CSE 332:
Intro to Parallelism: Multithreading and Fork-Join

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Announcements (2/19/2014)

- HW #5 due today
- HW #6 out today
- Read Grossman 2.1-3.4

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Sequential

- Sum up N numbers in an array
  - Complexity?

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Parallel Sum

- Sum up N numbers in an array
  - with two processors

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Parallel Sum

- Sum up N numbers in an array
  - with N processors?

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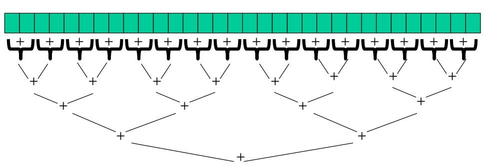
Parallel Sum

- Sum up N numbers in an array



- Complexity?
- How many processors?
- Faster with infinite processors?

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Changing a Major Assumption

- So far, we have assumed:

  One thing happens at a time

- Called sequential programming
- Dominated until roughly 2005
  - what changed?

A Simplified History

From roughly 1980-2005, desktop computers got exponentially faster at running sequential programs
- About twice as fast every couple years

Writing parallel (multi-threaded) code is harder than sequential
- Especially in common languages like Java and C

But nobody knows how to continue this
- Increasing clock rate generates too much heat
- Relative cost of memory access is too high
- But we can keep making “wires exponentially smaller” (Moore’s “Law”), so put multiple processors on the same chip (“multicore”)

Who Implements Parallelism

- User
- Application
- Operating System
- Programming Language, Compiler
- Algorithm
- Processor Hardware

Parallelism vs. Concurrency

Parallelism:
Use extra resources to solve a problem faster

Concurrency:
Manage access to shared resources

An analogy

A program is like a recipe for a cook
- Sequential: one cook who does one thing at a time

Parallelism: (Let’s get the job done faster!)
- Have lots of potatoes to slice?
- Hire helpers, hand out potatoes and knives
- But too many chefs and you spend all your time coordinating

Concurrency: (We need to manage a shared resource)
- Lots of cooks making different things, but only 4 stove burners
- Want to allow access to all 4 burners, but not cause spills or incorrect burner settings

Shared Memory with Threads

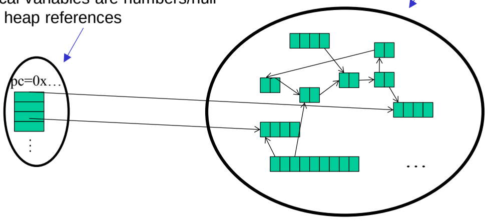
Old story: A running program has
- One *program counter* (current statement executing)
- One *call stack* (with each *stack frame* holding local variables)
- *Objects in the heap* created by memory allocation (i.e., `new`)
  - (nothing to do with data structure called a heap)
- *Static fields*

New story:
- A set of *threads*, each with its own program counter & call stack
  - No access to another thread’s local variables
- Threads can share static fields / objects
  - To *communicate*, write values to some shared location that another thread reads from

Old Story: one call stack, one pc

- Call stack with local variables
- pc determines current statement
- local variables are numbers/null or heap references

Heap for all objects and static fields



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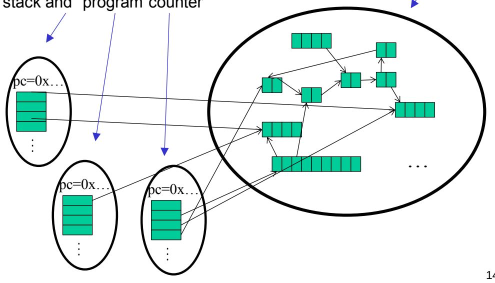
New Story: Shared Memory with Threads

Heap for all objects and static fields, *shared* by all threads

Threads, each with own *unshared* call stack and "program counter"



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Other models

We will focus on shared memory, but you should know several other models exist and have their own advantages **(see notes)**

- Message-passing: Each thread has its own collection of objects. Communication is via explicitly sending/receiving messages
  - Cooks working in separate kitchens, mail around ingredients
- Dataflow: Programmers write programs in terms of a DAG. A node executes after all of its predecessors in the graph
  - Cooks wait to be handed results of previous steps
- Data parallelism: Have primitives for things like "apply function to every element of an array in parallel"

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Our Needs

To write a shared-memory parallel program, need new primitives from a programming language or library

- Ways to create and *run multiple things at once*
  - Let's call these things **threads**
- Ways for threads to *share memory*
  - Often just have threads with references to the same objects
- Ways for threads to *coordinate (a.k.a. synchronize)*
  - For now, a way for one thread to wait for another to finish
  - Other primitives when we study concurrency

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Threads vs. Processors

What happens if you start 5 threads on a machine with only 4 processors (cores)?

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Threads vs. Processors

For sum operation:
- with 3 processors available, using 4 threads would take 50% more time than 3 threads

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Fork-Join Parallelism

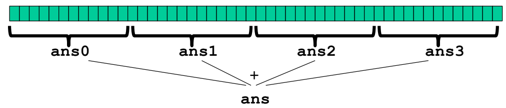
1. Define thread
   - Java: define subclass of `java.lang.Thread`, override `run`
2. Fork: instantiate a thread and start executing
   - Java: create thread object, call `start()`
3. Join: wait for thread to terminate
   - Java: call `join()` method, which returns when thread finishes

Above uses basic thread library build into Java

Later we'll introduce a better ForkJoin Java library designed for parallel programming

Sum with Threads

For starters: have 4 threads simultaneously sum ¼ of the array



- Create 4 *thread objects*, each given ¼ of the array
- Call `start()` on each thread object to run it in parallel
- Wait for threads to finish using `join()`
- Add together their 4 answers for the final result

Part 1: define thread class

```
class SumThread extends java.lang.Thread {

  int lo; // fields, passed to constructor
  int hi; // so threads know what to do.
  int[] arr;

  int ans = 0; // result

  SumThread(int[] a, int l, int h) {
    lo=l; hi=h; arr=a;
  }


  public void run() { //override must have this type
    for(int i=lo; i < hi; i++)
      ans += arr[i];
  }
}
```

Because we must override a no-arguments/no-result `run`, we use fields to communicate across threads

Part 2: sum routine

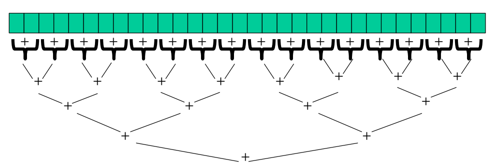
```
int sum(int[] arr){// can be a static method
  int len = arr.length;
  int ans = 0;
  SumThread[] ts = new SumThread[4];
  for(int i=0; i < 4; i++){// do parallel computations
    ts[i] = new SumThread(arr,i*len/4,(i+1)*len/4);
    ts[i].start();
  }
  for(int i=0; i < 4; i++) { // combine results
    ts[i].join(); // wait for helper to finish!
    ans += ts[i].ans;
  }
  return ans;
}
```


Parameterizing by number of threads

```
int sum(int[] arr, int numTs){
  int ans = 0;
  SumThread[] ts = new SumThread[numTs];
  for(int i=0; i < numTs; i++){
   ts[i] = new SumThread(arr,(i*arr.length)/numTs,
                            ((i+1)*arr.length)/numTs);
   ts[i].start();
  }
  for(int i=0; i < numTs; i++) {
    ts[i].join();
    ans += ts[i].ans;
  }
  return ans;
}
```


Recall: Parallel Sum

- Sum up N numbers in an array



- Let's implement this with threads...

Code looks something like this (using Java Threads)

```
class SumThread extends java.lang.Thread {
  int lo; int hi; int[] arr; // fields to know what to do
  int ans = 0; // result
  SumThread(int[] a, int l, int h) { … }
  public void run(){ // override
    if(hi – lo < SEQUENTIAL_CUTOFF)
      for(int i=lo; i < hi; i++)
        ans += arr[i];
    else {
      SumThread left = new SumThread(arr,lo,(hi+lo)/2);
      SumThread right= new SumThread(arr,(hi+lo)/2,hi);
      left.start();
      right.start();
      left.join(); // don't move this up a line – why?
      right.join();
      ans = left.ans + right.ans;
    }
  }
}
int sum(int[] arr){ // just make one thread!
   SumThread t = new SumThread(arr,0,arr.length);
   t.run();
   return t.ans;
}
```

Recursive problem decomposition

```
Thread: sum range [0,10)
  Thread: sum range [0,5)
     Thread: sum range [0,2)
        Thread: sum range [0,1) (return arr[0])
        Thread: sum range [1,2) (return arr[1])
        add results from two helper threads
     Thread: sum range [2,5)
        Thread: sum range [2,3) (return arr[2])
        Thread: sum range [3,5)
           Thread: sum range [3,4) (return arr[3])
           Thread: sum range [4,5) (return arr[4])
           add results from two helper threads
        add results from two helper threads
     add results from two helper threads
  Thread: sum range [5,10)
     Thread: sum range [5,7)
        Thread: sum range [5,6) (return arr[5])
        Thread: sum range [6,7) (return arr[6])
        add results from two helper threads
     Thread: sum range [7,10)
        Thread: sum range [7,8) (return arr[7])
        Thread: sum range [8,10)
           Thread: sum range [8,9) (return arr[8])
           Thread: sum range [9,10) (return arr[9])
           add results from two helper threads
        add results from two helper threads
  add results from two helper threads
```

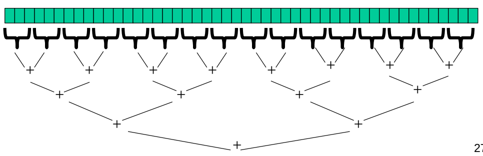
Example: summing an array with 10 elements. (too small to actually want to use parallelism)

The algorithm produces the following tree of recursion, where the range [i,j) includes i and excludes j:

Divide-and-conquer

Same approach useful for many problems beyond sum
- *If* you have enough processors, total time $O(\log n)$
- Next lecture: study reality of **P** << *n* processors

- Will write all our parallel algorithms in this style
  - But using a special fork-join library engineered for this style
    - Takes care of scheduling the computation well
  - Often relies on operations being associative (like +)



Thread Overhead

Creating and managing threads incurs cost

Two optimizations:

1. Use a *sequential cutoff*, typically around 500-1000
   - Eliminates lots of tiny threads
2. Do not create two recursive threads; create one thread and do the other piece of work "yourself"
   - Cuts the number of threads created by another 2x

Half the threads!

order of last 4 lines
Is critical – why?

```
// wasteful: don't
SumThread left  = …
SumThread right = …

left.start();
right.start();

left.join();
right.join();
ans=left.ans+right.ans;
```

```
// better: do!!
SumThread left  = …
SumThread right = …

left.start();
right.run();

left.join();
// no right.join needed
ans=left.ans+right.ans;
```

Note: `run` *is a normal function call! execution won't continue until we are done with* `run`

Better Java Thread Library

- Even with all this care, Java's threads are too "heavyweight"
  - Constant factors, especially space overhead
  - Creating 20,000 Java threads just a bad idea ☹
- The ForkJoin Framework is designed to meet the needs of divide-and-conquer fork-join parallelism
  - In the Java 7 standard libraries
    - (Also available for Java 6 as a downloaded `.jar` file)
  - Section will focus on pragmatics/logistics
  - Similar libraries available for other languages
    - C/C++: Cilk (inventors), Intel's Thread Building Blocks
    - C#: Task Parallel Library
    - …

Different terms, same basic idea

To use the ForkJoin Framework:

- A little standard set-up code (e.g., create a **`ForkJoinPool`**)

Don't subclass **`Thread`**	Do subclass **`RecursiveTask<V>`**
Don't override **`run`**	Do override **`compute`**
Do not use an **`ans`** field	Do return a **`V`** from **`compute`**
Don't call **`start`**	Do call **`fork`**
Don't *just* call **`join`**	Do call **`join`** (which returns answer)
Don't call **`run`** to hand-optimize	Do call **`compute`** to hand-optimize
Don't have a topmost call to **`run`**	Do create a pool and call **`invoke`**

See the web page for (linked in to project 3 description):
"A Beginner's Introduction to the ForkJoin Framework"

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Fork Join Framework Version:

```
class SumArray extends RecursiveTask<Integer> {
  int lo; int hi; int[] arr; // fields to know what to do
  SumArray(int[] a, int l, int h) { … }
  protected Integer compute(){// return answer
    if(hi – lo < SEQUENTIAL_CUTOFF) {
      int ans = 0; // local var, not a field
      for(int i=lo; i < hi; i++)
        ans += arr[i];
      return ans;
    } else {
      SumArray left = new SumArray(arr,lo,(hi+lo)/2);
      SumArray right= new SumArray(arr,(hi+lo)/2,hi);
      left.fork(); // fork a thread and calls compute
      int rightAns = right.compute();//call compute directly
      int leftAns  = left.join(); // get result from left
      return leftAns + rightAns;
    }
  }
}
static final ForkJoinPool fjPool = new ForkJoinPool();
int sum(int[] arr){
  return fjPool.invoke(new SumArray(arr,0,arr.length));
      // invoke returns the value compute returns
}
```