CSE 332: Data Structures & Parallelism

Lecture 19: Introduction to Graphs

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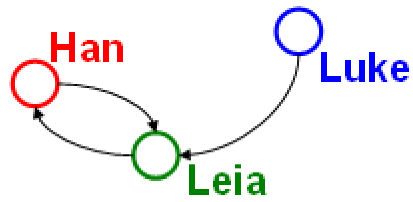
Today

- Graphs
  - Intro & Definitions

Graphs

- A graph is a formalism for representing relationships among items
  - Very general definition because very general concept
- A graph is a pair
  `G = (V,E)`
  - A set of vertices, also known as nodes
    $V = \{v_1, v_2, \ldots, v_n\}$
  - A set of edges
    $E = \{e_1, e_2, \ldots, e_m\}$
    - Each edge e_i is a pair of vertices (v_j, v_k)
    - An edge "connects" the vertices
- Graphs can be directed or undirected




```
V = {Han,Leia,Luke}
E = {(Luke,Leia),
     (Han,Leia),
     (Leia,Han)}
```

An ADT?

- Can think of graphs as an ADT with operations like `isEdge((v_j,v_k))` ($\texttt{isEdge}((v_j, v_k))$)
- But it is unclear what the "standard operations" are
- Instead we tend to develop algorithms over graphs and then use data structures that are efficient for those algorithms
- Many important problems can be solved by:
  1. Formulating them in terms of graphs
  2. Applying a standard graph algorithm
- To make the formulation easy and standard, we have a lot of *standard terminology* about graphs

Some graphs

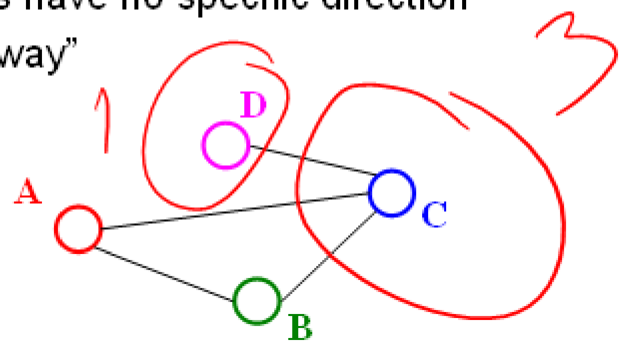
For each, what are the vertices and what are the edges?

- Web pages with links
- Facebook friends
- "Input data" for the Kevin Bacon game
- Methods in a program that call each other
- Road maps (e.g., Google maps)
- Airline routes
- Family trees
- Course pre-requisites
- …

Wow: Using the same algorithms for problems across so many domains sounds like "core computer science and engineering"

Undirected Graphs

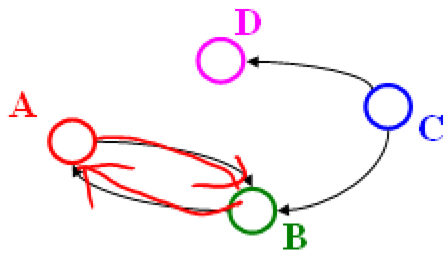
- In undirected graphs, edges have no specific direction
  - Edges are always "two-way"



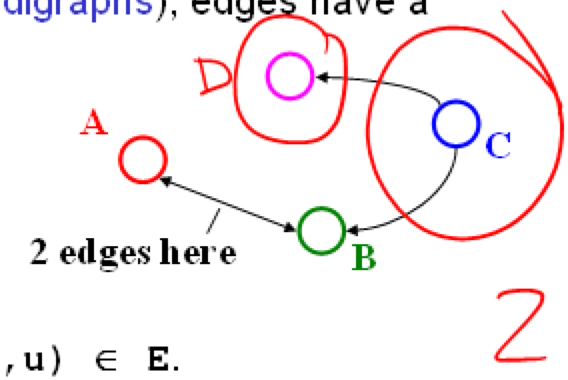
- Thus, `(u,v)` $\in$ **`E`** implies `(v,u)` $\in$ **`E`**.
  - Only one of these edges needs to be in the set; the other is implicit
- Degree of a vertex: number of edges containing that vertex
  - Put another way: the number of adjacent vertices

Directed Graphs

- In directed graphs (sometimes called digraphs), edges have a direction

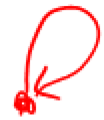


or



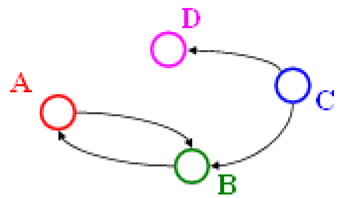
- Thus, $(u,v) \in E$ does *not* imply $(v,u) \in E$.
  - Let $(u,v) \in E$ mean $u \rightarrow v$
  - Call u the source and v the destination
- In-Degree of a vertex: number of in-bound edges, i.e., edges where the vertex is the destination
- Out-Degree of a vertex: number of out-bound edges i.e., edges where the vertex is the source

Self-edges, connectedness



- A self-edge a.k.a. a loop is an edge of the form `(u,u)`
  - Depending on the use/algorithm, a graph may have:
    - No self edges
    - Some self edges
    - All self edges (often therefore implicit, but we will be explicit)
- A node can have a degree / in-degree / out-degree of zero
- A graph does not have to be connected (In an undirected graph, this means we can follow edges from any node to every other node), even if every node has non-zero degree

More notation



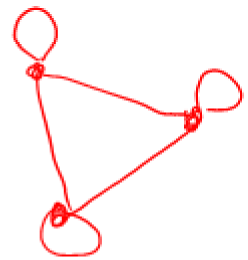
For a graph `G = (V,E)`:

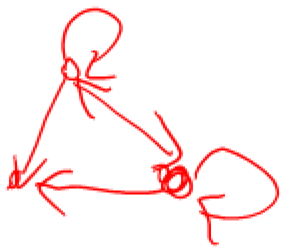
- `|V|` is the number of vertices
- `|E|` is the number of edges
  - Minimum? 0
  - Maximum for undirected? $\frac{V(V+1)}{2}$
  - Maximum for directed? V^2

$V = \{A, B, C, D\}$

$E = \{(C, B), (A, B), (B, A), (C, D)\}$

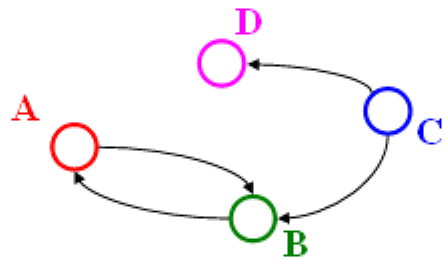
- If $(u,v) \in E$
  - Then v is a neighbor of u, i.e., v is adjacent to u
  - Order matters for directed edges
    - u is not adjacent to v unless $(v,u) \in E$





$3+2+1$

More notation



For a graph `G = (V,E)`:

- `|V|` is the number of vertices
- `|E|` is the number of edges
  - Minimum? 0
  - Maximum for undirected? `|V||V+1|/2` $\in$ $O(|V|^2)$
  - Maximum for directed? $|V|^2$ $\in$ $O(|V|^2)$

    (For both undirected and directed, assuming self-edges are allowed, else subtract `|V|` from the answers above)

- If `(u,v)` $\in$ `E`
  - Then `v` is a neighbor of `u`, i.e., `v` is adjacent to `u`
  - Order matters for directed edges
    - `u` is not adjacent to `v` unless `(v,u)` $\in$ `E`

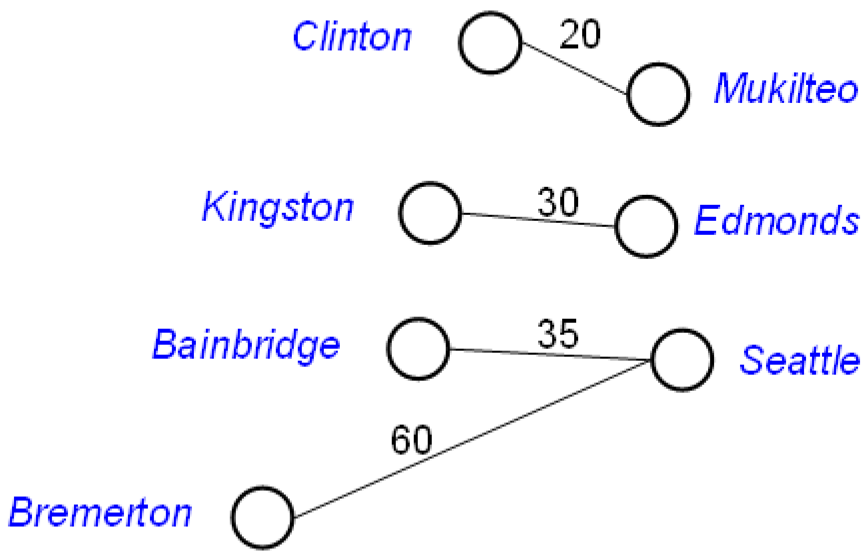
Examples again

Which would use directed edges? Which would have self-edges?
Which could have 0-degree nodes?

- Web pages with links
- Facebook friends ——— undirected
- "Input data" for the Kevin Bacon game
- Methods in a program that call each other
- Road maps (e.g., Google maps)
- Airline routes
- Family trees
- Course pre-requisites
- …

Weighted graphs

- In a weighed graph, each edge has a weight a.k.a. cost
  - Typically numeric (most examples will use ints)
  - *Orthogonal* to whether graph is directed
  - Some graphs allow *negative weights*; many don't



Examples

What, if anything, might weights represent for each of these? Do negative weights make sense?

- Web pages with links
- Facebook friends
- “Input data” for the Kevin Bacon game
- Methods in a program that call each other
- Road maps (e.g., Google maps)
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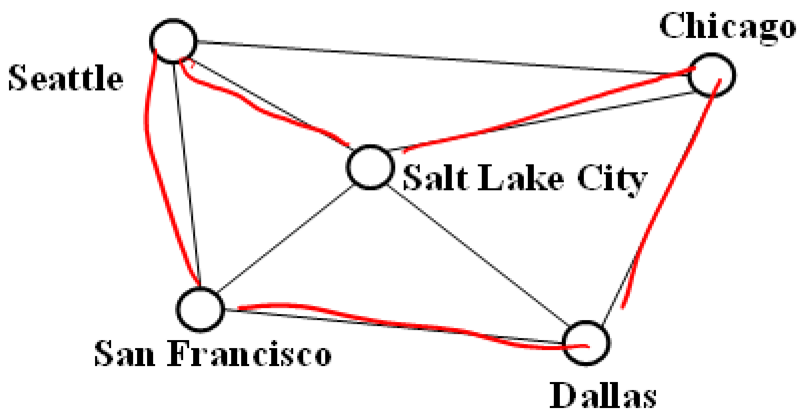
- …

Paths and Cycles

- A path is a list of vertices $[v_0, v_1, \ldots, v_n]$ such that $(v_i, v_{i+1}) \in E$ for all $0 \le i < n$. Say *"a path from* v_0 *to* v_n*"*

- A cycle is a path that begins and ends at the same node ($v_0 == v_n$)



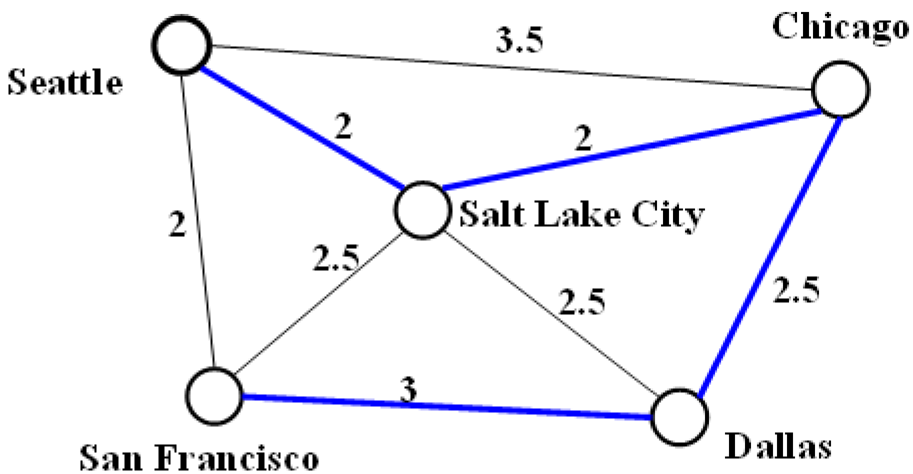
Example path (that also happens to be a cycle):
[Seattle, Salt Lake City, Chicago, Dallas, San Francisco, Seattle]

Path Length and Cost

- Path length: Number of *edges* in a path (also called "unweighted cost")
- Path cost: Sum of the weights of each edge

Example where:

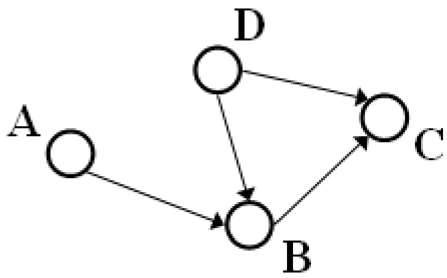
P= [Seattle, Salt Lake City, Chicago, Dallas, San Francisco]



length(P) = 4
cost(P) = 9.5

Paths/cycles in directed graphs

Example:



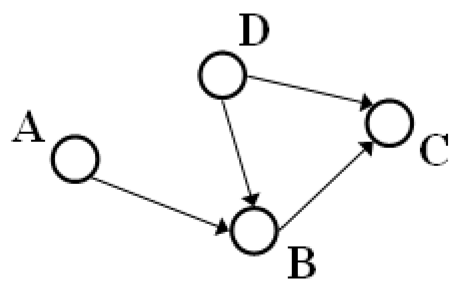
Undirected

Is there a path from A to D? No

Does the graph contain any cycles? No

Paths/cycles in directed graphs

Example:

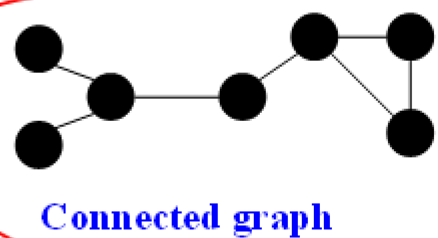


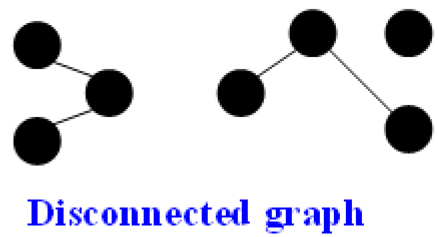
Is there a path from A to D? No

Does the graph contain any cycles? No

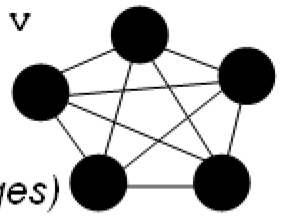
Undirected graph connectivity

- An undirected graph is connected if for all pairs of vertices u , v, there exists a *path* from u to v



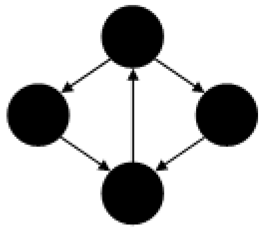


- An undirected graph is complete, a.k.a. fully connected if for all pairs of vertices u , v, there exists an *edge* from u to v

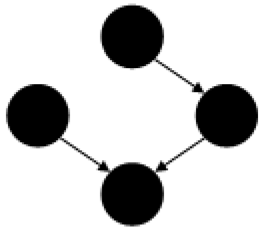


Directed graph connectivity

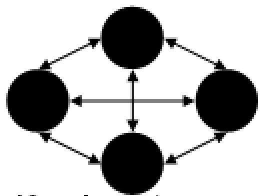
- A directed graph is strongly connected if there is a path from every vertex to every other vertex



- A directed graph is weakly connected if there is a path from every vertex to every other vertex *ignoring direction of edges*



- A complete a.k.a. fully connected directed graph has an edge from every vertex to every other vertex



(plus self edges)

Examples

For undirected graphs: connected?

For directed graphs: strongly connected? weakly connected?

- Web pages with links
- Facebook friends
- "Input data" for the Kevin Bacon game
- Methods in a program that call each other
- Road maps (e.g., Google maps)
- Airline routes
- Family trees
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- …

Trees as graphs

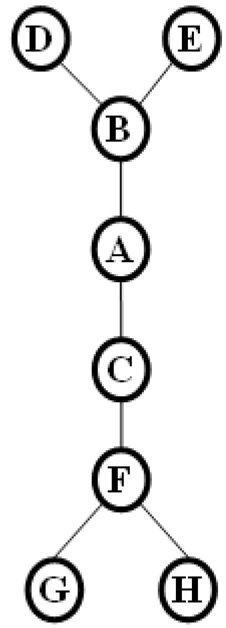
When talking about graphs,
we say a tree is a graph that is:

- undirected
- acyclic
- connected

So all trees are graphs, but not all graphs are trees

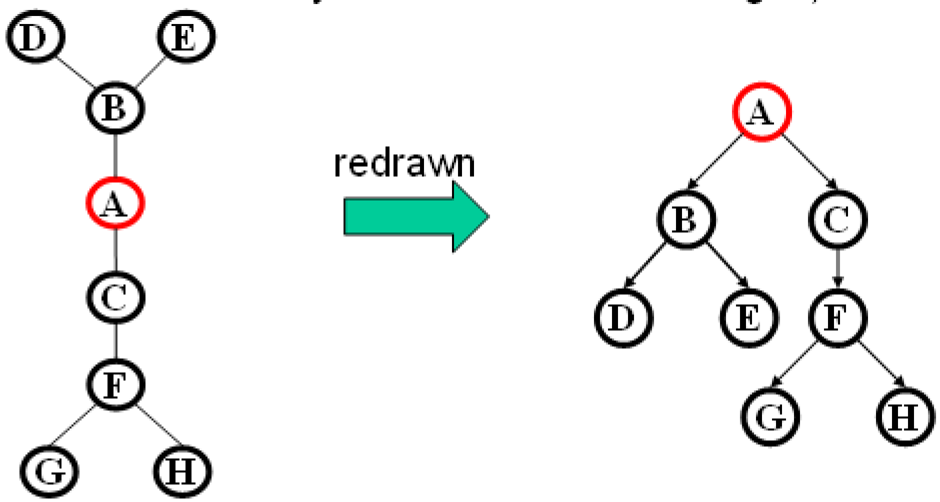
How does this relate to the trees we know and love?...

Example:



Rooted Trees

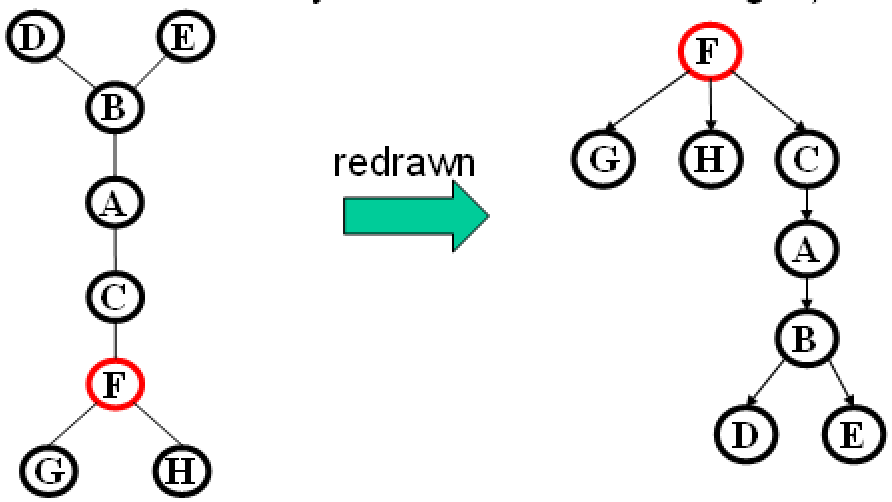
- We are more accustomed to rooted trees where:
  - We identify a unique ("special") root
  - We think of edges as **directed**: parent to children
- Given a tree, once you pick a root, you have a unique rooted tree (just drawn differently and with undirected edges)



Rooted Trees (Another example)

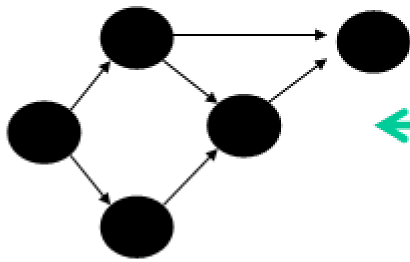
- We are more accustomed to rooted trees where:
  - We identify a unique (“special”) root
  - We think of edges as **directed**: parent to children

- Given a tree, once you pick a root, you have a unique rooted tree (just drawn differently and with undirected edges)



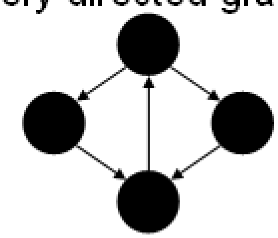
Directed acyclic graphs (DAGs)

- A DAG is a directed graph with no (directed) cycles
  - Every rooted directed tree is a DAG
    - But not every DAG is a rooted directed tree:



Not a rooted directed tree,
Has a cycle (in the
undirected sense)

  - Every DAG is a directed graph
    - But not every directed graph is a DAG:

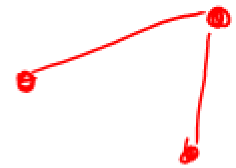


Examples

Which of our directed-graph examples do you expect to be a DAG?

- Web pages with links
- "Input data" for the Kevin Bacon game
- Methods in a program that call each other
- Airline routes
- Family trees
- Course pre-requisites
- …

Density / sparsity



- Recall: In an undirected graph, $0 \leq |E| < |V|^2$
- Recall: In a directed graph: $0 \leq |E| \leq |V|^2$
- So for any graph, $|E|$ is $O(|V|^2)$
- One more fact: If an undirected graph is *connected*, then $|E| \geq |V|-1$
- Because $|E|$ is often much smaller than its maximum size, we do not always approximate as $|E|$ as $O(|V|^2)$
  – This is a correct bound, it just is often not tight
  – If it is tight, i.e., $|E|$ is $\Theta(|V|^2)$ we say the graph is dense
    - More sloppily, dense means "lots of edges"
  – If $|E|$ is $O(|V|)$ we say the graph is sparse
    - More sloppily, sparse means "most (possible) edges missing"

What is the Data Structure?

- So graphs are really useful for lots of data and questions
  - For example, “what’s the lowest-cost path from x to y”

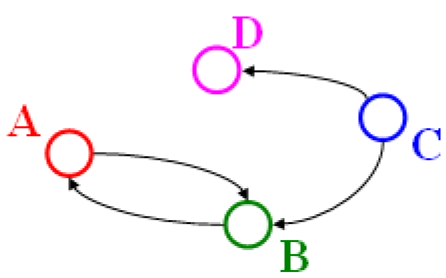
- But we need a data structure that represents graphs

- The “best one” can depend on:
  - Properties of the graph (e.g., dense versus sparse)
  - The common queries (e.g., “is `(u,v)` an edge?” versus “what are the neighbors of node `u`?”)

- So we’ll discuss the two standard graph representations
  - Adjacency Matrix and Adjacency List
  - Different trade-offs, particularly time versus space

Adjacency matrix

- Assign each node a number from `0` to `|V|-1`
- A |V| x |V| matrix (i.e., 2-D array) of Booleans (or 1 vs. 0)
  - If `M` is the matrix, then `M[u][v] == true` means there is an edge from `u` to `v`



To

From	A	B	C	D
A	F	T	F	F
B	T	F	F	F
C	F	T	F	T
D	F	F	F	F

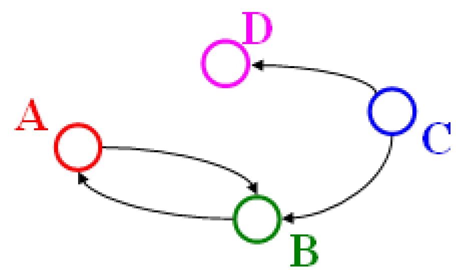
Adjacency Matrix Properties

- Running time to:
  - Get a vertex's out-edges:
  - Get a vertex's in-edges:
  - Decide if some edge exists:
  - Insert an edge:
  - Delete an edge:

- Space requirements:

- Best for sparse or dense graphs?

	A	B	C	D
A	F	T	F	F
B	T	F	F	F
C	F	T	F	T
D	F	F	F	F



Adjacency Matrix Properties

- Running time to:
  - Get a vertex's out-edges: $O(|V|)$
  - Get a vertex's in-edges: $O(|V|)$
  - Decide if some edge exists: $O(1)$
  - Insert an edge: $O(1)$
  - Delete an edge: $O(1)$

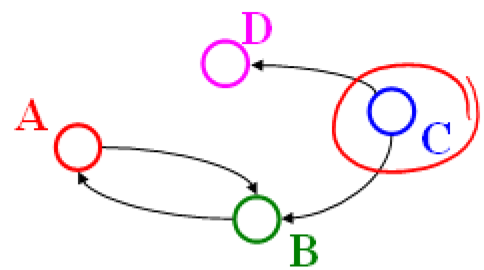
- Space requirements:
  - $|V|^2$ bits

- Best for sparse or dense graphs?
  - Best for dense graphs

To

from

	A	B	C	D
A	F	T	F	F
B	T	F	F	F
C	F	T	F	T
D	F	F	F	F



Adjacency Matrix Properties

- How will the adjacency matrix vary for an *undirected graph*?

- How can we adapt the representation for *weighted graphs*?

	A	B	C	D
A	F	T	F	F
B	T	F	F	F
C	F	T	F	T
D	F	F	F	F

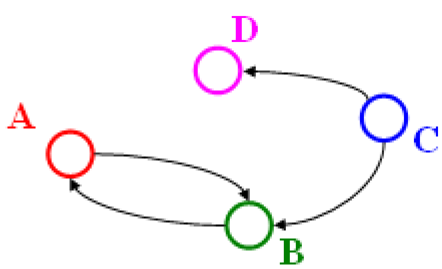
Adjacency Matrix Properties

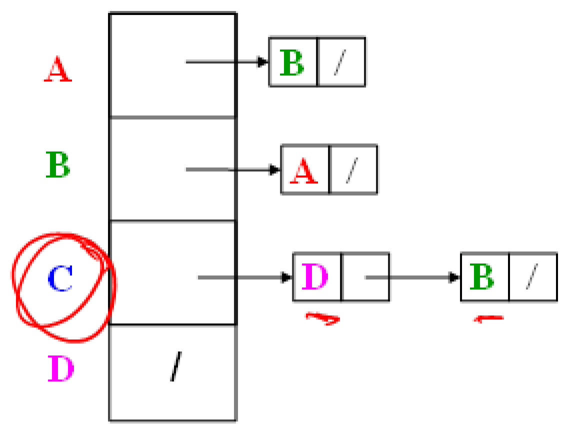
- How will the adjacency matrix vary for an *undirected graph*?
  - Undirected will be symmetric about diagonal axis
- How can we adapt the representation for *weighted graphs*?
  - Instead of a Boolean, store a number in each cell
  - Need some value to represent 'not an edge'
    - In *some* situations, 0 or -1 works

	A	B	C	D
A	F	T	F	F
B	T	F	F	F
C	F	T	F	T
D	F	F	F	F

Adjacency List

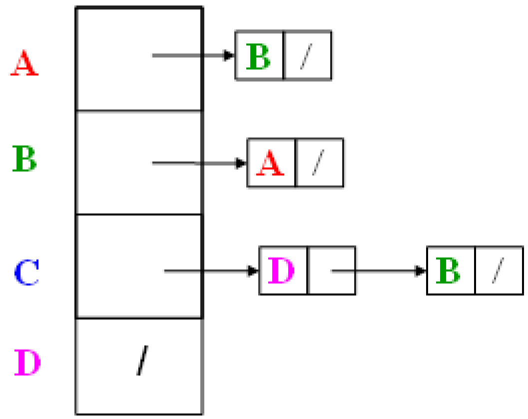
- Assign each node a number from `0` to `|V|-1`
- An array of length `|V|` in which each entry stores a list of all adjacent vertices (e.g., linked list)

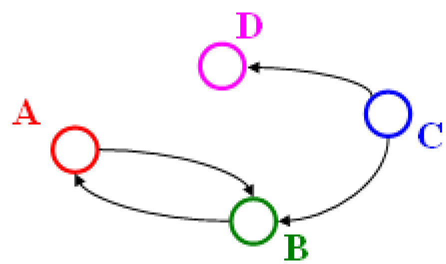




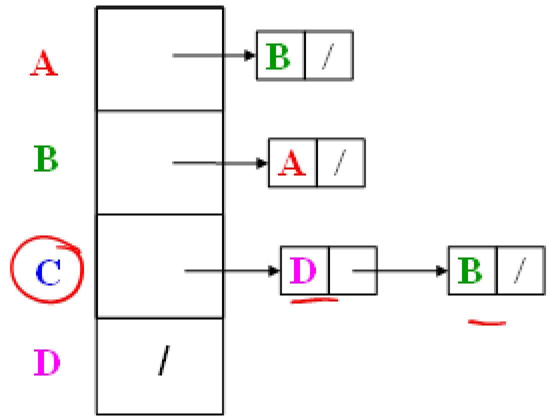
Adjacency List Properties

- Running time to:
  - Get all of a vertex's out-edges:
  - Get all of a vertex's in-edges:
  - Decide if some edge exists:
  - Insert an edge:
  - Delete an edge:
- Space requirements:
- Best for dense or sparse graphs?

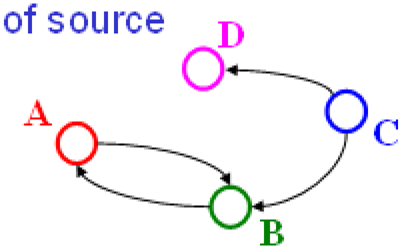




Adjacency List Properties



- Running time to:
  - Get all of a vertex's out-edges:
    $O(d)$ where d is out-degree of vertex
  - Get all of a vertex's in-edges:
    $O(|V| + |E|)$ (but could keep a second adjacency list for this!)
  - Decide if some edge exists:
    $O(d)$ where d is out-degree of source
  - Insert an edge: $O(1)$ (unless you need to check if it's there)
  - Delete an edge: $O(d)$ where d is out-degree of source



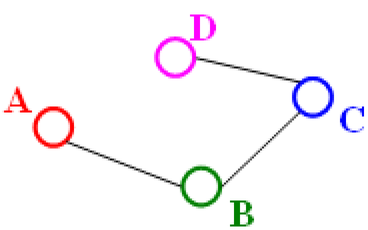
- Space requirements:
  - $O(|V|+|E|)$
- Best for dense or sparse graphs?
  - Best for sparse graphs, so usually just stick with linked lists

Undirected Graphs

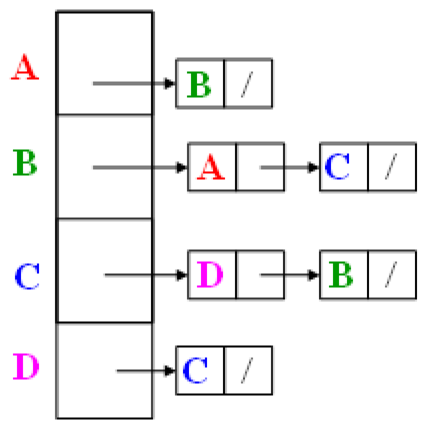
Adjacency matrices & adjacency lists both do fine for undirected graphs

- Matrix: Can save roughly ½ the space
  - But may slow down operations in languages with “proper” 2D arrays (not Java, which has only arrays of arrays)
  - How would you “get all neighbors”?
- Lists: Each edge in two lists to support efficient “get all neighbors”

Example:



	A	B	C	D
A	F	T	F	F
B	T	F	T	F
C	F	T	F	T
D	F	F	T	F



Which is better?

Graphs are often sparse:

- Streets form grids
  - every corner is not connected to every other corner
- Airlines rarely fly to all possible cities
  - or if they do it is to/from a hub rather than directly to/from all small cities to other small cities

Adjacency lists should generally be your default choice

- Slower performance compensated by greater space savings

Next…

Okay, we can represent graphs

Now let's implement some useful and non-trivial algorithms

- Topological sort: Given a DAG, order all the vertices so that every vertex comes before all of its neighbors
- Shortest paths: Find the shortest or lowest-cost path from x to y
  - Related: Determine if there even is such a path