CSE 331 Software Design & Implementation

Spring 2023 Section 4 – Functional Programming II

Administrivia

- HW4 released later today, due Wednesday (4/24) at **11pm**
 - Please! start early and be prepared for a challenge!
 - Give yourself time to come to OH and ask questions on Ed
 - Working on the same issue for hours when you're stuck won't help, ask for help!
- Remember to check the autograder to make sure you pass your own tests! You won't lose *autograder* points, but they should pass!
- Remember to look at your feedback, you may not have lost points on a question but still have helpful comments

Structural Induction – Review

- Let **P(S)** be the claim
- To Prove P(S) holds for any list S, we need to prove two implications: base case and inductive case
 - **Base Case:** prove P(nil)
 - Use any know facts and definitions
 - Inductive Hypothesis: assume P(L) is true for a L: List
 - Use this in the inductive step ONLY
 - Inductive Case: prove P(cons(x, L)) for any x : Z, L : List
 - Direct proof
 - Use know facts and definitions and Inductive Hypothesis
- Assuming we know P(S), if we prove P(cons(x, L)), we then prove recursively that P(S) holds for any List

Structural Induction - 331 Format

The following is the structural induction format we recommend for using in your homework (the staff solution also follows this format)

- 1) Introduction define P(L) to be what we are trying to prove
- 2) Base Case show P(nil) holds
- 3) Inductive Hypothesis assume P(L) is true for an arbitrary list
- 4) Inductive Step show P(cons(a, L)) holds
- 5) Conclusion "We have shown that P(L) holds for any list"

Note: You do not have to follow this format but your solution MUST include all the information above

Defining Function By Cases – Review

- Sometimes we want to define functions by cases
 - **Ex:** define f(n) where $n : \mathbb{Z}$

func

$$f(n) := 2n + 1$$
 if $n \ge 0$
 $f(n) := 0$
 if $n < 0$

- To use the definition f(m), we need to know if m > 0 or not
- This new code structure requires a new proof structure

Proof By Cases – Review

- Split a proof into cases:
 - **Ex:** a = True and a = False or $n \ge 0$ and n < 0
 - These cases needs to be exhaustive
- Ex: func f(n) := 2n + 1 if $n \ge 0$ f(n) := 0 if n < 0Prove that $f(n) \ge n$ for any $n : \mathbb{Z}$

Case $n \ge 0$:
f(n) = 2n + 1 def of f (since $n \ge 0$)
> nSince these 2
cases are
exhaustive,
f(n) >= nCase n < 0:
f(n) = 0 def of f (since n < 0)
 $\ge n$ since n < 0Since these 2
hese 2
cases are
exhaustive,
f(n) >= n
holds in general

pseudo-sort: takes a list of numbers as an argument, "looks at the first two numbers in the list, orders the pair to place the smaller of the two in the front, and then continues through the following pairs in the list after the first element"



(a) Write a formal definition using recursion

(b) Show by example that pseudo-sort does not actually sort the list

```
const s = sum(L);
...
return 2 * s; // = sum(twice(L))
```

Prove this code is correct by showing that **sum(twice(S)) = 2 sum(S)** holds for any list S by structural induction.

 $\begin{aligned} & \textbf{func swap(nil)} & \coloneqq & \text{nil} \\ & \text{swap(cons}(a, \text{nil})) & \coloneqq & \text{cons}(a, \text{nil}) & \text{for any } a : \mathbb{Z} \\ & \text{swap(cons}(a, \text{cons}(b, L))) & \coloneqq & \text{cons}(b, \text{cons}(a, \text{swap}(L))) \\ & & \text{for any } a, b : \mathbb{Z} \text{ and } L : \text{List} \end{aligned}$

Prove by cases that swap(cons(a, L)) \neq nil for any integer a : \mathbb{Z} and list L.

sum(twice-evens(S)) + sum(twice-odds(S)) = 3 sum()

Prove that this holds for any list *S* by structural induction.