## The Algorithm

1. For each node $\mathbf{v}$, set $\mathbf{v}$.cost $=\infty$ and $\mathbf{v}$. known $=$ false
2. Set source. cost $=0$
3. While there are unknown nodes in the graph
a) Select the unknown node $\mathbf{v}$ with lowest cost
b) Mark vas known
c) For each edge ( $\mathbf{v}, \mathbf{u}$ ) with weight $\mathbf{w}$, if $\mathbf{u}$ is unknown,
$\mathrm{c} 1=\mathrm{v}$. cost $+\mathrm{w} / /$ cost of best path through $v$ to $u$
if ( $\mathrm{c} 1<\mathrm{c} 2$ ) $\{/ /$ if the path through v is better
u.cost $=\mathrm{c} 1$
u .path $=\mathrm{v} / /$ for computing actual paths
\}

Example \#1


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## Example \#2



Order Added to Known Set:

$\square$

## Efficiency, first approach

Use pseudocode to determine asymptotic run-time - Notice each edge is processed only once
dijkstra (Graph G, Node start)
for each node: $x$.cost=infinity, $x$.known=false tart.cost = 0
while(not all nodes are known)
b $=$ find unknown node with smallest cost
b. known = true
for each edge ( $b, a$ ) in $G$
if(!a.known)
if (b.cost + weight((b,a)) < a.cost)
a.cost $=b \cdot \operatorname{cost}+$ weight ( $(b, a))$
a. path $=b$
\}
\}
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## Efficiency, second approach

Use pseudocode to determine asymptotic run-time
dijkstra(Graph G, Node start) {
dijkstra(Graph G, Node start) {
for each node: x.cost=infinity, x.known=false
for each node: x.cost=infinity, x.known=false
start.cost = 0
start.cost = 0
build-heap with all nodes
build-heap with all nodes
while(heap is not empty)
while(heap is not empty)
known = true
known = true
for each edge
for each edge
for each edge (b,a) in G
for each edge (b,a) in G
if(b.cost +
if(b.cost +
(b.cost + weight((b,a)) < a.cost) {
(b.cost + weight((b,a)) < a.cost) {
a.path = b
a.path = b
}
}
,
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