CSE332: Data Abstractions

Lecture 9: B Trees

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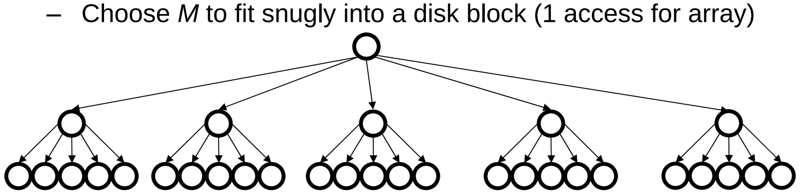
Our goal

- Problem: A dictionary with so much data most of it is on disk
- Desire: A balanced tree (logarithmic height) that is even shallower than AVL trees so that we can minimize disk accesses and exploit disk-block size
- A key idea: Increase the branching factor of our tree

M-ary Search Tree

- Build some sort of search tree with branching factor *M*:
  - Have an array of sorted children (**`Node[]`**)
  - Choose *M* to fit snugly into a disk block (1 access for array)



Perfect tree of height *h* has (M^{h+1}-1)/(*M*-1) nodes (textbook, page 4)

\# hops for **`find`**: If balanced, using $\log_M n$ instead of $\log_2 n$

- If *M*=256, that's an 8x improvement
- Example: *M* = 256 and *n* = 2^{40} that's 5 instead of 40

Runtime of **`find`** if balanced: $O(\log_2 M \log_M n)$ (binary search children)

Problems with M-ary search trees

- What should the order property be?
- How would you rebalance (ideally without more disk accesses)?
- Any "useful" data at the internal nodes takes up some disk-block space without being used by finds moving past it

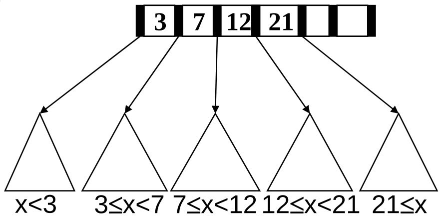
So let's use the branching-factor idea, but for a different kind of balanced tree

- Not a binary search tree
- But still logarithmic height for any *M* > 2

B+ Trees (we and the book say "B Trees")

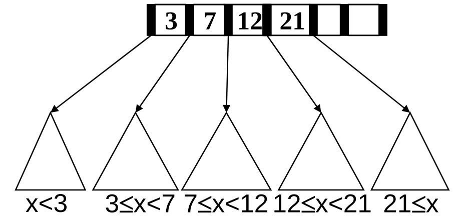
- Each internal node has room for up to *M-1* keys and *M* children
  - No other data; all data at the leaves!
- Order property:

  Subtree **between** keys *x* and *y* contains only data that is ≥ ***x*** and < ***y*** (notice the ≥)
- Leaf nodes have up to *L* sorted data items



- As usual, we'll ignore the "along for the ride" data in our examples
  - Remember no data at non-leaves

Find



- This is a new kind of tree
  - We are used to data at internal nodes
- But **`find`** is still an easy root-to-leaf recursive algorithm
  - At each internal node do binary search on the ≤ M-1 keys
  - At the leaf do binary search on the ≤ L data items
- But to get logarithmic running time, we need a balance condition…

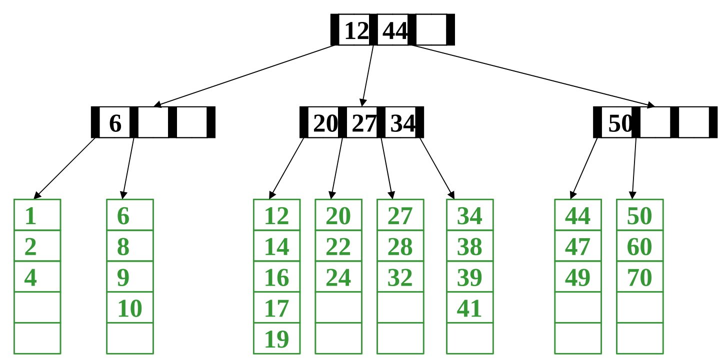
Structure Properties

- Root (special case)
  - If tree has ≤ *L* items, root is a leaf (very strange case)
  - Else has between 2 and *M* children
- Internal nodes
  - Have between ⌈*M*/2⌉ and *M* children, i.e., at least half full
- Leaf nodes
  - All leaves at the same depth
  - Have between ⌈*L*/2⌉ and *L* data items, i.e., at least half full

(Any *M* > 2 and *L* will work; picked based on disk-block size)

Example

Suppose *M*=4 (max children) and *L*=5 (max items at leaf)
- All internal nodes have at least 2 children
- All leaves have at least 3 data items (only showing keys)
- All leaves at same depth



Balanced enough

Not hard to show height *h* is logarithmic in number of data items *n*

- Let $M > 2$ (if $M = 2$, then a list tree is legal – no good!)
- Because all nodes are at least half full (except root may have only 2 children) and all leaves are at the same level, the minimum number of data items *n* for a height $h>0$ tree is…

$$n \geq \underbrace{2 \lceil M/2 \rceil^{h-1}}_{\text{minimum number of leaves}} \underbrace{\lceil L/2 \rceil}_{\text{minimum data per leaf}}$$

Exponential in height because $\lceil M/2 \rceil > 1$

B-Tree vs. AVL Tree

Suppose we have 100,000,000 items

- Maximum height of AVL tree?
- Maximum height of B tree with M=128 and L=64?

B-Tree vs. AVL Tree

Suppose we have 100,000,000 items

- Maximum height of AVL tree?
  - Recall S(h) = 1 + S(h-1) + S(h-2)
  - lecture7.xlsx reports: 47
- Maximum height of B tree with M=128 and L=64?
  - Recall $(2 + \lceil M/2 \rceil^{h-1}) \lceil L/2 \rceil$
  - lecture9.xlsx reports: 5 (and 4 is more likely)
  - Also not difficult to compute via algebra

Disk Friendliness

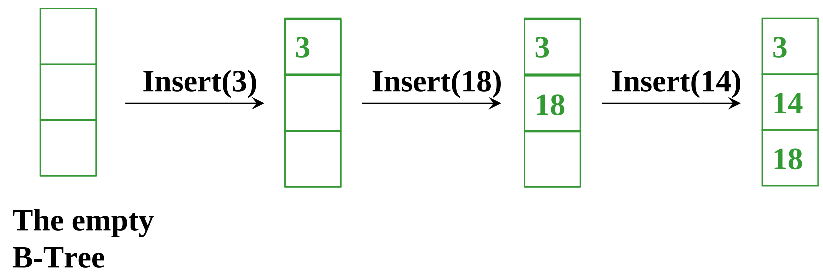
What makes B trees so disk friendly?

- Many keys stored in one node
  - All brought into memory in one disk access
  - Pick *M* wisely. Example: block=1KB, then *M*=128
  - Makes the binary search over *M*-1 keys totally worth it
- Internal nodes contain only keys
  - Any `find` wants only one data item
  - So only bring one leaf of data items into memory
  - Data-item size doesn't affect what *M* is

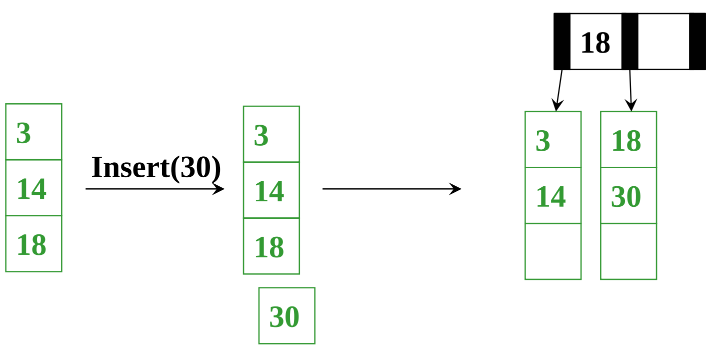
Maintaining balance

- So this seems like a great data structure (and it is)
- But we haven't implemented the other dictionary operations yet
  - **`insert`**
  - **`delete`**
- As with AVL trees, the hard part is maintaining structure properties
  - Example: for **`insert`**, there might not be room at the correct leaf

Building a B-Tree (insertions)



***M* = 3 *L* = 3**



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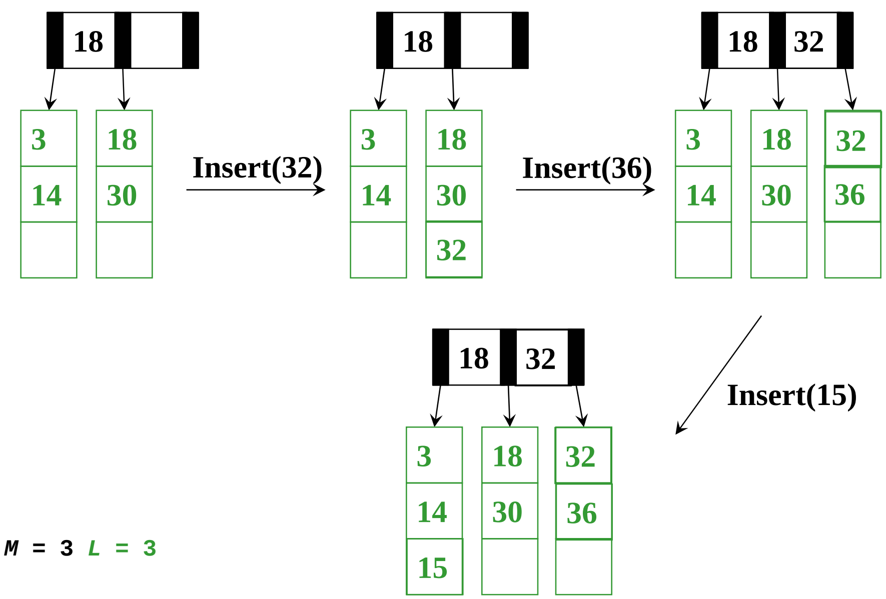

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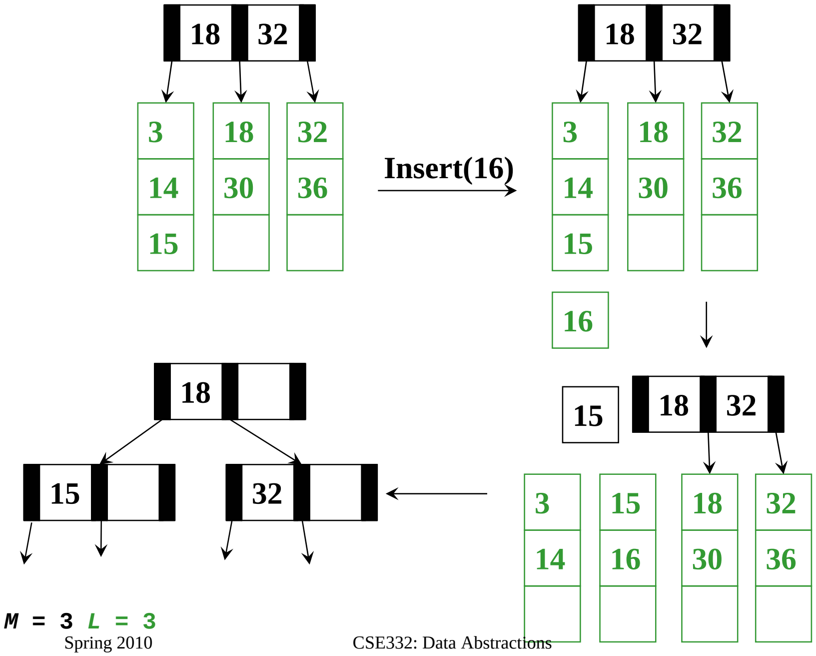

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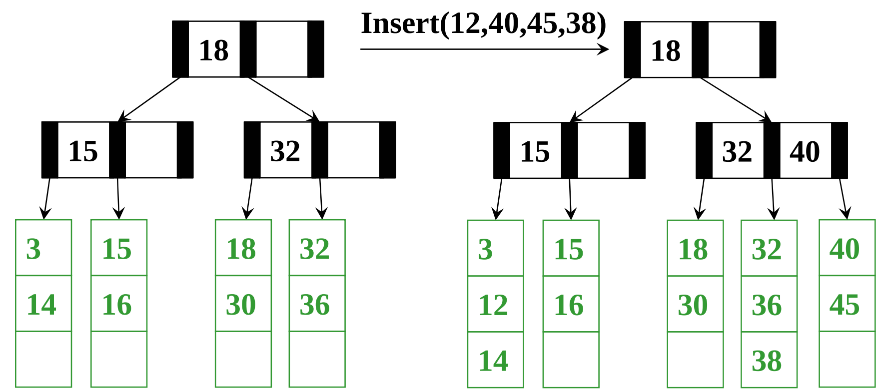

M = 3 *L* = 3

Insertion Algorithm

1. Insert the data in its leaf in sorted order

2. If the leaf now has *L*+1 items, *overflow!*
   - Split the leaf into two nodes:
     - Original leaf with $\lceil (L+1)/2 \rceil$ smaller items
     - New leaf with $\lfloor (L+1)/2 \rfloor = \lceil L/2 \rceil$ larger items
   - Attach the new child to the parent
     - Adding new key to parent in sorted order

3. If step (2) caused the parent to have *M*+1 children, *overflow!*
   - …

Insertion algorithm continued

3. If an internal node has *M*+1 children
   - Split the node into two nodes
     - Original node with $\lceil (M+1)/2 \rceil$ smaller items
     - New node with $\lfloor (M+1)/2 \rfloor = \lceil M/2 \rceil$ larger items
   - Attach the new child to the parent
     - Adding new key to parent in sorted order

Splitting at a node (step 3) could make the parent overflow too
- So repeat step 3 up the tree until a node doesn't overflow
- If the root overflows, make a new root with two children
  - This is the only case that increases the tree height

Efficiency of insert

- Find correct leaf: $O(\log_2 M \log_M n)$
- Insert in leaf: $O(L)$
- Split leaf: $O(L)$
- Split parents all the way up to root: $O(M \log_M n)$

Total: $O(L + M \log_M n)$

But it's not that bad:

- Splits are not that common (have to fill up nodes)
- Remember disk accesses were the name of the game: $O(\log_M n)$