

CSE 332

Data Abstractions

P vs. NP: Efficient Reductions Between Problems

More Graph Problems

1

Let's consider the **longest path** problem on a graph.

Remember, we were able to do **shortest paths** using Dijkstra's.

Take a few minutes to try to solve the **longest path** problem.

Decision Problems

2

Definition (Decision Problem)

A **decision problem** (or **language**) is a set of strings ($L \subseteq \Sigma^*$).

An algorithm (from Σ^* to `boolean`) solves a decision problem when it outputs `true` iff the input is in the set.

PRIMES

Input(s): Number x
Output: `true` iff x is prime

An Algorithm that solves PRIMES

```

1 isPrime(x) {
2   for (i = 2; i < x; i++) {
3     if (x % i == 0) {
4       return false;
5     }
6   }
7   return true;
8 }
```

Efficient?

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In this lecture, we'll be talking about **efficient reductions**. So, naturally, we have to answer two questions:

- What is an efficient algorithm?
- What is a reduction?

Efficient Algorithm

We say an algorithm is **efficient** if the worst-case analysis is a **polynomial**. Okay, but...

- $n^{10000000}$ is polynomial
- $3000000000000000n^3$ is polynomial

Are those really efficient?

Well, no, but, in practice...

when a polynomial algorithm is found the constants are actually low

Polynomial runtime is a **very** low bar, if we can't even get that...

Reductions

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This lecture is about exposing **hidden** similarities between problems.

We will show that problems that are **cosmetically different** are **substantially the same!**

Our main tool to do this is called a **reduction**:

Reductions

We have two **decision problems**, **A** and **B**. To show that **A** is "at least as hard as" **B**, we

- Suppose we can solve **A**
- Create an algorithm, which calls **A** as a method, to solve **B**

To show they're the same, we have to do both directions.

Longest Paths and HAM!

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Two New Computational Problems

LONG-PATH

Input(s): Unweighted Graph G ; Number k
Output: true iff G has a path with k edges

HAM-PATH

Input(s): Unweighted Graph G
Output: true iff G has a path using all vertices

Suppose we could solve LONG-PATH...

"Algorithm"

```
1 HAM-PATH(G) {
2   return LONG-PATH(G, |V| - 1)
3 }
```

Suppose we could solve HAM-PATH...

"Algorithm"

```
1 LONG-PATH(G, k) {
2   for ( $G' = (v_1, v_2, \dots, v_k)$  in  $G$ ) {
3     if (HAM-PATH( $G'$ )) {
4       return true;
5     }
6   }
7   return false;
8 }
```

A 2-CRAYOLA Question

6

Definition (k -coloring)

A k -coloring of a graph G is an assignment of k colors to vertices such that no two adjacent vertices have the same color.

2-COLOR

Input(s): Graph G
Output: true iff G has a valid 2-coloring

Can we solve this?

Can we solve this efficiently?

Algorithm For 2-COLOR

Try all 2^n possible colorings of the input graph!

Efficient Algorithm For 2-COLOR

Do a dfs on the graph! Every time we hit a vertex, assign it the opposite color from the vertex we just visited. If there's a color conflict, output false. If we finish with no color conflict, output true.

A 3-CRAYOLA Question

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Definition (k -coloring)

A k -coloring of a graph G is an assignment of k colors to vertices such that no two adjacent vertices have the same color.

3-COLOR

Input(s): Graph G
Output: true iff G has a valid 3-coloring

Inefficient Algorithm For 3-COLOR

Try all 3^n possible colorings of the input graph!

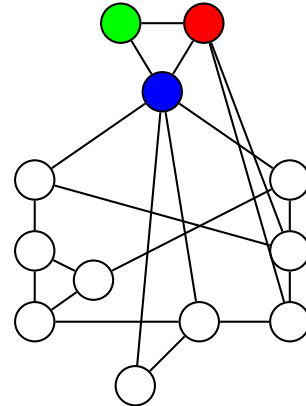
Efficient Algorithm For 3-COLOR

UNKNOWN

A Graph Called "Gadget"

8

Find a valid 3-coloring of this graph. To orient ourselves, I've started it:



Another Decision Problem!

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CIRCUITSAT

Input(s): n -Input/1-Output Circuit C
Output: true iff C has a satisfying assignment

Inefficient Algorithm For CIRCUITSAT

Try all 2^n possible assignments of variables

Efficient Algorithm For CIRCUITSAT

UNKNOWN

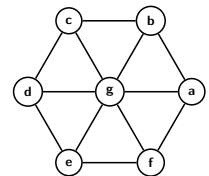
Suspicious...

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CIRCUITSAT

X Y Z
 OR NOT
 OR
 OUT

3-COLOR

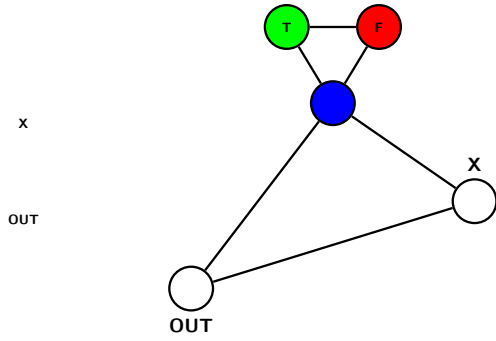


We don't know how to solve either of these problems...

Could they be the same problem in disguise?

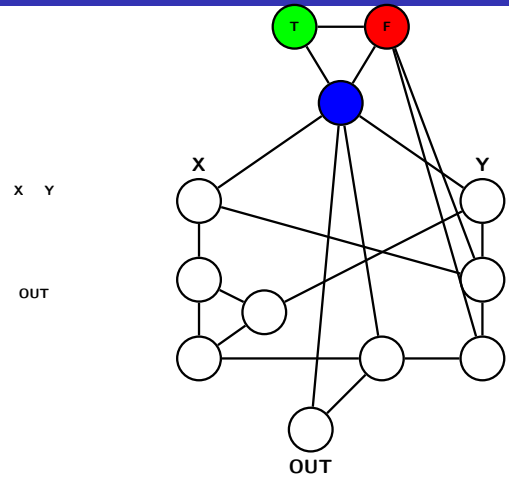
Not Gadget with Labels

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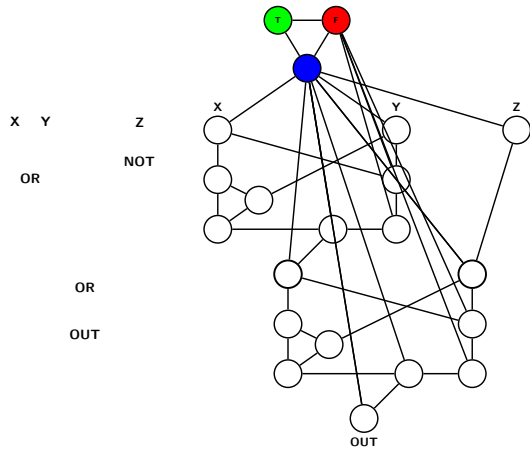
Or Gadget with Labels

12



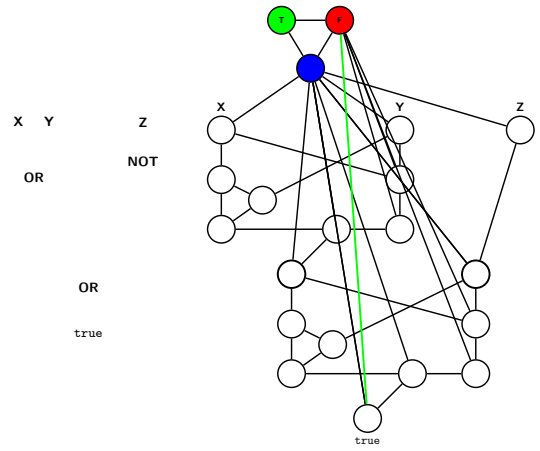
Circuit

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SATISFIABLE Circuit

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Lesson

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We found a way to "emulate" circuit satisfiability using three coloring!

If we can find a solution to 3-COLOR, we can solve CIRCUITSAT quickly.

These problems are substantially the same