

QuickSort complete parallel

iteration

$i=0$

$$T(n) = T\left(\frac{n}{2}\right) + \log(n)$$

$i=1$

1. Expansion

$$\rightarrow T\left(\frac{n}{2}\right) = T\left(\frac{n}{4}\right) + \log\left(\frac{n}{2}\right)$$

$$= T\left(\frac{n}{4}\right) + \log\left(\frac{n}{2}\right) + \log(n)$$

$i=2$

$$\rightarrow T\left(\frac{n}{4}\right) = T\left(\frac{n}{8}\right) + \log\left(\frac{n}{4}\right)$$

$$= T\left(\frac{n}{8}\right) + \log\left(\frac{n}{4}\right) + \log\left(\frac{n}{2}\right) + \log(n)$$

$i=3$

$$\log\left(\frac{a}{b}\right) = \log(a) - \log(b) \quad \text{expand logs}$$

$$= T\left(\frac{n}{8}\right) + (\log(n) - \log(4)) + (\log(n) - \log(2)) + \log n$$

$$\text{simplify logs} \quad 2 \quad \text{b/c } \log_2 2^i = i \quad 1$$

$$= T\left(\frac{n}{8}\right) + 3\log(n) - (2+1)$$

↓ summation

2. Find pattern

$$= T\left(\frac{n}{2^i}\right) + i\log(n) - \sum_{j=0}^{i-1} j$$

$$\rightarrow \text{Solve } T(1) = T\left(\frac{n}{2^i}\right) = C_0 \rightarrow \frac{n}{2^i} = 1$$

3. Solve for i

$$\rightarrow i = \log n$$

4. substitute in i & simplify

$$= C_0 + \log(n) \cdot \log(n) - \sum_{j=0}^{\log n - 1} j$$

→ solve Gauss's summation (see handout)

$$= \frac{(\log n - 1)(\log n)}{2}$$

$$= C_0 + \log^2(n) - \frac{1}{2} (\log(n)\log(n) - \log(n))$$

$$= \frac{1}{2} \log^2(n) + \frac{1}{2} \log(n)$$

→ worse runtime

$$= O(\log^2(n))$$