

# CSE 331

## Type Polymorphism

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# Type Polymorphism

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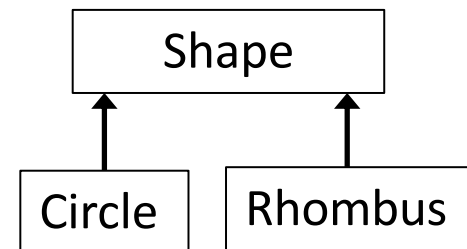
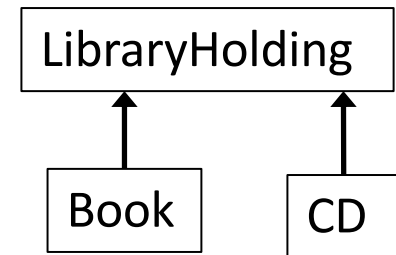
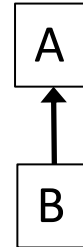
- Last design topic will be **"type polymorphism"**
  - allows code to work correctly with more than one type
- We will look at two instances of this
  - **subtypes**: can be used in where supertype expected
  - **generics**: can use type with different instantiation of its type parameters

# Subtypes

# What Is a Subtype?

---

- Sometimes "every B is an A"
  - every book is a library holding
  - every circle is a shape
- Denote with an upward arrow
- Not always so clear
  - would like a **formal** definition



## Recall: Comparing Specifications

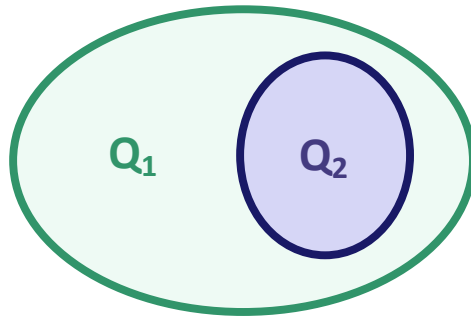
---

- Would like a similar definition for specifications
- Specification  $S_1$  is **stronger** than  $S_2$ ...
  - whenever  $S_1$  is satisfied,  $S_2$  is also satisfied
  - i.e., satisfying  $S_1$  implies satisfying  $S_2$
- Code written for  $S_2$  also works with  $S_1$
- But what does this mean?
  - specifications have a **precondition** and **postcondition**

# Recall: Stronger Assertions

---

- **Assertion** is **stronger** iff it holds in a subset of states

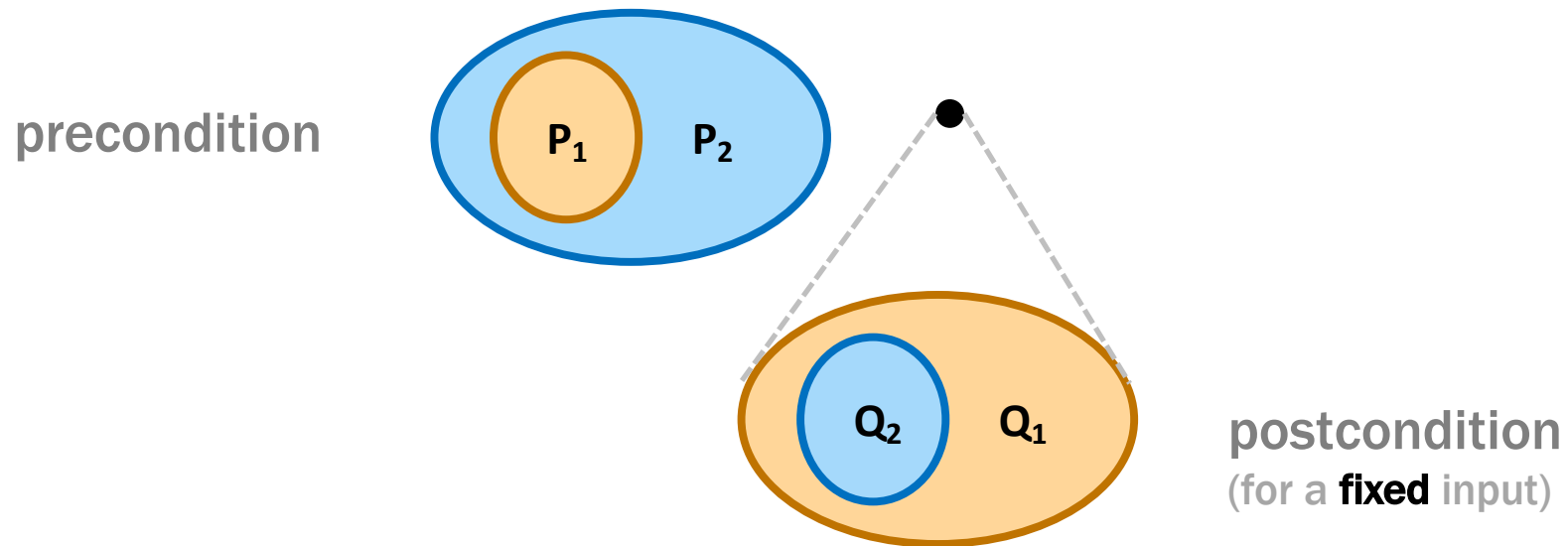


- **Stronger** assertion implies the **weaker** one
  - stronger is a synonym for “implies”
  - weaker is a synonym for “is implied by”

# Recall: Comparing Function Specifications

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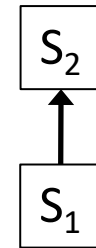
- Specification  $S_1$  is **stronger** than  $S_2$  if it has...
  - a **stronger** postcondition and the same precondition
  - a **weaker** precondition and the same postcondition
  - (or both)



# What Is a Subtype?

---

- Would like a similar definition for types
- ADT  $S_1$  is a **subtype** (stronger) than  $S_2$ ...
  - when is  $S_1$  satisfied,  $S_2$  is also satisfied
  - i.e., satisfying  $S_1$  implies satisfying  $S_2$
- Code written for  $S_2$  also works with  $S_1$
- Alternatively, subtypes are "**substitutable**"
  - called the "Liskov substitution principle?"

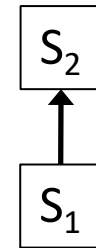




# What Is a Subtype?

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  - when is  $S_1$  satisfied,  $S_2$  is also satisfied
  - i.e., satisfying  $S_1$  implies satisfying  $S_2$
- Code written for  $S_2$  also works with  $S_1$
- But what does this mean?



# Comparing ADT Specifications

---

- ADT  $S_1$  is a **subtype** (stronger) than  $S_2$  if...
  - $S_1$  has all the methods of  $S_2$
  - each of method of  $S_1$  has a stronger specification than the corresponding method of  $S_2$   
stronger or the same specification

- For example:

```
interface A {  
    int foo(String s);  
    Object bar();  
}
```

```
interface B {  
    int foo(String s);  
    String bar();  
    int baz();  
}
```

foo exists with same spec  
bar exists with stronger spec

okay to have an extra method

B is a subtype but Java will not let you substitute

# Comparing ADT Specifications

---

- ADT  $S_1$  is a **subtype** (stronger) than  $S_2$  if...
  - $S_1$  has all the methods of  $S_2$
  - each of method of  $S_1$  has a stronger specification than the corresponding method of  $S_2$   
stronger or the same specification

- For example:

Java requires B to extend A

```
interface A {  
    int foo(String s);  
    Object bar();  
}
```

```
interface B extends A{  
    int foo(String s);  
    String bar();  
    int baz();  
}
```

Java mistakenly equates subtypes and subclasses

# Recall: Subclasses

---

- Subclassing is a means of sharing code
  - subclass gets parent fields & methods (unless overridden)

```
class Product {  
    private String name;  
    private int price;  
    public String getName() {return name; }  
    public int getPrice() { return price; }  
}  
  
class SaleProduct extends Product {  
    private float discount;  
    public int getPrice() {  
        return (1 - discount) * super.getPrice();  
    }  
}
```

# Subclassing ≠ Subtyping

---

- Subclassing is a means of code sharing
  - all fields of the superclass
  - all methods of superclass copied into subclass
    - unless overridden in the subclass
  - ensures it has all the methods of the parent class
- Subtyping is about specifications
  - each method specification must be stronger
    - weaker **precondition** and/or stronger **postcondition**
  - mostly in the **comments!**
    - which the Java compiler does not read

# Example 1

---

```
// An integer value that represents ...
```

```
interface NumberA {
```

```
    // @requires obj is present in A
```

```
    // @returns an index i such that A[i] = obj
```

```
    int indexOf(int[] A);
```

```
}
```

```
interface NumberB extends NumberA {
```

```
    // @requires obj is present in A
```

```
    // @returns the smallest index i such that A[i] = obj
```

```
    int indexOf(int[] A);
```

```
}
```

Would Java notice if we swapped these?

No! Compiler doesn't read the comments.

– can see that `NumberB` is a subtype of `NumberA`

`NumberB` has a stronger postcondition than `NumberA`

## Example 2

---

```
// An integer value that represents ...
interface NumberA {
    // @returns the smallest i such that A[i] = obj
    //      or -1 if obj is not present in A
    int indexOf(int[] A);
}

interface NumberB extends NumberA {
    // @requires obj is present in A
    // @returns the smallest index i such that A[i] = obj
    int indexOf(int[] A);
}
```

- can see that `NumberB` is not a subtype of `NumberA`  
    `NumberB` allows *fewer* inputs (stronger precondition) than `NumberA`
- but Java allows it anyway

# Subtyping in the Type Checker

---

```
class NumberA { ... }  
class NumberB extends NumberA { ... }  
  
public void foo(NumberA n) { ... }  
  
NumberB m = ...  
foo(m);           // Java allows this call
```

- Java allows you to pass the subclass
  - it allows substitution of subclasses
  - it **assumes** that subclasses are subtypes
- Subtyping shows up when you make method calls
  - both the arguments passed in and the return value



# Recall: Subclasses

---

- Subclassing is a surprisingly dangerous feature
- Subclassing tends to break modularity
  - creates **tight coupling** between super- and sub-class
  - often see the “fragile base class” problem
    - changes to super class often break subclasses
- New: Java **assumes** subclasses are subtypes
  - Java will let you pass subclass where supertype expected
  - no way for it to check that it is really a subtype!
  - code will break in strange ways if it's not true

# Subtyping in the Type Checker

---

```
class NumberA { ... }  
class NumberB extends NumberA { ... }  
  
    public void foo(NumberA n) {  
        ... in here ...  
    }
```

- In the body of the method `foo`, variable `n` will be instance of `NumberA` or a subclass (e.g. `NumberB`)
- It will have all the methods of `NumberA`
  - any subclass gets those methods copied into it
  - rules out many bugs!

# Java Does Some Checks

---

- Java checks the return types:

```
interface A {  
    String foo();  
}
```

```
interface B extends A {  
    Object foo(); // error!  
}
```

- subclass wants to return non-String values
- checks the part of the **postcondition** visible in the **types**

# Java Does Some Checks

---

- Java checks most exceptions:

```
interface A {  
    String foo() throws IOException;  
}
```

```
interface B extends A {  
    String foo() throws Exception; // error!  
}
```

- subclass wants to throw non-IOException exceptions
- checks the part of the **postcondition** visible in the **types**
- only applies to "checked" exceptions  
does not check RuntimeException or Error

# Java Cannot Handle More Inputs

---

- Java checks most exceptions:

```
interface A {  
    int foo(String s);  
}
```

Java (and C++) identify methods by  
their signature (name + argument types)

```
interface B extends A {  
    int foo(Object s); // doesn't work!  
}
```

- this is a strengthening (more allowed inputs)  
but it will not work properly...
- this is overloading not overriding  
B has two methods named "foo", not one

# TypeScript Can Handle More Inputs

---

- This works properly in TypeScript

```
interface A {  
    foo(s: string): number;  
}
```

```
interface B extends A {  
    foo(s: string | number): number; // okay!  
}
```

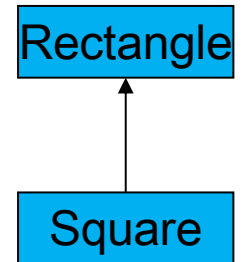
- TypeScript has only one method with a given name
- also not necessary to say "extends"

TypeScript has a structural, not nominal, type system

# Squares and Rectangles

---

- Is a Square a Rectangle?
  - yes
- Is a Square a subtype of Rectangle?
  - seems like it should be...



# Squares and Rectangles

---

- Consider the following method of Rectangle:

```
// @modifies obj
// @effects obj.width = w and obj.height = h
void setSize(int w, int h)
```

- How do we implement this in Square?

```
// @requires w = h
// @modifies obj
// @effects obj.width = w and obj.height = h
void setSize(int w, int h)
```

This is a weakening (stronger precondition)!



# Squares and Rectangles

---

- Consider the following method of Rectangle:

```
// @modifies obj
// @effects obj.width = w and obj.height = h
void setSize(int w, int h)
```

- How do we implement this in Square?

```
// @modifies obj
// @effects obj.width = w and obj.height = h
// @throws BadSize if w != h
void setSize(int w, int h)
```

This is an incomparable spec

# Squares and Rectangles

---

- Consider the following method of Rectangle:

```
// @modifies obj
// @effects obj.width = w and obj.height = h
void setSize(int w, int h)
```

- How do we implement this in Square?

```
// @modifies obj
// @effects obj.width = w and obj.height = w
void setSize(int w, int h)
```

This is an incomparable spec

# Squares and Rectangles

---

- Consider the following method of Rectangle:

```
// @modifies obj
// @effects obj.width = w and obj.height = h
void setSize(int w, int h)
```

- How do we implement this in Square?

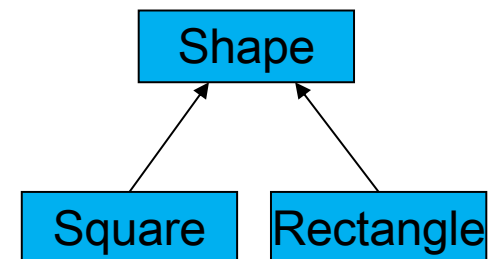
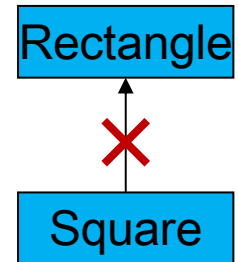
```
// @modifies obj
// @effects obj.width = obj.height = sideLength
void setSize(int sideLength)
```

This isn't the same method  
(overloading not overriding)

# Squares and Rectangles

---

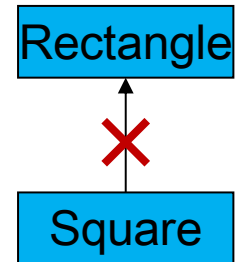
- **Square is not a subtype of Rectangle**
  - subtyping can be *unintuitive*
- **Solution 1: make them siblings**
  - common parts in the parent **class** Shape
  - **cannot substitute** Square **for** Rectangle



# Squares and Rectangles

---

- **Square is not a subtype of Rectangle**
  - subtyping can be *unintuitive*
- **Solution 2: make them immutable**
  - problem in `setSize` because it mutates
  - would be no problems if we did not allow it
- **Will see more examples of this later on...**
  - reading and writing operations are different



# Benefits of Immutability

---

- **No worries about representation exposure**
  - do not need to copy in & copy out
- **No worries about key mutation errors**
  - one of the worst bugs out there
- **Subtyping usually works the way you expect**
  - squares are subtypes of rectangles

# Inappropriate Subtyping in the JDK

---

```
class Hashtable {  
    public void put(Object key, Object value) { ... }  
    public Object get(Object key) { ... }  
}  
  
class Properties extends Hashtable {  
    public void setProperty(String key, String val) {  
        put(key, val);  
    }  
  
    public String getProperty(String key) {  
        return (String) get(key);  
    }  
}
```

- can cast `Properties` to `Hashtable` (but it's not a good idea!)

# Inappropriate Subtyping in the JDK

---

```
class Hashtable {  
    public void put(Object key, Object value) { ... }  
    public Object get(Object key) { ... }  
}  
  
class Properties extends Hashtable {  
    public void setProperty(String key, String val) {  
        put(key, val);  
    }  
  
    public String getProperty(Object key) {  
        return (String) get(key);  
    }  
}
```

```
Properties p = new Properties();  
Hashtable h = p;  
h.put("One", 1);  
p.getProperty("One"); // crash!
```



# Inappropriate Subtyping in the JDK

---

- The documentation says not to do this:

*"Because Properties inherits from Hashtable, the put and putAll methods can be applied to a Properties object. ... If the store or save method is called on a "compromised" Properties object that contains a non-String key or value, **the call will fail**."*

- Problem solved?
  - no! someone will still mess this up
  - this is **bad design**
- What should they do?

# Composition

---

```
class Properties {  
    private Hashtable tbl;  
  
    public void setProperty(String key, String val) {  
        tbl.put(key, val);  
    }  
  
    public String getProperty(Object key) {  
        return (String) tbl.get(key);  
    }  
}
```

- Can no longer be misused
  - has no `get` or `put` methods at all
- Solution you already know is called "composition"

# Generic Types

# The Many Varieties of Lists

---

**type** List := nil | cons( $x : \mathbb{Z}$ , L : List)

**list of integers**

**type** BList := nil | cons( $f : \mathbb{B}$ , L : List)

**list of booleans**

**type** PList := nil | cons( $p : \mathbb{Z} \times \mathbb{Z}$ , L : List)

**list of pairs**

— **common pattern to all of these:**

all the same except what **data** we store in them

# The Many Varieties of Lists

---

**type** List := nil | cons( $x : \mathbb{Z}$ , L : List)

list of integers

**type** BList := nil | cons( $f : \mathbb{B}$ , L : List)

list of booleans

**type** PList := nil | cons( $p : \mathbb{Z} \times \mathbb{Z}$ , L : List)

list of pairs

- Can have one definition for all lists:

**type** List $\langle A \rangle$  := nil | cons( $x : A$ , L : List)

- those above are List $\langle \mathbb{Z} \rangle$ , List $\langle \mathbb{B} \rangle$ , and List $\langle \mathbb{Z} \times \mathbb{Z} \rangle$
- "A" is a type argument
- "List" is a function from types to types

# Java Before Generics

---

- **Collections allowed any kind of object**

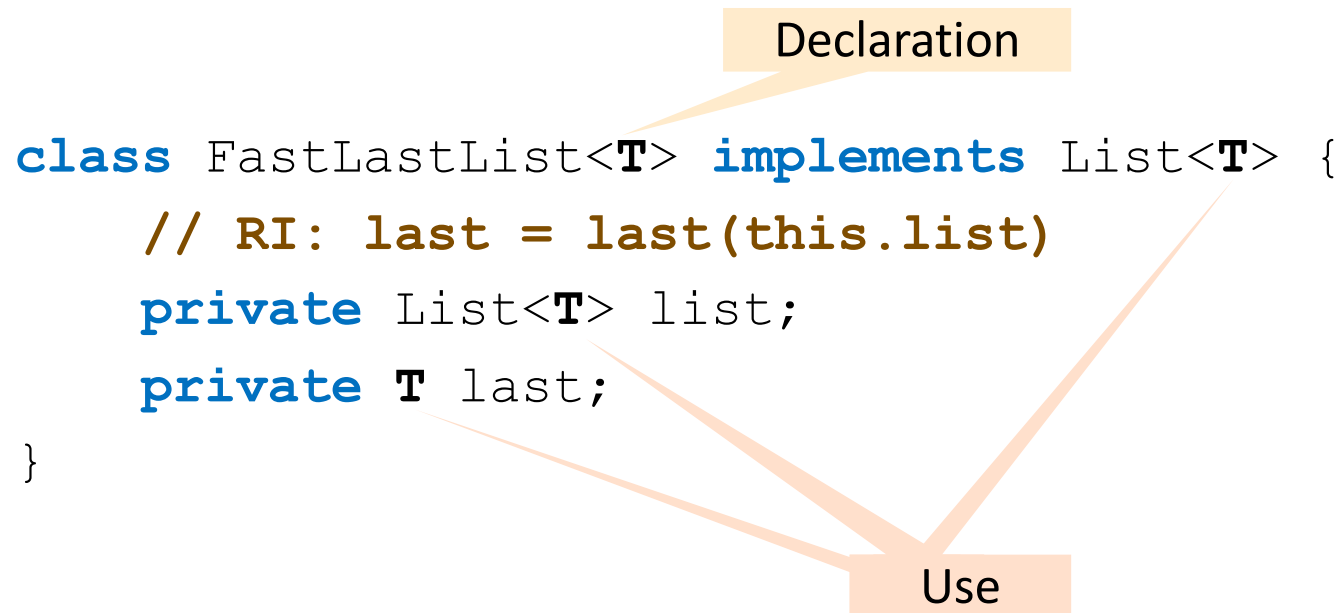
```
Hashtable h = new Hashtable();  
h.put("abc", new Integer(3));
```

```
Integer v = (Integer) h.get("abc");
```

- **only type checking occurs at run-time**  
relying on your unit tests & code reviews to catch everything
- **programmers frequently make mistakes here**

# Java Generics

---



- "T" is an argument to the type declaration
  - also called a "type parameter"

# Java Generics

---

```
class Name<A, B, C, D> { ... }
```

```
interface Name<A, B, C, D> { ... }
```

- **Declarations can have any number of arguments**

- **Java style is to use short names**

- e.g. "E" for element, "K" for key, "V" for value

- **Must fill in the arguments to use it**

```
Name<Integer, Boolean, String, Object> n;
```

- **Java only *warns* if you leave off < . . > but don't do it!**



# Type Constraints

---

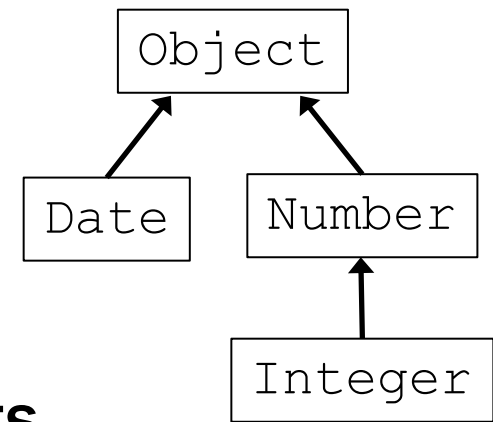
- Type declarations constrain the values passed in

```
boolean add1 (Object elt);
```

```
boolean add2 (Number elt);
```

```
add1 (new Date()); // okay
```

```
add2 (new Date()); // error!
```



- Same idea applies to type arguments
  - here they are called "bounds" on the type

# Type Bounds

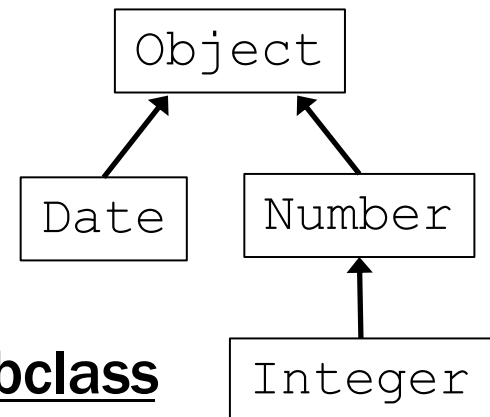
---

- Type Bounds constrain the types passed in

```
interface List1<E extends Object> { ... }  
interface List2<E extends Number> { ... }
```

```
List1<Date> L1;    // okay  
List2<Date> L2;    // error!
```

- these are called "upper bounds"
- type argument can be that type or a subclass



# Java Generics (Take Two)

---

```
class Name<T1 extends B1, ..., Tn extends Bn>  
interface Name<T1 extends B1, ..., Tn extends Bn>
```

- **Declarations can have any number of arguments**
  - each argument has an optional upper bound  
if not provided, it is "Object"
- **Must fill in the arguments to use it**

```
Name<Integer, Boolean, String, Object> n;
```

- Java only **warns** if you leave off <..>
- Java gives an **error** if the type does not meet the bound

# Uses of Type Arguments

---

- Code can only use methods from the type bound

```
class Foo1<E extends Object> {  
    public int m(E arg) {  
        return arg.intValue(); // error!  
    }  
}  
  
class Foo2<E extends Number> {  
    public int m(E arg) {  
        return arg.intValue(); // okay!  
    }  
}
```

- can only call methods guaranteed to be there

# More Examples of Generic Classes

---

- These can look pretty crazy at first:

```
class Graph<N> implements Iterable<N> {  
    private Map<N, Set<N>> node2neighbors;  
    public Graph(Set<N> nodes, Set<Pair<N, N>> edges) {  
        ...  
    }  
}
```

```
interface Path<N, P> extends Path<N, P>>  
    extends Iterable<N>, Comparable<Path<N, P>> {  
    public Iterator<N> iterator();  
}
```

# More Examples of Generic Classes

---

- Type argument is in scope immediately after ","
- Often see variable used in its bound:

```
class TreeSet<T extends Comparable<T>>

interface Comparable<C> {
    int compareTo(C other);
}
```

- so `Comparable<T>` will have:

```
int compareTo(T other); // compare two Ts
```

# More Bounds

---

```
class Name<A extends B>
```

- "B" is an upper bound

```
class Name<A extends B & C & D>
```

- can include multiple upper bounds with "&"
- these can be classes or interfaces

# **Generic Methods**



## Example 3

---

```
class Utils {  
    static double sum(List<Number> list) {  
        double result = 0;  
        for (Number n : list) {  
            result += n.doubleValue();  
        }  
        return result;  
    }  
    ...  
}
```

- would like `sum` to work with any type of number  
e.g., want to pass `List<Double>` or `List<Integer>`
- that will not work for reasons we will see later

## Example 4

---

```
class Utils {  
    ...  
    static Object choose(List<Object> list) {  
        int index = (int) (list.size() * Math.random());  
        return list.get(index);  
    }  
}
```

- **would like** choose **to work with any element type**  
e.g., want to pass `List<Double>` or `List<String>`
- **would like** `choose(List<Double>)` **to return** `Double`

# Generic Methods in Java

---

```
class Utils {  
    static <T extends Number> double sum(List<T> list) {  
        double result = 0;  
        for (T n : list) {  
            result += n.doubleValue();  
        }  
        return result;  
    }  
  
    static <T> T choose(List<T> list) {  
        int index = (int) (list.size() * Math.random());  
        return list.get(index);  
    }  
}
```

The diagram illustrates the Declaration and Use of generic methods in Java. It features two orange callout boxes labeled "Declaration" and "Use".

- The first "Declaration" box points to the `static <T extends Number> double` part of the `sum` method signature.
- The first "Use" box points to the `double` return type in the `sum` method signature.
- The second "Declaration" box points to the `static <T> T` part of the `choose` method signature.
- The second "Use" box points to the `int` return type in the `choose` method signature.

## Example 4 (Updated)

---

```
class Utils {  
    ...  
  
    static <T> T choose(List<T> list) {  
        int index = (int) (list.size() * Math.random());  
        return list.get(index);  
    }  
}
```

- can now call like this:

```
List<String> list = ...;  
String s = choose(list); // result is a String
```

- no need to fill in type parameters to the method  
they are always "inferred"

## Example 3 (Updated)

---

```
class Utils {  
    static <T extends Number> double sum(List<T> list) {  
        double result = 0;  
        for (T n : list) {  
            result += n.doubleValue();  
        }  
        return result;  
    }  
    ...  
}
```

– since T **extends** Number, we can call doubleValue

# Generic Classes and Methods

---

```
class Name<A> {  
    public <B> int foo(A a, B b) {  
        ...  
    }  
}
```

- both the class and the method have type params
- all of those are in scope within the method

# More Examples of Generic Methods

---

```
<T extends Comparable<T>> T max(Collection<T> c) {  
    ...  
}
```

```
<T extends Comparable<T>> void sort(List<T> list) {  
    ... // can use List.get() and T.compareTo(T)  
}
```

```
<T> void copyTo(List<T> dest, List<T> src) {  
    for (T t : src)  
        dest.add(t);  
}
```

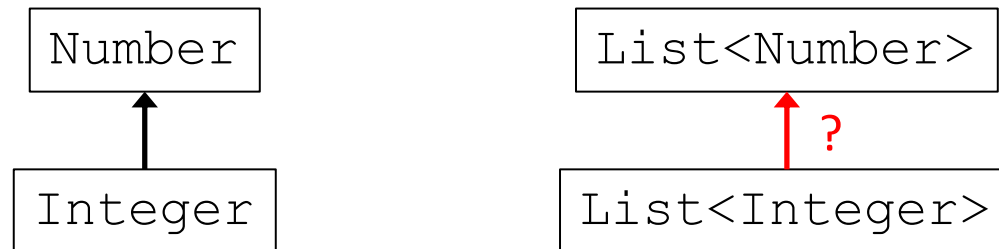
– last method can be improved further...

# Generics & Subtyping



# How Does Subtyping Work With Generics?

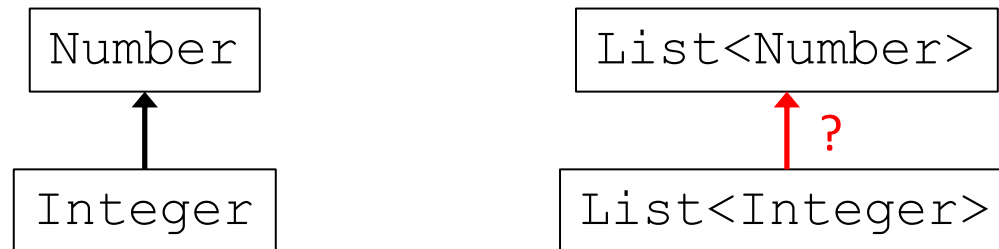
---



- Integer **is a subtype of** Number
  - can safely be substituted where Number is expected  
(note that both classes are immutable...)
- **Is** List<Integer> **a subtype of** List<Number>?
  - can we safely substitute the latter for the former?

# How Does Subtyping Work With Generics?

---



- If true, `List<T>` would be **covariant** in its type arg ↑
- Would be **contravariant** if arrow was reversed ↓

# List Is Neither Covariant Nor Contravariant

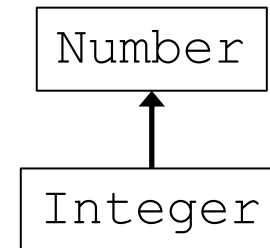
---

```
List<Number> numList = new ArrayList<Number>();  
List<Integer> intList = new ArrayList<Integer>();  
  
// cannot substitute numList for intList  
intList.add(new Integer(3));  
    becomes numList.add(new Integer(3));    // okay  
Integer n = intList.get(0);  
    becomes Integer n = numList.get(0);    // error!  
  
// cannot substitute intList for numList  
Number n = numList.get(0);  
    becomes Number n = intList.get(0);    // okay  
numList.add(new Double(3));  
    becomes intList.add(new Double(3));    // error!
```

# List Is Invariant

---

```
interface List<T> {  
    void add(T elem);  
    T get(int index);  
}
```



- becomes these two interfaces

```
interface List<Number> {  
    void add(Number elem);  
    Number get(int index);  
}
```

```
interface List<Integer> {  
    void add(Integer elem);  
    Integer get(int index);  
}
```

- add prevents covariance, get prevents contravariance
- List is "invariant" with respect to its type arg
- Java assumes all generic types are invariant

# Read-Only List is Covariant

---

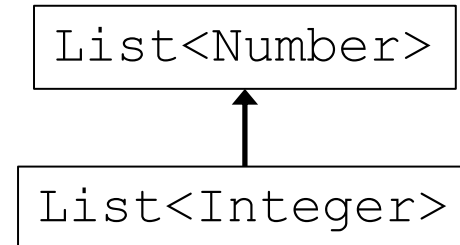
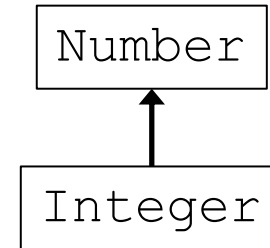
```
interface List<T> {  
    T get(int index);  
}
```

- becomes these two interfaces

```
interface List<Number> {  
    Number get(int index);  
}
```

```
interface List<Integer> {  
    Integer get(int index);  
}
```

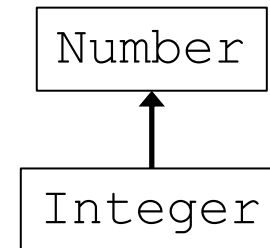
- this is **covariant**



# Write-Only List Is Contravariant

---

```
interface List<T> {  
    void add(T elem);  
}
```

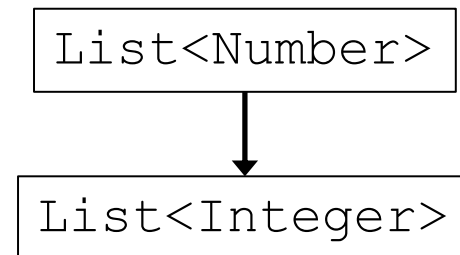


- becomes these two interfaces

```
interface List<Number> {  
    void add(Number elem);  
}
```

```
interface List<Integer> {  
    void add(Integer elem);  
}
```

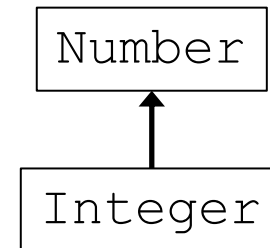
- this is **contravariant**



# List Is Invariant

---

```
interface List<T> {  
    void add(T elem);  
    T get(int index);  
}
```



- would be **covariant** if  $T$  is only a *return value*
  - would be **contravariant** if  $T$  is only an *argument*
- Java does not see these distinctions
    - all generic types are **invariant**
    - other languages do (e.g., C# and Scala)

# Java Generic Invariance

---

- **Cannot pass** `List<Integer>` **for** `List<Number>` **or vice versa!**
  - saw this before with `sum` **and** `choose`
  - generic methods are needed to overcome invariance

- Still have subtyping on classes themselves

```
interface A<T> { ... }  
interface B<T> extends A<T> { ... }
```

- can pass `B<Integer>` **where** `A<Integer>` **expected**
- cannot pass `B<Number>` **where** `A<Integer>` **expected**  
once a parameter changes, the classes are unrelated



## Example 5

---

```
interface Set<E> {  
    // @modifies obj  
    // @effects obj = c ++ obj_0  
    void addAll(_____ c)  
}
```

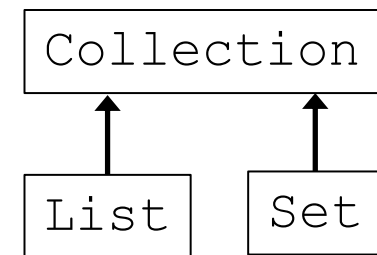
- what is best argument type?  
try to make this as general as possible...

## Example 5

---

```
interface Set<E> {  
    // @modifies obj  
    // @effects obj = c ++ obj_0  
    void addAll(Set<E> c)  
}
```

- how about this?
- it is too restrictive!  
cannot pass `List<E>` for example



# Example 5

---

```
interface Set<E> {  
    // @modifies obj  
    // @effects obj = c ++ obj_0  
    void addAll(Collection<E> c)  
}
```

- how about this?
- still too restrictive!

cannot pass `List<Integer>` to `addAll(Collection<Number>)`  
as we would have on `Set<Number>`

- prevented by **invariance**

## Example 5

---

```
interface Set<E> {  
    // @modifies obj  
    // @effects obj = c ++ obj_0  
    <T extends E> void addAll(Collection<T> c)  
}
```

- this is the most general

can pass `List<Integer>` to `addAll(Collection<T>)`

allowed on `Set<Number>` since `Integer` extends `Number`

- generic methods work around **invariance**

# Recall: More Examples of Generic Methods

---

```
<T extends Comparable<T>> T max(Collection<T> c) {  
    ...  
}
```

```
<T extends Comparable<T>> void sort(List<T> list) {  
    ... // can use List.get() and T.compareTo(T)  
}
```

```
<T> void copyTo(List<T> dest, List<T> src) {  
    for (T t : src)  
        dest.add(t);  
}
```

- last method can be improved further...

# More Examples (Updated)

---

```
<T, S extends T> void copyTo(List<T> dest, List<S> src) {  
    for (S t : src)  
        dest.add(t);  
}
```

- any valid  $S$  is castable to  $T$  since it is a superclass  
also works if  $S = T$

**Wildcards**

# Wildcard Example

---

- More concise way of writing some generics
  - this earlier example:

```
interface Set<E> {  
    <T extends E> void addAll(Collection<T> c);  
}
```

- can be written equivalently as:

```
interface Set<E> {  
    void addAll(Collection<? extends E> c);  
}
```

- wildcard is an anonymous type variable
    - automatically transformed into above with some name like "T"



# Wildcards

---

- **More concise way of writing some generics**
  - **"? extends E"** is an anonymous subclass of E  
or E itself
  - **"?"** is an anonymous subclass of Object  
or Object itself

## ? vs Object

---

- **Do not confuse** `List<?>` **with** `List<Object>`
  - former allows the latter
  - but also allows `List<Integer>`, `List<String>`, **etc.**
- **Cannot pass** `List<Integer>` **as** `List<Object>`
  - prevented by **invariance**
- **Can pass** `List<Integer>` **as** `List<?>`
  - allowed by generic methods

# Example

---

```
void foo(List<?> list1, List<?> list2) {  
    ...  
}
```

- each "?" is its own anonymous variable
- so this example becomes

```
<T1, T2> void foo(List<T1> list1, List<T2> list2) {  
    ...  
}
```

- if you want both to be the same type, you need this

```
<T> void foo(List<T> list1, List<T> list2) {  
    ...  
}
```

# Non-Examples

---

```
<T> T choose(Collection<T> c) {  
    ...  
}
```

- cannot be translated into a wildcard
- need both "T"s to be the same type  
type returned is whatever was in the list

- another non-example:

```
<T extends Comparable<T>> T max(Collection<T> c) {  
    ...  
}
```

# Wildcards

---

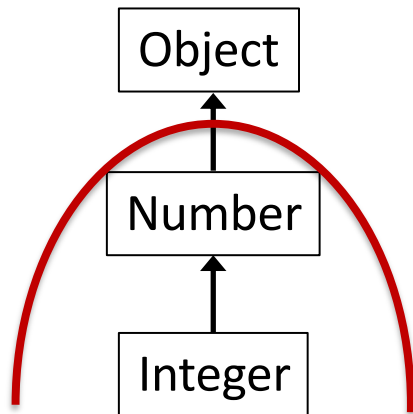
- More concise way of writing some generics
  - "**?** **extends** E" is an anonymous subclass of E  
or E itself
  - "**?**" is an anonymous subclass of `Object`  
or `Object` itself
  - "**?** **super** E" is an anonymous **superclass** of E  
or E itself
- No way to do this without wildcards!
  - no theoretical reason not to allow it  
presumably just cut due to lack of time

# Type Bounds

---

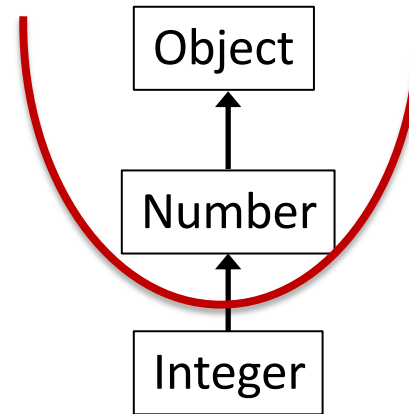
Upper Bound

? **extends** Number



Lower bound

? **super** Number



# Recall: More Examples (Updated)

---

```
<T, S extends T> void copyTo(List<T> dest, List<S> src) {  
    for (S t : src)  
        dest.add(t);  
}
```

- any valid  $S$  is castable to  $T$  since it is a superclass  
also works if  $S = T$

# More Examples (Updated More)

---

```
<T, S extends T> void copyTo(List<T> dest, List<S> src)
```

- can write this with some wildcards

```
<T> void copyTo(List<? super T> dest,  
               List<? extends T> src) {  
    for (T t : src)  
        dest.add(t);  
}
```

- still need one variable to connect the two types
- `dest` is anything that can accept (consume) "T"s
- `src` is anything that can give out (produce) "T"s



# Producer Extends, Consumer Super (PECS)

---

```
<T> void copyTo(List<? super T> dest,  
                List<? extends T> src) {  
    for (T t : src)  
        dest.add(t);  
}
```

- **Should you use "extends" or "super"?**
  - use "**? extends T**" when you get values (it's a **producer**)  
fine if it gives you a subclass [covariant case]
  - use "**? super T**" when you put values (it's a **consumer**)  
fine if it accepts a superclass [contravariant case]

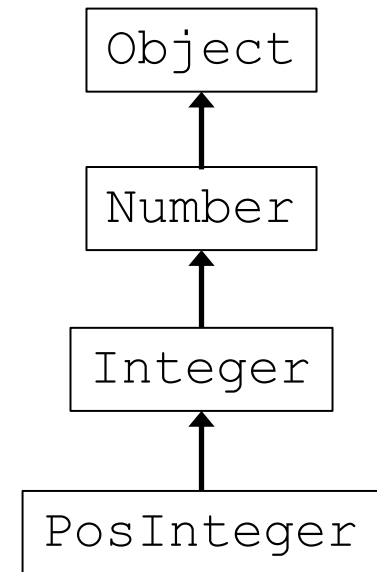
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? extends Integer> list;
```

- Which of these lines is legal?

```
list = new ArrayList<Object>();  
list = new ArrayList<Number>();  
list = new ArrayList<Integer>();  
list = new ArrayList<PosInteger>();  
list = new ArrayList<NegInteger>();
```



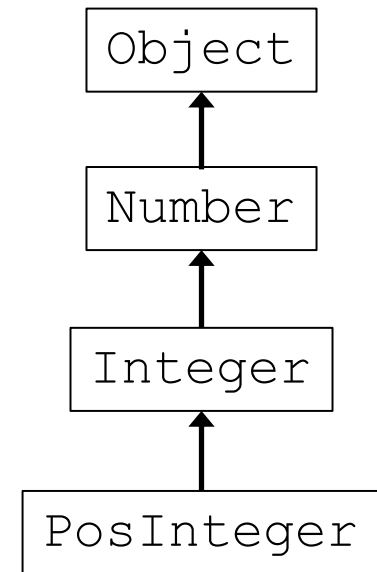
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? extends Integer> list;
```

- Which of these lines is legal?

```
o = list.get(0);  
n = list.get(0);  
i = list.get(0);  
p = list.get(0);
```



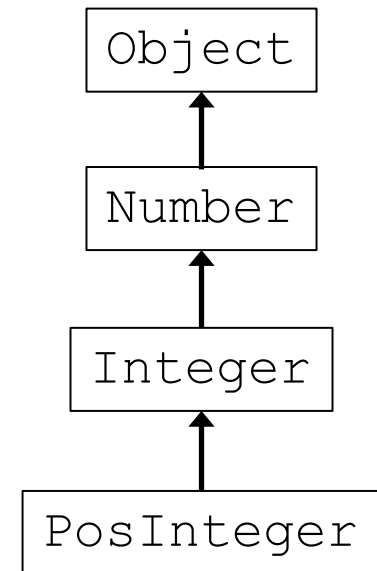
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? extends Integer> list;
```

- Which of these lines is legal?

```
list.add(o);  
list.add(n);  
list.add(i);  
list.add(p);  
list.add(null);
```



**extends** is for producers  
almost no consuming is legal

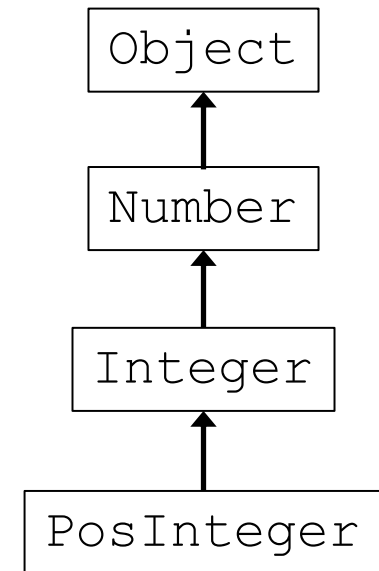
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? super Integer> list;
```

- Which of these lines is legal?

```
list = new ArrayList<Object>();  
list = new ArrayList<Number>();  
list = new ArrayList<Integer>();  
list = new ArrayList<PosInteger>();  
list = new ArrayList<NegInteger>();
```



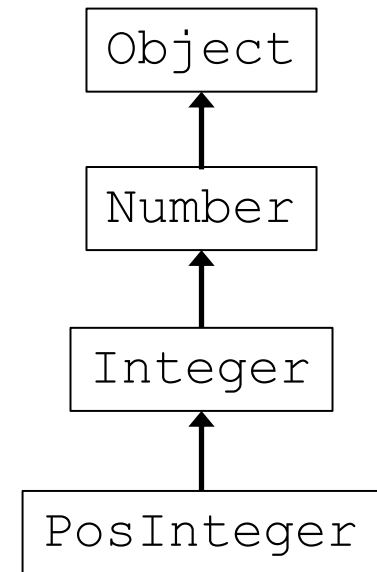
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? super Integer> list;
```

- Which of these lines is legal?

```
list.add(o);  
list.add(n);  
list.add(i);  
list.add(p);  
list.add(null);
```



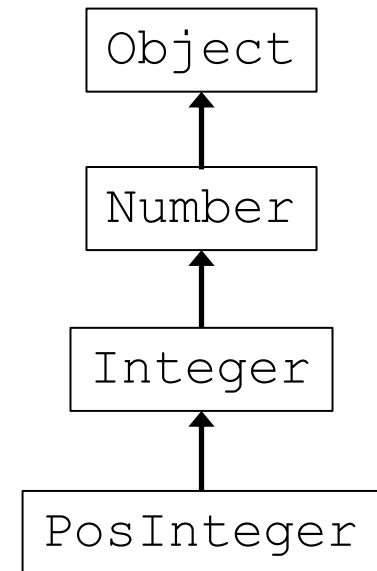
# Legal Operations With Wildcards

---

```
Object o;  
Number n;  
Integer i;  
PosInteger p; // just pretend  
List<? super Integer> list;
```

- Which of these lines is legal?

```
o = list.get(0);  
n = list.get(0);  
i = list.get(0);  
p = list.get(0);
```



**super** is for consumers  
almost no producing is legal

# The Depths of Java Generics





# Arrays

---

- **Arrays are conceptually like** `ArrayList`:

```
Integer[] vals = ...;  
println(vals[0]);  
vals[0] = 10;
```

```
ArrayList<Integer> vals = ...;  
println(vals.get(0));  
vals.set(0, 10);
```

- **operation to get a value by index**
- **operation to set a value by index**

# Arrays

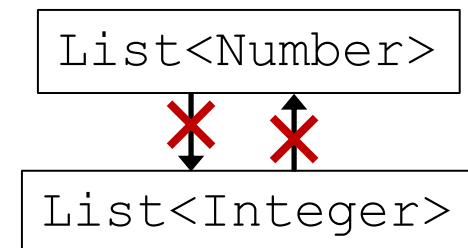
---

- **Arrays are conceptually like** `ArrayList`:

```
Integer[] vals = ...;  
println(vals[0]);  
vals[0] = 10;
```

```
ArrayList<Integer> vals = ...;  
println(vals.get(0));  
vals.set(0, 10);
```

- **How does subtyping work?**
  - **saw** `ArrayList` is invariant



# Arrays

---

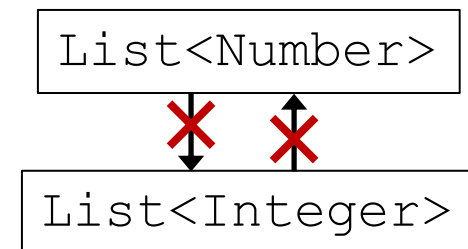
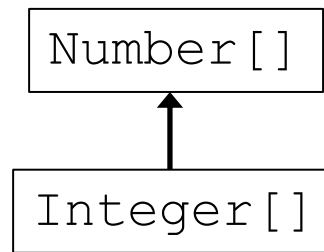
- **Arrays are conceptually like** `ArrayList`:

```
Integer[] vals = ...;  
println(vals[0]);  
vals[0] = 10;
```

```
ArrayList<Integer> vals = ...;  
println(vals.get(0));  
vals.set(0, 10);
```

- **How does subtyping work?**

- **saw** `ArrayList` is invariant
- arrays are **covariant!**  
but... how?



# Arrays Covariance Is Useful

---

- Necessary for things like sorting
  - one example:

```
void swap(LibraryHolding[] arr, int i, int j) {  
    LibraryHolding temp = arr[i];  
    arr[i] = arr[j];  
    arr[j] = temp;  
}
```

```
Book[] books = ...  
maybeSwap(books, 0, 10); // should work & does
```

# Arrays Aren't Really Covariant

---

- Somewhere, this code must fail

```
void replace17 (LibraryHolding[] arr,  
               LibraryHolding h) {  
    arr[17] = h;  
}  
  
Book[] books = ...  
LibraryHolding theWall = new CD("Pink Floyd", ...);  
  
replace17(books, theWall);  
Book b = books[17];    // not really a book...  
b.getChapters();       // doesn't have this method
```

- the last line can't work, so where does it fail?

# Arrays Aren't Really Covariant

---

- Somewhere, this code must fail

```
void replace17 (LibraryHolding[] arr,  
               LibraryHolding h) {  
    arr[17] = h;  
}
```

- fails on this line, attempting to write a `CD` into `Book[]`
- Java checks this at **runtime**, not compile time
  - every array remember its element type
  - all writes are type checked at runtime
    - pay a performance penalty for this

# Type Erasure

---

- **Type parameters become `Object` when compiled**
  - e.g., the following declaration

```
List<String> list = new ArrayList<String>();
```

- **becomes**

```
List<Object> list = new ArrayList<Object>();
```

- **Generics are purely a **compiler** feature**
  - Java cannot double-check type information at runtime...

# Checking Type Parameters at Runtime

---

- Java cannot check type information at runtime:

```
Collection<?> cs = new ArrayList<String>();  
if (cs instanceof Collection<String>) {  
    ...  
}
```

- Java can check that it is a `Collection<?>`
- but not that it is a `Collection<String>`

- As a result, this check is **illegal!**



# Checking Type Parameters at Runtime

---

- Java cannot check type information at runtime:

```
List<?> list = new ArrayList<String>();  
List<String> list2 = (List<String>) list;
```

- Java can check that it is a `List<?>`
  - but not that it is a `List<String>`
- This should be **illegal**, but instead it's a **warning**
    - compiler flag will make this an error
    - these should always be treated as errors...

# Checking Type Parameters at Runtime

---

- Ignore "unchecked cast" warnings at your peril
- Can seriously break the type system

```
public static <T> magicCast(T t, Object o) {  
    return (T) o;  
}
```

```
String s = "abc";  
Integer n = magicCast(3, s); // why not
```

- can turn any type into any other type!
- will result in incredibly **painful** debugging

# Example: Equals in Java

---

```
// Represents an amount of time measured in seconds
class Duration {

    // RI: 0 <= sec < 60
    // AF: obj = 60 * this.min + this.sec
    private int min;
    private int sec;

    public boolean equals(Object o) {
        if (!(o instanceof Duration))
            return false;

        Duration d = (Duration) o;
        return this.min == d.min && this.sec == d.sec;
    }
}
```

- Correct and idiomatic Java

# Generics Causes Problems in Equals

---

```
class Node<E> {  
  
    private E data;  
  
    public boolean equals(Object o) {  
        if (!(o instanceof Node<E>))  
            return false;  
  
        Node<E> n = (Node<E>) o;  
        return this.data.equals(n.data);  
    }  
}
```

- This does not compile!
  - cannot perform these type checks at runtime
- So how do we fix it?

# Generics Causes Problems in Equals

---

```
class Node<E> {  
  
    private E data;  
  
    public boolean equals(Object o) {  
        if (!(o instanceof Node<?>))  
            return false;  
  
        Node<?> n = (Node<?>) o;  
        return this.data.equals(n.data);  
    }  
}
```

- The call to `this.data.equals` will check types
  - not necessary for us to do it again here

# Type Erasure and Arrays

---

```
class Foo<E> {  
    private E[] data;  
  
    public Foo() {  
        this.data = new E[10];  
    }  
}
```

- This is **illegal!**
  - "E" becomes `Object` when compiled but...
  - arrays need to know the element type to check writes
- What should you do?
  - just use `ArrayList` instead