CSE332: Data Abstractions

Lecture 2: Math Review; Algorithm Analysis

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Announcements

Project 1 posted
- Section materials on using Eclipse will be very useful if you have never used it
- (Could also start in a different environment if necessary)
- Section material on generics will be very useful for Phase B

Homework 1 posted

Feedback on typos is welcome
- Won't announce every time I update posted materials with minor fixes

Today

- Finish discussing queues
- Review math essential to algorithm analysis
  - Proof by induction
  - Powers of 2
  - Exponents and logarithms
- Begin analyzing algorithms
  - Using asymptotic analysis (continue next time)

Mathematical induction

Suppose $P(n)$ is some predicate (mentioning integer n)
- Example: $n \geq n/2 + 1$

To prove $P(n)$ for all integers $n \geq c$, it suffices to prove
1. $P(c)$ – called the "basis" or "base case"
2. If $P(k)$ then $P(k+1)$ – called the "induction step" or "inductive case"

Why we will care:

To show an algorithm is correct or has a certain running time *no matter how big a data structure or input value is*

(Our "n" will be the data structure or input size.)

Example

P(n) = "the sum of the first n powers of 2 (starting at 0) is 2^n-1"

Theorem: *P(n)* holds for all $n \geq 1$
Proof: By induction on n

- Base case: n=1. Sum of first 1 power of 2 is 2^0 , which equals 1.
  And for n=1, 2^n-1 equals 1.
- Inductive case:
  - Assume the sum of the first k powers of 2 is 2^k-1
  - Show the sum of the first (k+1) powers of 2 is 2^{k+1}-1

  Using assumption, sum of the first (k+1) powers of 2 is
  $(2^k-1) + 2^{(k+1)-1} = (2^k-1) + 2^k = 2^{k+1}-1$

Powers of 2

- A bit is 0 or 1
- A sequence of n bits can represent 2^n distinct things
  - For example, the numbers 0 through 2^n-1
- 2^{10} is 1024 ("about a thousand", kilo in CSE speak)
- 2^{20} is "about a million", mega in CSE speak
- 2^{30} is "about a billion", giga in CSE speak

Java: an **`int`** is 32 bits and signed, so "max int" is "about 2 billion"
a **`long`** is 64 bits and signed, so "max long" is 2^{63}-1

Therefore…

Could give a unique id to…

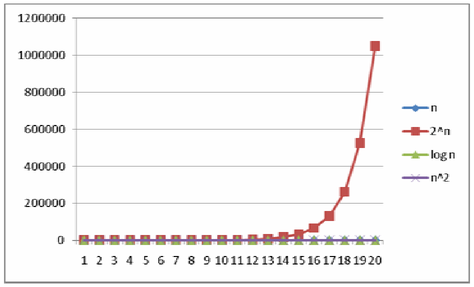
- Every person in the U.S. with 29 bits
- Every person in the world with 33 bits
- Every person to have ever lived with 38 bits (estimate)
- Every atom in the universe with 250-300 bits

So if a password is 128 bits long and randomly generated, do you think you could guess it?

Logarithms and Exponents

- Since so much is binary in CS **`log`** almost always means $\mathbf{log_2}$
- Definition: $\mathbf{log_2\ x = y}$ if $\mathbf{x = 2^y}$
- So, $\mathbf{log_2}$ 1,000,000 = "a little under 20"
- Just as exponents grow *very* quickly, logarithms grow *very* slowly

See Excel file for plot data – play with it!



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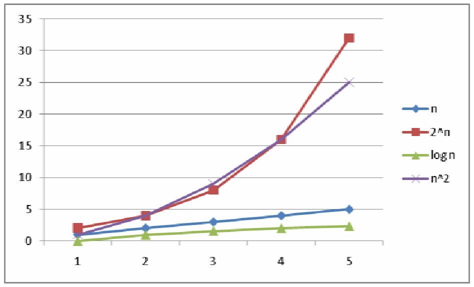

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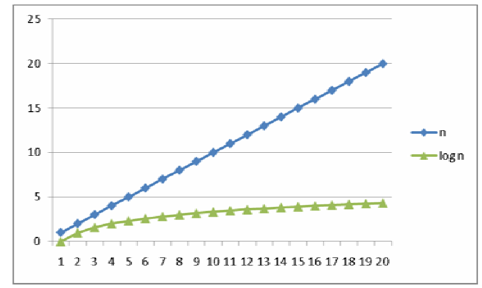

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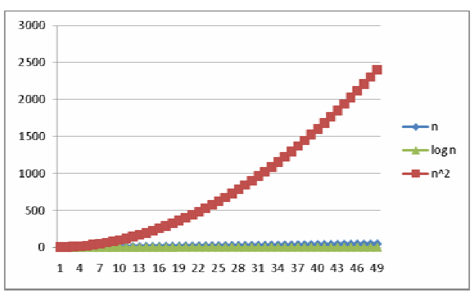


Properties of logarithms

- $\log(A*B) = \log A + \log B$
  - So $\log(N^k) = k \log N$
- $\log(A/B) = \log A - \log B$
- $\log(\log x)$ is written $\log \log x$
  - Grows as slowly as 2^{2^y} grows fast
- $(\log x)(\log x)$ is written $\log^2 x$
  - It is greater than $\log x$ for all $x > 2$

Log base doesn't matter much!

"Any base *B* log is equivalent to base 2 log within a constant factor"

- And we are about to stop worrying about constant factors!
- In particular, $\log_2 x = 3.22 \log_{10} x$
- In general,

  $\log_B x = (\log_A x) / (\log_A B)$

Algorithm Analysis

As the "size" of an algorithm's input grows
(integer, length of array, size of queue, etc.):

- How much longer does the algorithm take (time)
- How much more memory does the algorithm need (space)

Because the curves we saw are so different, we often only care about "which curve we are like"

Separate issue: Algorithm *correctness* – does it produce the right answer for all inputs

- Usually more important, naturally

Example

- What does this pseudocode return?

```
x := 0;
for i=1 to N do
  for j=1 to i do
     x := x + 3;
return x;
```

- Correctness: For any N ≥ 0, it returns…

Example

- What does this pseudocode return?

```
x := 0;
for i=1 to N do
  for j=1 to i do
     x := x + 3;
return x;
```

- Correctness: For any N ≥ 0, it returns 3N(N+1)/2
- Proof: By induction on *n*
  - *P(n)* = after outer for-loop executes *n* times, **x** holds 3n(n+1)/2
  - Base: n=0, returns 0
  - Inductive: From *P(k)*, **x** holds 3k(k+1)/2 after *k* iterations. Next iteration adds 3(k+1), for total of 3k(k+1)/2 + 3(k+1) = (3k(k+1) + 6(k+1))/2 = (k+1)(3k+6)/2 = 3(k+1)(k+2)/2

Example

- How long does this pseudocode run?

```
x := 0;
for i=1 to N do
  for j=1 to i do
     x := x + 3;
return x;
```

- Running time: For any N ≥ 0,
  - Assignments, additions, returns take "1 unit time"
  - Loops take the sum of the time for their iterations
- So: 2 + 2*(number of times inner loop runs)
  - And how many times is that…

Example

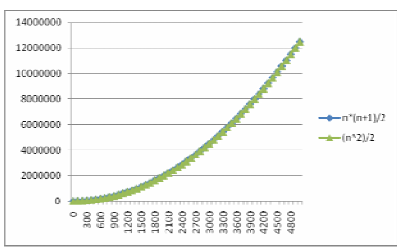
- How long does this pseudocode run?

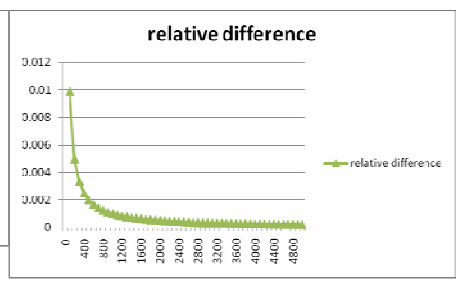
```
x := 0;
for i=1 to N do
  for j=1 to i do
     x := x + 3;
return x;
```

- The total number of loop iterations is N*(N+1)/2
  - This is a very common loop structure, worth memorizing
  - Proof is by induction on N, known for centuries
  - This is *proportional to* N^2 , and we say $O(N^2)$, "big-Oh of"
    - For large enough N, the N and constant terms are irrelevant, as are the first assignment and return
    - See plot… N*(N+1)/2 vs. just $N^2/2$

Lower-order terms don't matter

N*(N+1)/2 vs. just $N^2/2$





Recurrence Equations

- For running time, what the loops did was irrelevant, it was how many times they executed.
- Running time as a function of input size *n* (here loop bound):

  $T(n) = n + T(n-1)$

  (and $T(0)$ = 2ish, but usually implicit that $T(0)$ is some constant)
- Any algorithm with running time described by this formula is $O(n^2)$
- "Big-Oh" notation also ignores the constant factor on the high-order term, so $3N^2$ and $17N^2$ and $(1/1000)\ N^2$ are all $O(N^2)$
  - As N grows large enough, no smaller term matters
  - Next time: Many more examples + formal definitions

Big-O: Common Names

$O(1)$	constant (same as $O(k)$ for constant k)
$O(\log n)$	logarithmic
$O(n)$	linear
$O(n \log n)$	"$n \log n$"
$O(n^2)$	quadratic
$O(n^3)$	cubic
$O(n^k)$	polynomial (where is k is an constant)
$O(k^n)$	exponential (where k is any constant > 1)

Pet peeve: "exponential" does not mean "grows really fast", it means "grows at rate proportional to k^n for some k>1"

- A savings account accrues interest exponentially (k=1.01?)
- If you don't know k, you probably don't know it's exponential