

CSE 331

Trees

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What We Get from Reasoning

- **If the proof works, the code is correct**
 - why reasoning is useful for finding bugs
- **If the code is incorrect, the proof will not work**
- **If the proof does not work, the code is probably wrong**
 - could potentially be an issue with the proof (e.g., two “<”s)
 - but that is a rare occurrence

Proof by Calculation

Finding Facts at a Return Statement

- Consider this code

```
// Inputs a and b must be integers.  
// Returns a non-negative integer.  
const f = (a: number, b: number): number => {  
  const L: List = cons(a, cons(b, nil));  
  if (a >= 0 && b >= 0)  
    return sum(L);  
  ...  
}
```

find facts by reading along path
from top to return statement

- Known facts include “ $a \geq 0$ ”, “ $b \geq 0$ ”, and “ $L = \text{cons}(\dots)$ ”

Proving Correctness with Conditionals

```
// Inputs x and y are integers.
// Returns a number less than x.
const f = (x: number, y, number): number => {
  if (y < 0) {
    return x + y;
  } else {
    return x - 1;
  }
};
```

- Known fact in then (top) branch: “ $y \leq -1$ ”

$x + y$

Proving Correctness with Conditionals

```
// Inputs x and y are integers.
// Returns a number less than x.
const f = (x: number, y, number): number => {
  if (y < 0) {
    return x + y;
  } else {
    return x - 1;
  }
};
```

- Known fact in then (top) branch: “ $y \leq -1$ ”

$$\begin{array}{ll} x + y & \leq x + -1 & \text{since } y \leq -1 \\ & < x + 0 & \text{since } -1 < 0 \\ & = x & \end{array}$$

Proving Correctness with Conditionals

```
// Inputs x and y are integers.
// Returns a number less than x.
const f = (x: number, y, number): number => {
  if (y < 0) {
    return x + y;
  } else {
    return x - 1;
  }
};
```

- Known fact in else (bottom) branch: “ $y \geq 0$ ”

$x - 1$

Proving Correctness with Conditionals

```
// Inputs x and y are integers.
// Returns a number less than x.
const f = (x: number, y, number): number => {
  if (y < 0) {
    return x + y;
  } else {
    return x - 1;
  }
};
```

- Known fact in else (bottom) branch: “ $y \geq 0$ ”

$$\begin{array}{lll} x - 1 & < x + 0 & \text{since } -1 < 0 \\ & = x & \end{array}$$

Proving Correctness with Conditionals

```
// Inputs x and y are integers.
// Returns a number less than x.
const f = (x: number, y, number): number => {
  if (y < 0) {
    return x + y;
  } else {
    return x - 1;
  }
};
```

- **Conditionals give us extra known facts**

- get known facts from

1. specification
2. conditionals
3. constant declarations

find facts by reading along path
from top to the return statement

Proving Correctness with Multiple Claims

- Need to check the claim from the spec at each `return`
- If spec claims multiple facts, then we must prove that each of them holds

```
// Inputs x and y are integers with x < y - 1
// Returns a number less than y and greater than x.
const f = (x: number, y, number): number => { .. };
```

- multiple known facts: $x : \mathbb{Z}$, $y : \mathbb{Z}$, and $x < y - 1$
- multiple claims to prove: $x < r$ and $r < y$
where “r” is the return value
- requires *two* calculation blocks

Recall: Max With an Imperative Specification

```
// Returns a if a >= b and b if a < b
const max = (a: number, b, number): number => {
  if (a >= b) {
    return a;
  } else {
    return b;
  }
};
```

Level 0

Example Correctness with Conditionals

```
// Returns r with (r=a or r=b) and r >= a and r >= b
const max = (a: number, b, number): number => {
  if (a >= b) {
    return a;
  } else {
    return b;
  }
};
```

Level 1

- Three different facts to prove at each **return**
- Two known facts in each branch (return value is “r”):
 - then branch: $a \geq b$ and $r = a$
 - else branch: $a < b$ and $r = b$

Example Correctness with Conditionals

```
// Returns r with (r=a or r=b) and r >= a and r >= b
const max = (a: number, b, number): number => {
  if (a >= b) {
    return a;           Know  $a \geq b$  and  $r = a$ 
  } else {
    return b;
  }
};
```

- **Correctness of return in “then” branch:**
 - $r = a$ holds so “ $r = a$ or $r = b$ ” holds,
 - $r = a$ holds so “ $r \geq a$ ” holds, and

$r = a$
 $\geq b$ since $a \geq b$

Example Correctness with Conditionals

```
// Returns r with (r=a or r=b) and r >= a and r >= b
const max = (a: number, b, number): number => {
  if (a >= b) {
    return a;
  } else {
    return b;          Know a < b and r = b
  }
};
```

- **Correctness of return in “else” branch:**

- $r = b$ holds so “ $r = a$ or $r = b$ ” holds,
- $r = b$ holds so “ $r \geq b$ ” holds, and
- $r \geq a$ holds since we have $r > a$:

$r = b$
 $> a$ **since $a < b$**

Sum of a List

```
// a and b must be integers
const f = (a: number, b: number): number => {
  const L: List = cons(a, cons(b, nil));
  const s: number = sum(L); // = a + b
  ...
};
```

- Can prove the claim in the comments by calculation

sum(cons(a, cons(b, nil)))	
= a + sum(cons(b, nil))	def of sum
= a + b + sum(nil)	def of sum
= a + b	def of sum

func sum(nil) := 0
sum(cons(x, L)) := x + sum(L) for any $x \in \mathbb{Z}$ and any $L \in \text{List}$

Sum of a List

```
// a and b must be integers
const f = (a: number, b: number): number => {
  const L: List = cons(a, cons(b, nil));
  const s: number = sum(L); // = a + b
  ...
}
```

- Can prove the claim in the comments by calculation

$$\text{sum}(\text{cons}(a, \text{cons}(b, \text{nil}))) = \dots = a + b$$

- For which values of a and b does this hold?

holds for any $a \in \mathbb{Z}$ and $b \in \mathbb{Z}$

What We Have Proven

- We proved by calculation that

$$\text{sum}(\text{cons}(a, \text{cons}(b, \text{nil}))) = a + b$$

- This holds for any $a \in \mathbb{Z}$ and $b \in \mathbb{Z}$
- We have proven *infinitely* many facts
 - $\text{sum}(\text{cons}(3, \text{cons}(5, \text{nil}))) = 8$
 - $\text{sum}(\text{cons}(-5, \text{cons}(2, \text{nil}))) = -3$
 - ...
 - replacing all the ‘a’s and ‘b’s with those numbers gives a calculation proving the “=” for those numbers

What We Have Proven

- We proved by calculation that

$$\text{sum}(\text{cons}(a, \text{cons}(b, \text{nil}))) = a + b \quad \text{for any } a, b \in \mathbb{Z}$$

- We can use this fact for any a and b we choose
 - our proof is a “recipe” that can be used for any a and b
 - just as a function can be used with any argument values, our proof can be used with any values for the “any” variables (any values satisfying the specification)
 - use “for any ...” to make clear which things are variables
- This is called a “direct proof” of the “for any” claim

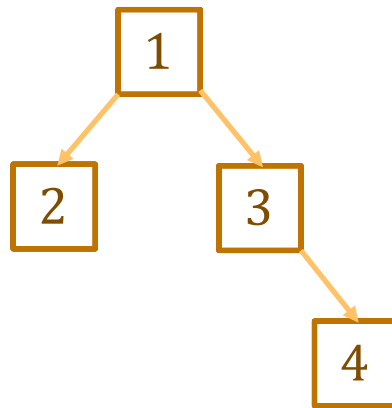
Binary Trees

Binary Trees

```
type Tree := empty | node(x :  $\mathbb{Z}$ , L : Tree, R : Tree)
```

- **Inductive definition of binary trees of integers**

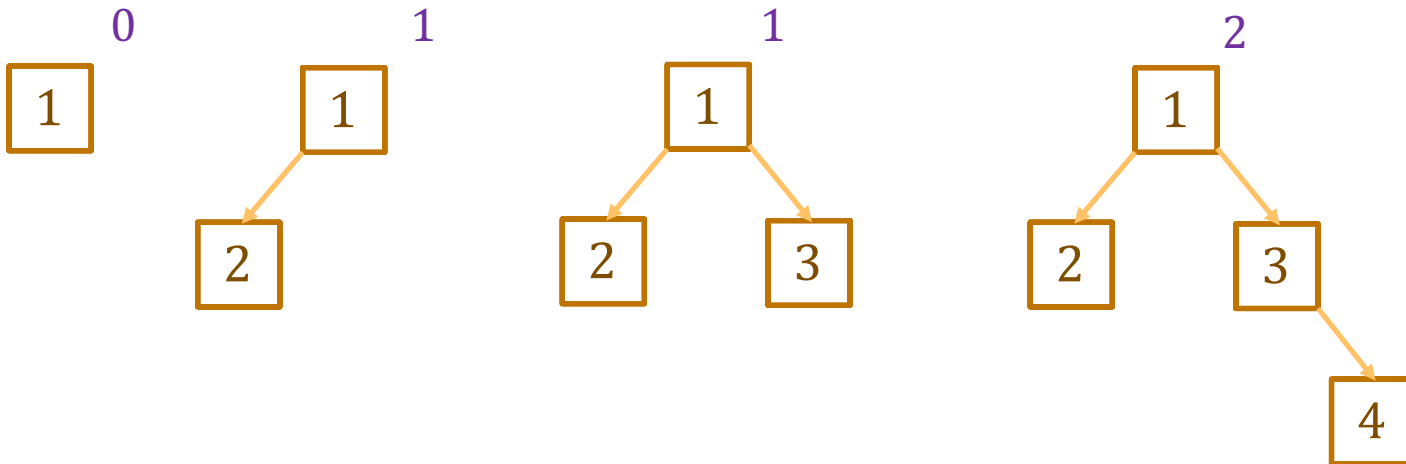
```
node(1, node(2, empty, empty), node(3, empty, node(4, empty, empty)))
```



Height of a Tree

`type Tree := empty | node(x: \mathbb{Z} , L: Tree, R: Tree)`

- Height of a tree: “maximum steps to get to a leaf”

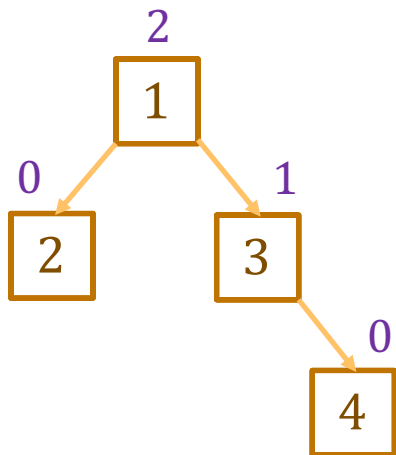


Height of a Tree

`type Tree := empty | node(x: \mathbb{Z} , L: Tree, R: Tree)`

- **Mathematical definition of height**

`func height(empty) :=`
`height(node(x, L, R)) :=`



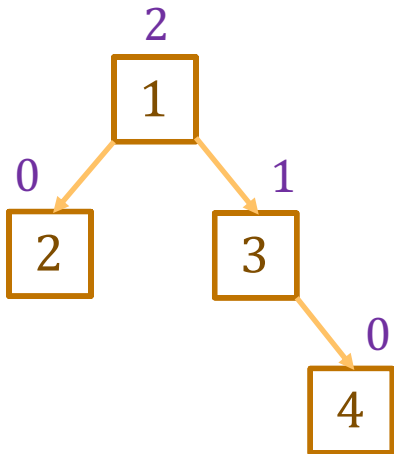
for any $x \in \mathbb{Z}$ and any $L, R \in \text{Tree}$

Height of a Tree

`type Tree := empty | node(x: \mathbb{Z} , L: Tree, R: Tree)`

- **Mathematical definition of height**

`func height(empty) := -1`
`height(node(x, L, R)) := 1 + max(height(L), height(R))`
for any $x \in \mathbb{Z}$ and any $L, R \in \text{Tree}$



Using Definitions in Calculations

```
func height(empty)      := -1
   height(node(x, L, R)) := 1 + max(height(L), height(R))
                           for any  $x \in \mathbb{Z}$  and any  $L, R \in \text{Tree}$ 
```

- **Suppose** “ $T = \text{node}(1, \text{empty}, \text{node}(2, \text{empty}, \text{empty}))$ ”
- **Prove that** $\text{height}(T) = 1$

$\text{height}(T) =$

Using Definitions in Calculations

func height(empty) := -1
 height(node(x, L, R)) := 1 + max(height(L), height(R))
 for any $x \in \mathbb{Z}$ and any $L, R \in \text{Tree}$

- **Suppose** “ $T = \text{node}(1, \text{empty}, \text{node}(2, \text{empty}, \text{empty}))$ ”
- **Prove that** $\text{height}(T) = 1$

height(T)	= height(node(1, empty, node(2, empty, empty)))	since T = ...
	= 1 + max(height(empty), height(node(2, empty, empty)))	def of height
	= 1 + max(-1, height(node(2, empty, empty)))	def of height
	= 1 + max(-1, 1 + max(height(empty), height(empty)))	def of height
	= 1 + max(-1, 1 + max(-1, -1))	def of height (x 2)
	= 1 + max(-1, 1 + -1)	def of max
	= 1 + max(-1, 0)	
	= 1 + 0	def of max
	= 1	

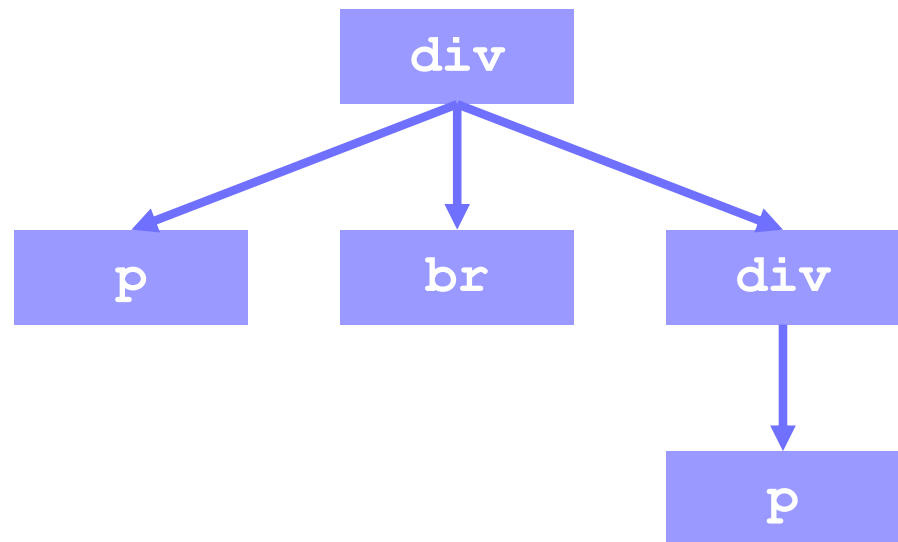
Trees

- **Trees are inductive types with a constructor that has 2+ recursive arguments**
- **These come up all the time...**
 - no constructors with recursive arguments = “generalized enums”
 - constructor with 1 recursive arguments = “generalized lists”
 - constructor with 2+ recursive arguments = “generalized trees”
- **Some prominent examples of trees:**
 - HTML: used to describe UI
 - JSON: used to describe just about any data

Recall: HTML

- Nesting structure describes the tree

```
<div>
  <p id="firstParagraph"> Some Text </p>
  <br>
  <div>
    <p>Hello</p>
  </div>
</div>
```



Custom Tags

- The React library lets you write “custom tags”
 - functions that return HTML

```
return (  
  <div>  
    <p>Hi, Alice!</p>  
    <p>Hi, Bob!</p>  
  </div>);
```

can become

```
return (  
  <div>  
    <SayHi name={ "Alice" } />  
    <SayHi name={ "Bob" } />  
  </div>);
```

Custom Tags

- The React library lets you write “custom tags”

```
return (  
  <div>  
    <SayHi name={"Alice"}/>  
    <SayHi name={"Bob"}/>  
  </div>);
```

makes two calls to this function

```
const SayHi = (props: {name: string}): JSX.Element => {  
  return <p>Hi, {props.name}</p>;  
};
```

- attributes are passed as a record argument (“props”)

Custom Tags

```
return (  
  <div>  
    <SayHi name={"Alice"} lang={"es"}/>  
    <SayHi name={"Bob"}/>  
  </div>);
```

makes two calls to this function

```
type SayHiProps = {name: string, lang?: string};  
  
const SayHi = (props: SayHiProps): JSX.Element => {  
  if (props.lang === "es") {  
    return <p>Hola, {props.name}</p>;  
  } else {  
    return <p>Hi, {props.name}</p>;  
  }  
};
```

Custom Tags

- The React library lets you write “custom tags”
 - attributes are passed as a record argument (“props”)
- In `render`, React will paste the parts together:

```
<div>
  <SayHi name={"Alice"} lang={"es"}/>
  <SayHi name={"Bob"}/>
</div>
```

becomes

```
<div>
  <p>Hola, Alice!</p>
  <p>Hi, Bob!</p>
</div>
```

Custom Tags

- HTML literal syntax allows any tags

```
return (  
  <div>  
    <SayHi name={"Alice"} lang={"es"}/>  
    <SayHi name={"Bob"}/>  
  </div>);
```

- evaluates to a tree with two nodes with tag name “SayHi”
 - this matters when *testing* (comes up in HW3)
- React’s `render` method is what calls `SayHi`
 - HTML returned is *substituted* where the “SayHi” tag was

React Render

- React's `render` pastes strings together

```
const name: String = "Fred";  
return <p>Hi, {name}</p>;
```

returns a different tree than

```
return <p>Hi, Fred</p>;
```

- in first tree, “p” tag has one child
 - in second tree, “p” tag has two children
 - render method concatenates text children into one string
- These differences matter for **testing!**

React Render

- React's `render` pastes arrays into child list

```
const L = [<span>Hi</span>, <span>Fred</span>];  
return <p>{L}</p>;
```

returns a different tree than

```
return <p><span>Hi</span><span>Fred</span></p>;
```

- in first tree, “p” tag has one child
 - in second tree, “p” tag has two children
 - render method turns the first into the second
- These differences matter for **testing!**