i++) {
7          result *= 4
8      }
9      return result + f(n/2)
10 }
```

a) Find a recurrence $T(n)$ modeling the *worst-case runtime complexity* of $f(n)$

We look at the three separate components (base case, non-recursive work, recursive work). The base case is a constant amount of work, because we only do a return statement. We'll label it c_0. The non-recursive work is a constant amount of work (we'll call it c_1) for the assignments and `if` tests and a constant (we'll call c_2) multiple of n for the loops. The recursive work is $2T\left(\frac{n}{2}\right)$.

Putting these together, we get:

$T(n) = c_0$, if 1

$T(n) = 2T\left(\frac{n}{2}\right) + c_2 n + c_1$, otherwise

$$T(n) = \begin{cases} c_0 & n = 1 \\ 2T\left(\frac{n}{2}\right) + c_2 n + c_1 & otherwise \end{cases}$$

b) Find a closed form for $T(n)$

The recursion tree has $\lg(n)$ height, each non-leaf node of the tree does $c_2 \frac{n}{2^i} + c_1$ work, each leaf node does c_0 work, and each level has 2^i nodes.

So, the total work is

$$\left(\sum_{i=0}^{\lg(n)-1} 2^i\left(c_1 + c_2 \frac{n}{2^i}\right)\right) + c_0 \cdot 2^{\lg(n)}$$

$$= \left(\sum_{i=0}^{\lg(n)-1} 2^i c_1 + c_2 n\right) + c_0 \cdot (n)$$

$$= c_1 \frac{1-2^{\lg(n)}}{1-2} + c_2 n \lg(n) + c_0 n$$

$$= c_1(n-1) + c_2 n \lg(n) + c_0 n$$

3. Big-Oof Bounds

Consider the function $f(n)$. Find a recurrence modeling the worst-case runtime of this function and then find a Big-Oh bound for this recurrence.

```
1  f(n) {
2      if (n == 1) {
3          return 0
4      }
5
6      int result = 0
7      for (int i = 0; i < n; i++) {
8          for (int j = 0; j < i; j++) {
9              result += j
10
11         }
12     }
13     return f(n/2) + result + f(n/2)
14 }
```

a) Find a recurrence $T(n)$ modeling the *worst-case runtime complexity* of $f(n)$

$T(n) = c_0$, if $n = 1$

$T(n) = 2T\left(\frac{n}{2}\right) + c_2\frac{n(n-2)}{2} + c_1$, otherwise

$$T(n) = \begin{cases} c_0 & \text{if } n = 1 \\ 2T\left(\frac{n}{2}\right) + c_2\frac{n(n-1)}{2} + c_1 & \text{otherwise} \end{cases}$$

b) Find a Big-Oh bound for your recurrence.

Since we only want a Big-Oh, we can actually leave off lower-order terms when doing our analysis, as they won't affect the runtime bounds; so, we can ignore the constants c_1 and c_2 in our analysis.

Note that $\frac{n(n-1)}{2} = \frac{n^2}{2} - \frac{n}{2} \in \mathcal{O}(n^2)$. We can, again, ignore the lower-order term ($\frac{n}{2}$) since we only want a Big-Oh bound.

The recursion tree has $\lg(n)$ height, each non-leaf node of the tree does $\left(\frac{n}{2^i}\right)^2$ work, each leaf node does c_0 work, and each level has 2^i nodes.

So, the total work is:

$$\sum_{i=0}^{\lg(n)-1} 2^i\left(\frac{n}{2^i}\right)^2 + c_0 \cdot 2^{\lg n} = n^2 \sum_{i=0}^{\lg(n)-1} \frac{2^i}{4^i} + c_0 n < n^2 \sum_{i=0}^{\infty} \frac{1}{2^i} + c_0 n = \frac{n^2}{1-\frac{1}{2}} + c_0 n$$

This expression is upper-bounded by n^2 so $T \in \mathcal{O}(n^2)$.

4. Odds Not in Your Favor

Consider the function $g(n)$. Find a recurrence modeling the worst-case runtime of this function and then find a Big-Oh bound for this recurrence.

```
1  g(n) {
2      if (n <= 1) {
3         return 1000
4      }
5      if (g(n/3) > 5) {
6          for (int i = 0; i < n; i++) {
7              println("Yay!")
8          }
9          return 5 * g(n/3)
10     } else {
11         for (int i = 0; i < n * n; i++) {
12             println("Yay!")
13         }
14         return 4 * g(n/3)
15     }
16 }
```

a) Find a recurrence $T(n)$ modeling the *worst-case runtime complexity* of $f(n)$

$T(n) = c_0$, if $n \leq 1$

$T(n) = 2T\left(\frac{n}{3}\right) + c_1 n + c_2$, otherwise

$$T(n) = \begin{cases} c_0 & \text{if } n \leq 1 \\ 2T\left(\frac{n}{3}\right) + c_1 n + c_2 & \text{otherwise} \end{cases}$$

b) Find a closed form for $T(n)$

The recursion tree has height $\log_3(n)$, each non-leaf level i has work $\left(\frac{c_1 n}{3^i} + c_2\right)2^i$, and the leaf level has work $c_0 2^{\log_3(n)}$. Putting this together, we have:

$$\sum_{i=0}^{\log_3(n)-1} \left(\left(\frac{c_1 n}{3^i} + c_2\right)2^i\right) + c_0 2^{\log_3(n)}$$

$$= \sum_{i=0}^{\log_3(n)-1} \left(\frac{c_1 n 2^i}{3^i} + c_2 2^i\right) + c_0 2^{\log_3(n)}$$

$$= c_1 n\left(\sum_{i=0}^{\log_3(n)-1} \left(\frac{2}{3}\right)^i\right) + c_2\left(\sum_{i=0}^{\log_3(n)-1} 2^i\right) + c_0 2^{\log_3(n)}$$

Using the finite geometric series,

$$= c_1 n\left(\frac{1-\left(\frac{2}{3}\right)^{\log_3(n)}}{1-\frac{2}{3}}\right) + c_2\left(\frac{1-2^{\log_3(n)}}{1-2}\right) + c_0 2^{\log_3(n)}$$

$$= 3c_1 n\left(1 - \left(\frac{2}{3}\right)^{\log_3(n)}\right) + c_2\left(2^{\log_3(n)} - 1\right) + c_0 2^{\log_3(n)}$$

$$= 3c_1 n\left(1 - \frac{n^{\log_3(2)}}{n}\right) + c_2\left(n^{\log_3(2)} - 1\right) + c_0 n^{\log_3(2)}$$

$$= 3c_1 n - 3c_1 n^{\log_3(2)} + c_2 n^{\log_3(2)} - c_2 + c_0 n^{\log_3(2)}$$

$$= 3c_1 n + \left(c_0 + c_2 - 3c_1\right)n^{\log_3(2)} - c_2$$