Solving Recurrences

Data Structures and Parallelism

Warm Up

Write a recurrence to represent the running time of this code

```
int Mystery(int n){
   if(n <= 5)
      return 1;
   for(int i=0; i<n; i++){
      for(int j=0; j<n; j++){
         System.out.println("hi");
      }
   }
   return n*Mystery(n/2);
}
```

Outline

Last Time:

- We wrote recurrences to describe the running times of recursive functions.

Wednesday:

- How do we turn a recurrence into a big-Θ bound?

Friday:

- A big hammer for getting big-Θ bounds
- Amortized Bounds

Tree Method

Idea:

- Since we're making recursive calls, let's just draw out a tree, with one node for each recursive call.
- Each of those nodes will do some work, and (if they make more recursive calls) have children.
- If we can just add up all the work, we can find a big-Θ bound.

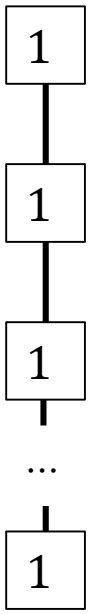
Solving Recurrences I: Binary Search

$$T(n) \equiv \begin{cases} 1 \; when \; n \leq 1 \\ T\left(\frac{n}{2}\right) + 1 \; otherwise \end{cases}$$

0. Draw the tree.
1. What is the input size at level i?
2. What is the number of nodes at level i?
3. What is the work done at recursive level i?
4. What is the last level of the tree?
5. What is the work done at the base case?
6. Sum over all levels (using 3,5).
7. Simplify

Solving Recurrences I: Binary Search

$$T(n) \equiv \begin{cases} 1 \; when \; n \leq 1 \\ T\left(\frac{n}{2}\right) + 1 \; otherwise \end{cases}$$



0. Draw the tree.
1. What is the input size at level i?
2. What is the number of nodes at level i?
3. What is the work done at recursive level i?
4. What is the last level of the tree?
5. What is the work done at the base case?
6. Sum over all levels (using 3,5).
7. Simplify

Solving Recurrences I: Binary Search

$$T(n) = \begin{cases} 1 \; when \; n \leq 1 \\ T\left(\frac{n}{2}\right) + 1 \; otherwise \end{cases}$$

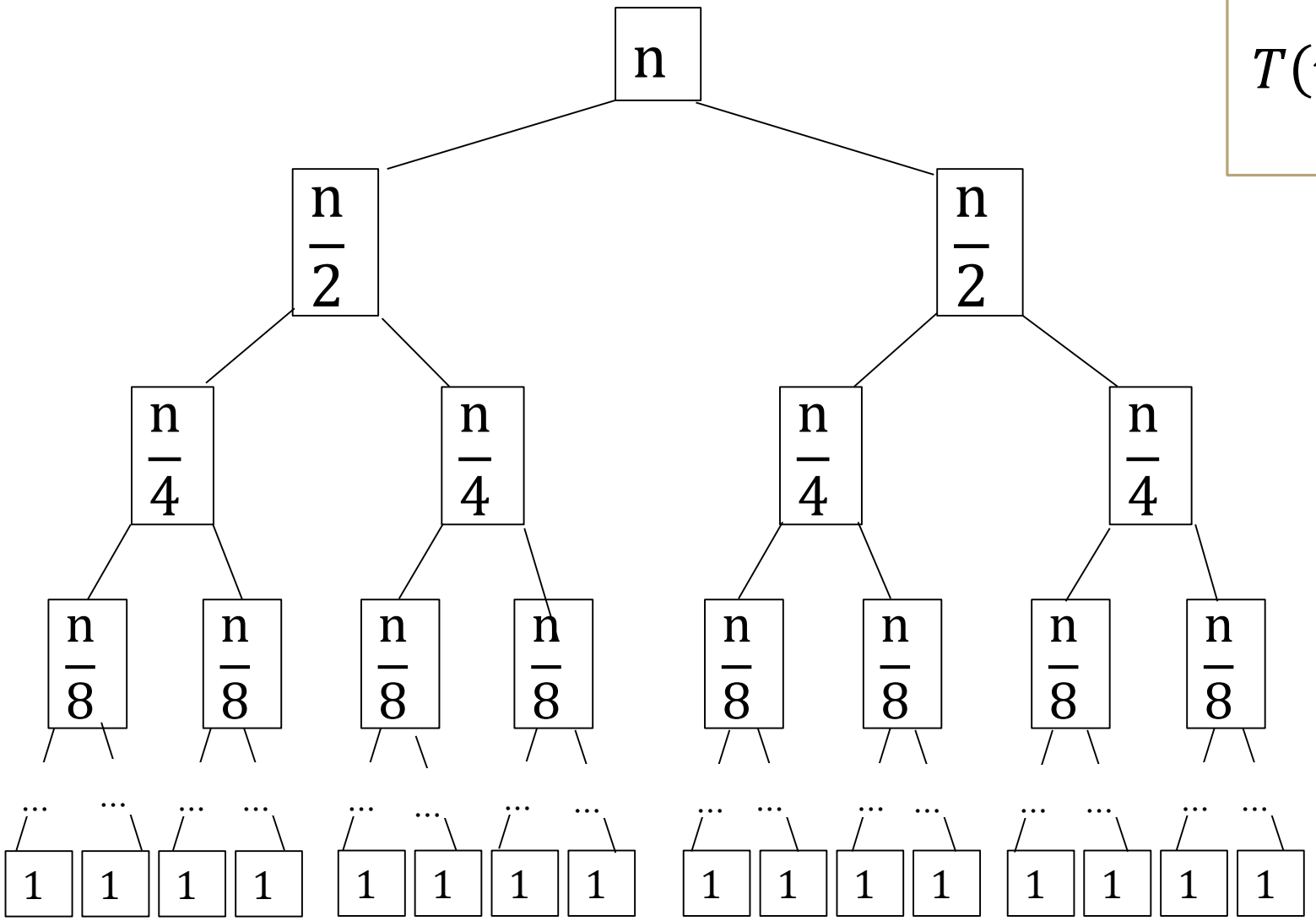
Level	Input Size	Work/call	Work/level
0	n	1	1
1	$n/2$	1	1
2	$n/2^2$	1	1
i	$n/2^i$	1	1
$\log_2 n$	1	1	1

0. Draw the tree.
1. What is the input size at level i?
2. What is the number of nodes at level i?
3. What is the work done at recursive level i?
4. What is the last level of the tree?
5. What is the work done at the base case?
6. Sum over all levels (using 3,5).
7. Simplify

$$\sum_{i=0}^{\log_2 n - 1} 1 + 1 = \log_2 n$$

Solving Recurrences II:

$$T(n) = \begin{cases} 1 \; when \; n \leq 1 \\ 2T\left(\frac{n}{2}\right) + n \; otherwise \end{cases}$$



Tree Method Formulas

$$T(n) = \begin{cases} 1 \text{ when } n \leq 1 \\ 2T\left(\frac{n}{2}\right) + n \text{ otherwise} \end{cases}$$

How much work is done by recursive levels (branch nodes)?

1. What is the input size at level i?
   - i = 0 is overall root level.

   $(n/2^i)$
2. At each level i, how many calls are there?

   2^i
3. At each level i, how much work is done??

   $2^i(n/2^i) = n$

$$Recursive\ work = \sum_{i=0}^{lastRecursiveLevel} branchNum(i)branchWork(i)$$

$$\sum_{i=0}^{\log_2 n - 1} 2^i \left(\frac{n}{2^i}\right)$$

How much work is done by the base case level (leaf nodes)?

4. What is the last level of the tree? $(n/2^i) = 1 \rightarrow 2^i = n \rightarrow i = \log_2 n$
5. What is the work done at the last level?

$$NonRecursive\ work = WorkPerBaseCase \times numberCalls$$

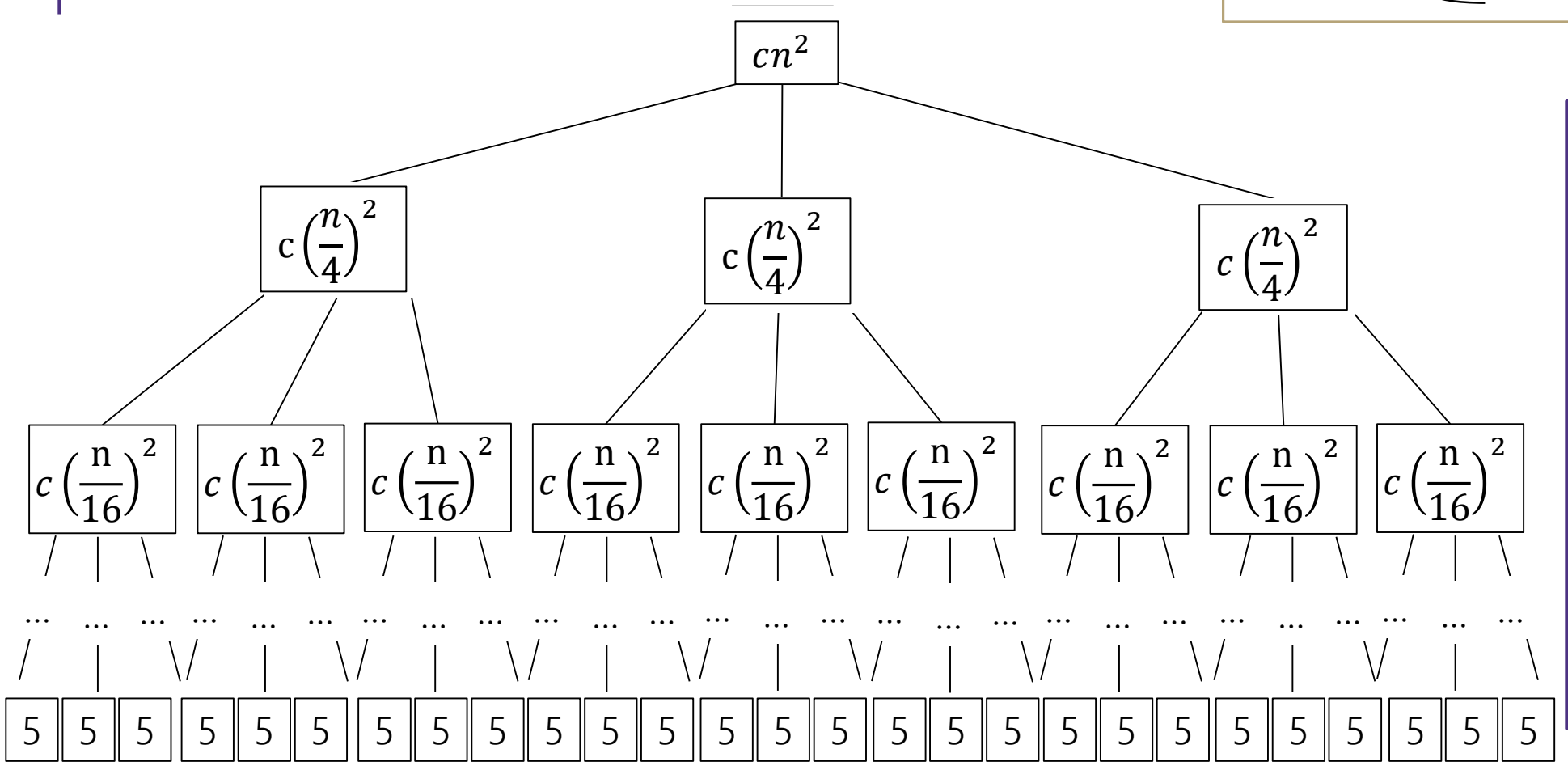
$$1 \cdot 2^{\log_2 n} = n$$

6. Combine and Simplify

$$T(n) = \sum_{i=0}^{\log_2 n - 1} 2^i \left(\frac{n}{2^i}\right) + n = n \log_2 n + n$$

Solving Recurrences III

$$T(n) = \begin{cases} 5 \; when \; n \leq 4 \\ 3T\left(\frac{n}{4}\right) + cn^2 \; otherwise \end{cases}$$



Answer the following questions:

1. What is input size on level i?
2. Number of nodes at level i?
3. Work done at recursive level i?
4. Last level of tree?
5. Work done at base case?
6. What is sum over all levels?

Solving Recurrences III

$$T(n) = \begin{cases} 5 \text{ when } n \leq 4 \\ 3T\left(\frac{n}{4}\right) + cn^2 \text{ otherwise} \end{cases}$$

1. Input size on level i? $\frac{n}{4^i}$

$$c\left(\frac{n}{4^i}\right)^2$$

2. How many calls on level i? 3^i

3. How much work on level i? $3^i c\left(\frac{n}{4^i}\right)^2 = \left(\frac{3}{16}\right)^i cn^2$

4. What is the last level? When $\frac{n}{4^i} = 4 \rightarrow \log_4 n - 1$

5. A. How much work for each leaf node? 5

B. How many base case calls? $3^{\log_4 n - 1} = \frac{3^{\log_4 n}}{3}$ power of a log $x^{\log_b y} = y^{\log_b x}$ $= \frac{n^{\log_4 3}}{3}$

Level (i)	Number of Nodes	Work per Node	Work per Level
0	1	cn^2	cn^2
1	3	$c\left(\frac{n}{4}\right)^2$	$\frac{3}{16}cn^2$
2	3^2	$c\left(\frac{n}{4^2}\right)^2$	$\left(\frac{3}{16}\right)^2 cn^2$
i	3^i	$c\left(\frac{n}{4^i}\right)^2$	$\left(\frac{3}{16}\right)^i cn^2$
Base = $\log_4 n - 1$	$3^{\log_4 n - 1}$	5	$\left(\frac{5}{3}\right) n^{\log_4 3}$

6. Combining it all together…

$$T(n) = \sum_{i=0}^{\log_4 n - 2} \left(\frac{3}{16}\right)^i cn^2 + \left(\frac{5}{3}\right) n^{\log_4 3}$$

Solving Recurrences III

7. Simplify…

$$T(n) = \sum_{i=0}^{\log_4 n - 2} \left(\frac{3}{16}\right)^i cn^2 + \left(\frac{5}{3}\right) n^{\log_4 3}$$

factoring out a constant

$$\sum_{i=a}^{b} cf(i) = c\sum_{i=a}^{b} f(i)$$

$$T(n) = cn^2 \sum_{i=0}^{\log_4 n - 2} \left(\frac{3}{16}\right)^i + \left(\frac{5}{3}\right) n^{\log_4 3}$$

finite geometric series

$$\sum_{i=0}^{n-1} x^i = \frac{x^n - 1}{x - 1}$$

Closed form:

$$T(n) = cn^2 \left(\frac{\frac{3}{16}^{\log_4 n - 1} - 1}{\frac{3}{16} - 1}\right) + \left(\frac{5}{3}\right) n^{\log_4 3}$$

If we're trying to prove upper bound…

$$T(n) \leq cn^2 \sum_{i=0}^{\infty} \left(\frac{3}{16}\right)^i + 4n^{\log_4 3}$$

infinite geometric series

$$\sum_{i=0}^{\infty} x^i = \frac{1}{1 - x}$$

when -1 < x < 1

$$T(n) \leq cn^2 \left(\frac{1}{1 - \frac{3}{16}}\right) + \left(\frac{5}{3}\right) n^{\log_4 3}$$

$$T(n) \in O(n^2)$$

Solving Recurrences II

Data Structures and Parallelism

Warm Up

Write a recurrence to describe the running time of Mystery.

If you have extra time, find the big-Θ running time.

```
Mystery(int[] arr){
    if(arr.length == 1)
        return arr[0];
    else if(arr.length == 2)
        return arr[0] + arr[1];
    //copies all but first two elements of arr.
    int[] smaller = Arrays.copyOfRange(arr, 2, arr.length);
    return a[0] + a[1] + Mystery(smaller);
}
```

Warm Up

$$T(n) = \begin{cases} 1 & \text{if } n = 1 \\ 2 & \text{if } n = 2 \\ T(n-2) + 4 & \text{otherwise} \end{cases}$$

Is there an easier way?

We do all that effort to get an exact formula for the number of operations,

But we usually only care about the Θ bound.

There must be an easier way

Sometimes, there is!

Master Theorem

Given a recurrence of the following form:

$$T(n) = \begin{cases} d & when\ n \leq \text{some constant} \\ aT\left(\frac{n}{b}\right) + n^c & \text{otherwise} \end{cases}$$

Where a, b, c, and d are all constants.
The big-theta solution always follows this pattern:

If $\log_b a < c$ then $T(n)$ is $\Theta(n^c)$

If $\log_b a = c$ then $T(n)$ is $\Theta(n^c \log n)$

If $\log_b a > c$ then $T(n)$ is $\Theta\left(n^{\log_b a}\right)$

Apply Master Theorem

Given a recurrence of the form:

$$T(n) = \begin{cases} d \text{ when } n \leq \text{some constant} \\ aT\left(\frac{n}{b}\right) + n^c \text{ otherwise} \end{cases}$$

If $\log_b a < c$ then $T(n)$ is $\Theta(n^c)$

If $\log_b a = c$ then $T(n)$ is $\Theta(n^c \log n)$

If $\log_b a > c$ then $T(n)$ is $\Theta\left(n^{\log_b a}\right)$

$$T(n) = \begin{cases} 1 \text{ when } n \leq 1 \\ 2T\left(\frac{n}{2}\right) + n \text{ otherwise} \end{cases}$$

a = 2
b = 2
c = 1
d = 1

$$\log_b a = c \Rightarrow \log_2 2 = 1$$

$$T(n) \text{ is } \Theta(n^c \log_2 n) \Rightarrow \Theta(n^1 \log_2 n)$$

Reflecting on Master Theorem

Given a recurrence of the form:

$$T(n) = \begin{cases} d \text{ when } n \leq \text{some constant} \\ aT\left(\frac{n}{b}\right) + n^c \text{ otherwise} \end{cases}$$

If $\log_b a < c$ then $T(n)$ is $\Theta(n^c)$

If $\log_b a = c$ then $T(n)$ is $\Theta(n^c \log n)$

If $\log_b a > c$ then $T(n)$ is $\Theta\left(n^{\log_b a}\right)$

$height \approx \log_b a$

$branchWork \approx n^c \log_b a$

$leafWork \approx d\left(n^{\log_b a}\right)$

The $\log_b a < c$ case

- Recursive case conquers work more quickly than it divides work
- Most work happens near "top" of tree
- Non recursive work in recursive case dominates growth, n^c term

The $\log_b a = c$ case

- Work is equally distributed across levels of the tree
- Overall work is approximately work at any level x height

The $\log_b a > c$ case

- Recursive case divides work faster than it conquers work
- Most work happens near "bottom" of tree
- Work at base case dominates.

Benefits of Solving By Hand

If we had the Master Theorem why did we do all that math???

Not all recurrences fit the Master Theorem.

- Recurrences show up everywhere in computer science.
- And they're not always nice and neat.

It helps to understand exactly where you're spending time.

- Master Theorem gives you a very rough estimate. The Tree Method can give you a much more precise understanding.

Unrolling

There's an alternative to the tree method called "unrolling"

Instead of going through the process of making the tree and finding the work at each level, sometimes we want to just plug the formula into itself until we see a pattern.

Benefits: No drawing, no big table.

Drawbacks: No drawing, no big table – harder to find patterns.

Next week's section handout will have more examples.

Solving Recurrences IV:

$$T(n,k) = \begin{cases} 2T\left(n - \left(\frac{n}{k}\right)^{\frac{3}{2}}, k\right) & \text{if } n \geq 2k \\ 2^{2k} & \text{otherwise} \end{cases}$$

Show $T(n^2, n) \leq 2^{cn}$ where c is some constant.

This recurrence is the main point of the paper Robbie's presenting at the conference.

You actually know enough to find a good upper bound!

(but not a $\Theta()$ bound)

Hint: don't try to find the exact number of levels.

Instead bound the number of levels needed to cut the first parameter in half.

Amortization

What's the worst case for inserting into an ArrayList?

- O(n). If the array is full.

Is O(n) a good description of the worst case behavior?

- If you're worried about a single insertion, maybe.
- If you're worried about doing, say, n insertions in a row. NO!

Amortized bounds let us study the behavior of a bunch of consecutive calls.

Amortization

The most common application of amortized bounds is for insertions/deletions and data structure resizing.

Let's see why we always do that doubling strategy.

How long in total does it take to do n insertions?

We might need to double a bunch, but the total resizing work is at most O(n)

And the regular insertions are at most $n \cdot O(1) = O(n)$

So n insertions take $O(n)$ work total

Or amortized $O(1)$ time.

Amortization

Why do we double? Why not increase the size by 10,000 each time we fill up?

How much work is done on resizing to get the size up to n?

Will need to do work on order of current size every 10,000 inserts

$$\sum_{i=0}^{\frac{n}{10000}} 10000i \approx 10{,}000 \cdot \frac{n^2}{10{,}000^2} = O(n^2)$$

The other inserts do $O(n)$ work total.

The amortized cost to insert is $O\left(\frac{n^2}{n}\right) = O(n)$.

Much worse than the $O(1)$ from doubling!

Project Checkpoint

Get into your project groups

Once you talk to one of us you're free to go.