Name: $\qquad$
Email address: $\qquad$
Quiz Section: $\qquad$

## CSE 332 Autumn 2013: Midterm Exam

(closed book, closed notes, no calculators)

Instructions: Read the directions for each question carefully before answering. We will give partial credit based on the work you write down, so show your work! Use only the data structures and algorithms we have discussed in class or that were mentioned in the book so far.

Note: For questions where you are drawing pictures, please circle your final answer for any credit.

Good Luck!

Total: 100 points. Time: 50 minutes.

| Question | Max Points | Score |
| :---: | :---: | :---: |
| 1 | 18 |  |
| 2 | 6 |  |
| 3 | 10 |  |
| 4 | 6 |  |
| 5 | 8 |  |
| 6 | 6 |  |
| 7 | 8 |  |
| 8 | 15 |  |
| 9 | 15 |  |
| 10 | 8 |  |
| Total | 100 |  |

## 1. (18 pts) Big-Oh

( 2 pts each) For each of the functions $f(N)$ given below, indicate the tightest bound possible (in other words, giving $\mathrm{O}\left(2^{\mathrm{N}}\right)$ as the answer to every question is not likely to result in many points). Unless otherwise specified, all logs are base 2. Your answer should be as "tight" and "simple" as possible. For questions that ask about running time of operations, assume that the most efficient implementation is used.

You do not need to explain your answer.
a) Find in a separate chaining hash table containing $N$ elements where each bucket points to a sorted linked list (worst case)
b) Preorder traversal on an AVL tree containing $N$ elements (worst case)
c) $f(N)=\mathrm{N}(\log \mathrm{N})^{2}+\mathrm{N} \log \log \mathrm{N}$
d) $f(N)=\mathrm{N} \log _{2}(\mathrm{~N}+\mathrm{N})$
e) Enqueue in a queue containing $N$ elements implemented using linked list nodes (worst case)
f) Determine what the maximum value is in a binary search tree containing $N$ elements (worst case) $\qquad$
g) DecreaseKey ( $k, v$ ) on a binary min heap containing $N$ elements. Assume you have a reference to the key $k$. $v$ is the amount that $k$ should be decreased. (worst case) $\qquad$
h) Merging two binary min heaps containing $N$ elements each. (worst case) $\qquad$
i) $T(N)=T(N-1)+5$ $\qquad$
2. ( 6 pts) Big-Oh and Run Time Analysis: Describe the worst case running time of the following pseudocode functions in Big-Oh notation in terms of the variable $n$. Your answer should be as "tight" and "simple" as possible.
Showing your work is not required

```
    I. void sunny (int n, int sum) {
    for (int k = n; k > 0; k = k/2) {
        for (int j = n; j > 0; j--)
                sum++;
    }
    }
II. int funny (int n) {
    if (n < 100)
        return n - 2;
    else {
            return funny (n - 2);
    }
}
III. int happy (int n, int sum) {
    for (int k = 0; k < n; k = k++) {
        for (int i = 0; i < k; i++)
                sum++;
        for (int j = n * n; j > 0; j--)
                sum++;
    }
    return sum;
}
```

3. ( 10 pts) Big-O, Big $\Omega, \operatorname{Big} \Theta$
(2 pts each) For parts (a) - (e) circle whether the statement is true or false. You do not need to show any work or give an explanation.

TRUE / FALSE
a) $\mathrm{N} \log \mathrm{N}+100 \mathrm{~N}=\Omega\left(\mathrm{N}^{2}\right)$

TRUE / FALSE
b) $50 \mathrm{~N}^{2}+\mathrm{N}^{3}+10 \mathrm{~N}^{4}=\Theta\left(\mathrm{N}^{3}\right)$

TRUE / FALSE
c) $5 \mathrm{~N}+(1 / 2) \mathrm{N}=\Omega(\mathrm{N} \log \mathrm{N})$

TRUE / FALSE
d) $\mathrm{N}^{3}+\mathrm{N}^{3}=\mathrm{O}\left(\mathrm{N}^{6}\right)$

TRUE / FALSE
e) $5 \mathrm{~N} \log \mathrm{~N}=\Omega\left(\mathrm{N}^{2}\right)$

## 4. (6 pts) Recurrence Relationships -

Suppose that the running time of an algorithm satisfies the recurrence relationship

$$
\mathrm{T}(1)=7 .
$$

and

$$
\mathrm{T}(\mathrm{~N})=\mathrm{T}(\mathrm{~N}-2)+4 \quad \text { for odd integers } \mathrm{N}>1
$$

Find the closed form for $\mathrm{T}(\mathrm{N})$ and show your work step by step. In other words express $\mathrm{T}(\mathrm{N})$ as a function of N . Your answer should not be in Big-Oh notation show the relevant exact constants in your answer (e.g. don't use "C" in your answer).
5. (8 pts) Draw the AVL tree that results from inserting the keys $5,1,4,7,8,9,2$ in that order into an initially empty AVL tree. You are only required to show the final tree, although if you draw intermediate trees, please circle your final result for ANY credit.
a) ( 2 pts) What is the minimum number of nodes in an AVL tree of height 5 ? (Hint: the height of a tree consisting of a single node is 0 ) Give an exact number not a formula.
b) ( 2 pts ) What is the minimum number of nodes than must be examined in order to find the minimum value in an AVL tree of height 5? Give an exact number not a formula.
c) ( 2 pts) Given a perfect binary tree of height h , what are the minimum and maximum number of nodes in the right subtree of the root?

Minimum:

Maximum:

## 7. (8 pts) B-trees

a) (4 pts) Given a B-tree of height H , give the tight big-O running time for a find operation in terms of M, L, N, and H. Do not simplify your answer - it is o.k. to ignore constants and smaller order terms other than those containing M, L, N, or H - but do not remove any occurrences of $\mathrm{M}, \mathrm{L}, \mathrm{N}$, or $\mathbf{H}$ from your answers.

Worst case running time for a find operation $=$

Best case running time for a find operation $=$
b) (4 pts) Given the following parameters for a B-tree with $\mathrm{M}=11$ and $\mathrm{L}=8$ :

Key Size $=10$ bytes
Pointer Size $=2$ bytes
Data Size $=16$ bytes per record (includes the key)
Assuming that M and L were chosen appropriately, what is the likely size of a disk block on the machine where this implementation will be deployed? Give a numeric answer and a short justification based on two equations using the parameter values above.

## 8. ( 15 pts) Hash Tables

For a) and b) below, insert the following elements in this order: 51, 30, 39, 10, 19, 40.
For each table, TableSize $=10$, and you should use the primary hash function $h(k)=$ $\mathrm{k} \% 10$. If an item cannot be inserted into the table, please indicate this and continue inserting the remaining values.
a) Quadratic probing hash table

| 0 |  |
| :--- | :--- |
| 1 |  |
| 2 |  |
| 3 |  |
| 4 |  |
| 5 |  |
| 6 |  |
| 7 |  |
| 8 |  |
| 9 |  |

b) Separate chaining hash table - use an unsorted linked list for each bucket, insert at front

| 0 |  |
| :--- | :--- |
| 1 |  |
| 2 |  |
| 3 |  |
| 4 |  |
| 5 |  |
| 6 |  |
| 7 |  |
| 8 |  |
| 9 |  |

c) What is the load factor in Table a)?
d) In a sentence or two, describe one advantage that separate chaining has over quadratic probing?
e) In a sentence or two, describe one advantage that linear probing has over quadratic probing?
9. (15 pts) Min Three Heaps - As discussed on homework 2, a three heap with $n$ elements can be stored in an array A, where $\mathrm{A}[0]$ contains the root of the tree.
a) Given indexing that starts at $\mathrm{A}[0]$, give the equations for how you would calculate the three children of the value stored at index i as well as the parent.

## Child 1: <br> Child 2: <br> Child 3:

Parent:
b) Implement the insertThreeHeap method. To avoid ambiguity, please give your answer in Java. You should implement a min heap. You can assume that the array A is big enough to contain all inserted values and that the values you will be inserting are ints >0 and will not contain duplicates. You do not need to optimize your code by "bubbling a hole" up/down, but it is fine to do so. You may NOT assume that other methods have been implemented (such as percolateUp, percolateDown, swap, etc.), however it is fine to implement such helper methods if you would like to. Assume that you can read and modify the size of the heap.
// Given that values are stored in A[0] to A[size-1],
// insert val into the min three heap.
void insertThreeHeap(int[] A, int val) \{
9. Min Three Heaps (continued) - feel free to use the space on this page as well.

## 10. (8 pts) B-tree Insertion

a) (2pts) In the B-Tree shown below, please write in the appropriate values for the interior nodes.
b) (2 pts) Based on the picture below, what are the values for M and L ? $\mathrm{M}=$

$$
\mathrm{L}=
$$

c) (4 pts) Starting with the B-tree shown below, insert 27. Draw and circle the resulting tree (including values for interior nodes) below. Use the method for insertion described in lecture and used on homework.


