

InsertionSort

- ❖ Idea: At step k , insert the k^{th} element in the correct position
 - Sort first two elements
 - Now insert 3rd element in order
 - ...
- ❖ Loop invariant (“when loop index is i ”):
 - First i elements are in sorted order
- ❖ Time:
Best-case: _____ Worst-case: _____ “Average” case: _____
- ❖ Characteristics:
Stable: _____ In-place: _____

Demo:

https://docs.google.com/presentation/d/10b9aRqpGJu8pUk8OpfqUIEEm8ou-zmmC7b_BE5wgNg0/present

SelectionSort

- ❖ Idea: At step k , select the smallest elt and put it at k^{th} position
 - Find smallest element, put it 1st
 - Find next smallest element, put it 2nd
 - ...
- ❖ Loop invariant (“when loop index is i ”):
 - First i elements are the i smallest elements in sorted order
- ❖ Time:
Best-case: _____ Worst-case: _____ “Average” case: _____
- ❖ Characteristics:
Stable: _____ In-place: _____

Demo:

<https://docs.google.com/presentation/d/1p6g3r9BpwTARjUylA0V0ysspP2temzHNJEJjCG41I4r0/edit>

In-place HeapSort

- ❖ Idea: Put everything in a **MAX** heap ; successively delete Max
 - insert each `arr[i]` –OR– better yet, use `buildHeap`
 - `for(i=0; i < arr.length; i++)`
`arr[arr.length - i] = deleteMax();`
- ❖ Loop invariant (“when loop index is **i**”): same as naïve version
- ❖ Time:
Best-case: _____ Worst-case: _____ “Average” case: _____
- ❖ Characteristics:
Stable: _____ In-place: _____

Demo:

<https://docs.google.com/presentation/d/1SzcQC48OB9agStD0dFRgccU-tyjD6m3esrSC-GLxmNc/present>

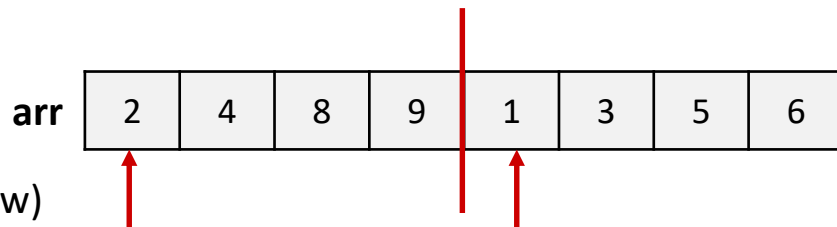
MergeSort: Merging Example (1 of 10)

❖ Start with:



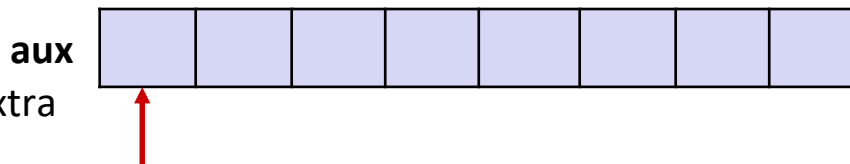
❖ Return from left and right recursion

- (pretend it works for now)



❖ Merge

- Use 3 cursors and an extra auxiliary array
- When done, copy the extra array back to the original



Lecture Outline

- ❖ Comparison-based Sorting
 - Review
 - Simple algorithms
 - InsertionSort
 - SelectionSort
 - Fancier Algorithms: HeapSort
 - Fancier algorithms using Divide-and-Conquer
 - Intro
 - MergeSort
 - **QuickSort**

Step #3: Recursion (2 of 3)

