
CSE 331

Software Design & Implementation

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Data Abstraction: Abstract Data Types (ADTs)

Administrivia

- HW3 due tomorrow night. When???
11 PM pacific time!
 - Please double check hw3-final tag is correct and be sure there are no gitlab runner bugs, etc.
- Sections tomorrow: HW4 – implement rational numbers and related ADTs given a detailed specification, test with supplied JUnit tests, and more...
 - Assignment posted later today
 - Starter code for hw4 will be pushed to repos later today or tonight

Administrivia (2)

- New on the web: writeups describing CSE 331 concepts and conventions for specifications, rep invariants, abstraction functions, and other topics.
 - (these include the “readings” for this part of the course)

Outline

This lecture:

1. What is an Abstract Data Type (ADT)?
2. How to specify an ADT?
 - Immutable
 - Mutable
3. Design methodology for ADTs

Very related next lectures:

- Representation invariants
- Abstraction functions

Two distinct, complementary ideas for reasoning about ADT implementations

Procedural and data abstractions

Procedural abstraction:

- Abstract from details of *procedures* (e.g., methods)
- A specification mechanism
- Satisfy the specification with an implementation

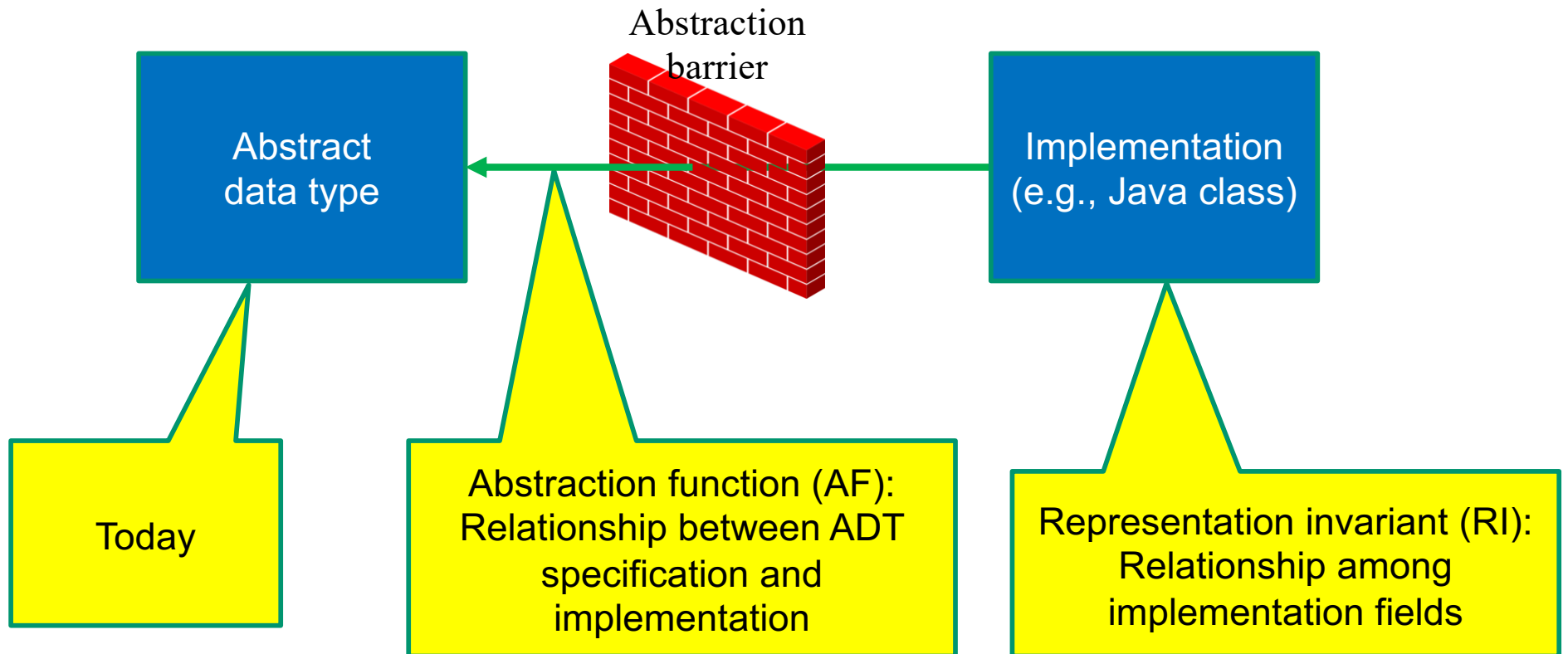
Data abstraction:

- Abstract from details of *data representation*
- Also a specification mechanism
 - And a way of thinking about programs and design
- Standard terminology: **Abstract Data Type**, or **ADT**

Outline of next 3 lectures

ADT
specification

ADT
implementation



Why we need Data Abstractions (ADTs)

Organizing and manipulating data is pervasive

- Inventing and describing algorithms is less common

Start your design by **designing data abstractions**

- How will relevant data be organized
- What operations will be permitted on the data by clients
- Secondary: how is data stored/represented? What algorithms manipulate the data?

Potential problems with choosing a data abstraction:

- Decisions about data structures often made too early
- Duplication of effort in creating derived data
- Very hard to change key data structures (modularity!)

An ADT is a set of operations

- ADT abstracts from the *organization* to *meaning* of data
- ADT abstracts from *structure* to *use*
- A **type** is a **set of operations**
 - `create, getBase, getAltitude, getBottomAngle, ...`
- Operations are the only way clients can access data
- Representation should not matter to the client
 - So hide it from the client

```
class RightTriangle {  
    private float base;  
    private float altitude;  
}
```

```
class RightTriangle {  
    private float base;  
    private float hypot;  
    private float angle;  
}
```


An abstract data type defines a class of abstract objects which is completely characterized by the operations available on those objects ...

When a programmer makes use of an abstract data object, he [sic] is concerned only with the behavior which that object exhibits but not with any details of how that behavior is achieved by means of an implementation...

-- Programming with Abstract Data Types, Barbara Liskov and Stephen Zilles 1974 (!)



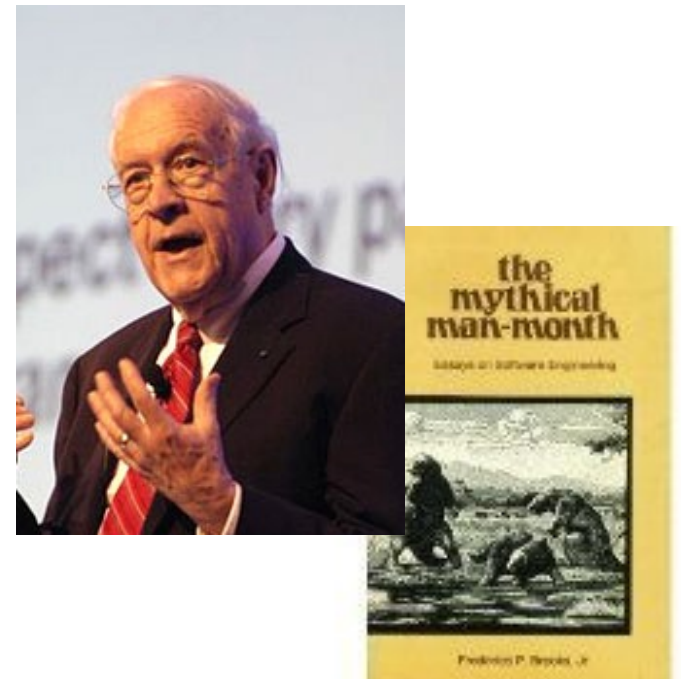
Bad programmers worry about the code. Good programmers worry about data structures and their relationships.

-- Linus Torvalds



Show me your flowcharts and conceal your tables, and I shall continue to be mystified. Show me your tables, and I won't usually need your flowcharts; they'll be obvious.

-- Fred Brooks



Are these classes the same?

```
class Point {
    public float x;
    public float y;
}

class Point {
    public float r;
    public float theta;
}
```

Different: cannot replace one with the other in a program

Same: both classes implement the concept “2-d point”

Goal of ADT methodology is to express the sameness:

- Clients depend only on the concept “2-d point”

Benefits of ADTs

If clients “respect” or “are forced to respect” data abstractions...

- For example, “it’s a 2-D point with these operations...”
- Can delay decisions on how ADT is implemented
- Can fix bugs by changing how ADT is implemented
- Can change algorithms
 - For performance
 - In general or in specialized situations
- ...

We talk about an “*abstraction barrier*”

- A good thing to have and not *cross* (also known as *violate*)

Concept of 2-d point, as an ADT

```
class Point {
    // A 2-d point exists in the plane, ...
    public float x();
    public float y();
    public float r();
    public float theta();

    // ... can be created, ...
    public Point(); // new point at (0,0)
    public Point centroid(Set<Point> points);

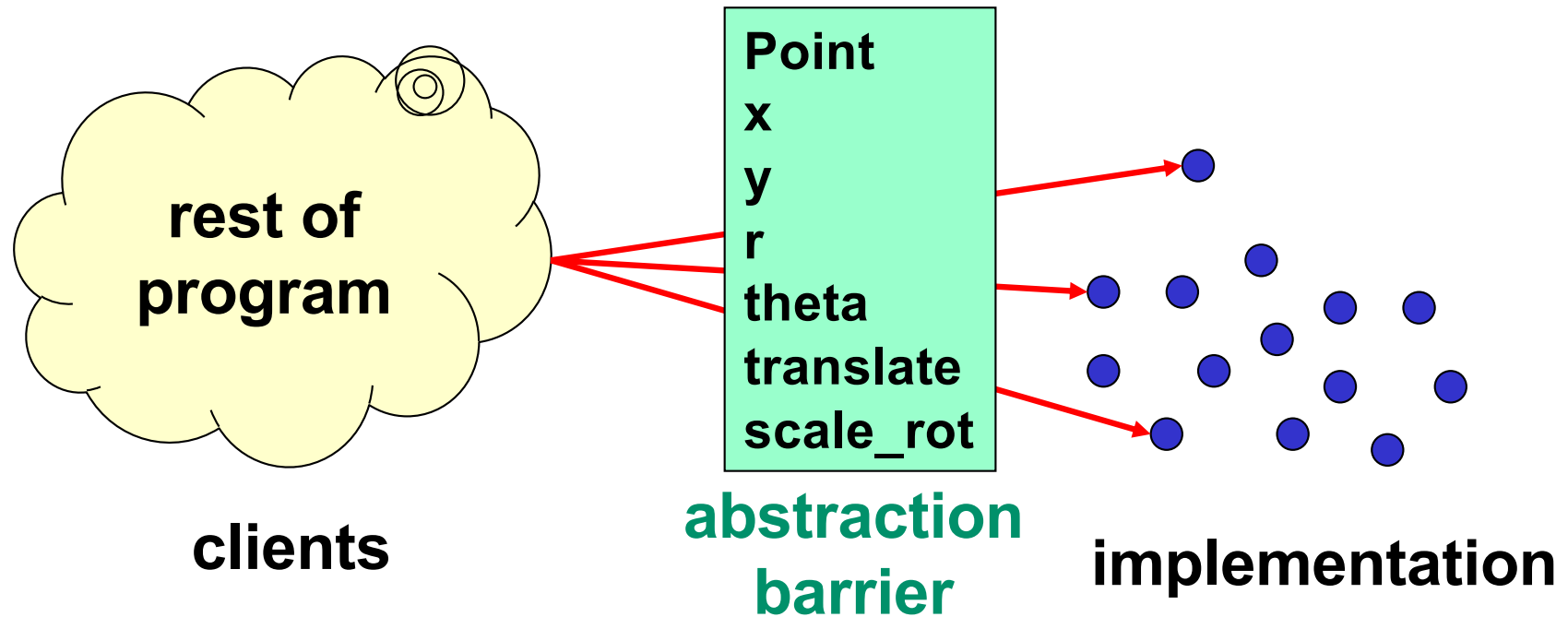
    // ... can be moved, ...
    public void translate(float delta_x,
                          float delta_y);
    public void scaleAndRotate(float delta_r,
                               float delta_theta);
}
```

Observers

Creators/
Producers

Mutators

Abstract data type = objects + operations



- Implementation is hidden
- The only operations on objects of the type are those provided by the abstraction

Specifying a data abstraction

- An *abstract state*
 - Not the (concrete) representation in terms of fields, objects, ...
 - Although some of the concrete state might coincide (implement directly) parts of the abstract state
 - “Does not exist” but used to specify the operations
- A collection of *operations* (procedural abstractions)
 - *Not* a collection of procedure implementations
 - Specified in terms of abstract state
 - No other way to interact with the data abstraction
 - Four types of operations: creators, observers, producers, mutators

Specifying an ADT

Immutable

1. overview
2. abstract state (fields)
3. creators
4. observers
5. producers
6. ~~mutators~~

Mutable

1. overview
2. abstract state (fields)
3. creators
4. observers
5. producers (rare)
6. mutators

- Creators: return new ADT values (e.g., Java constructors)
- Producers: ADT operations that return new ADT values
- Mutators: Modify a value of an ADT
- Observers: Return information about an ADT

Implementing an ADT

To implement a data abstraction (e.g., with a Java class):

- See next two lectures
- This lecture is just about specifying an ADT
- *Nothing* about the concrete representation appears in the specification

Poly, an immutable datatype: overview

```
/**
```

```
* A Poly is an immutable polynomial with  
* integer coefficients. A typical Poly is
```

```
*  $c_0 + c_1x + c_2x^2 + \dots$ 
```

```
**/
```

```
class Poly {
```

Abstract state (specification fields)



Overview:

- Always state whether mutable or immutable
- Define an abstract model for use in operation specifications
 - Difficult and vital!
 - Appeal to math if appropriate
 - Give an example (reuse it in operation definitions)
- State in specifications is *abstract*, not concrete

Poly: creators

```
// effects: makes a new Poly = 0
```

```
public Poly()
```

```
// effects: makes a new Poly =  $cx^n$ 
```

```
// throws: NegExponent if  $n < 0$ 
```

```
public Poly(int c, int n)
```

Creators

- New object, not part of pre-state: in **effects**, not **modifies**
- Overloading: distinguish procedures of same name by parameters (Example: two **Poly** constructors)

Footnote: slides omit full JavaDoc comments to save space; style might not be perfect either – focus on main ideas

Poly: observers

```
// returns: the degree of this,  
//   i.e., the largest exponent with a  
//   non-zero coefficient.  
//   Returns 0 if this = 0.  
public int degree()  
  
// returns: the coefficient of the term  
//   of this whose exponent is d  
// throws: NegExponent if d < 0  
public int coeff(int d)
```

Notes on observers

Observers

- Used to obtain information about objects of the type
- Return values of other types
- Never modify the abstract value
- Specification uses the abstraction from the overview

this

- The particular `Poly` object being accessed
- *Target* of the invocation
- Also known as the *receiver*

```
Poly x = new Poly(4, 3);  
int c = x.coeff(3);  
System.out.println(c);    // prints 4
```

Poly: producers

```
// returns: this + q (as a Poly)
```

```
public Poly add(Poly q)
```

```
// returns: the Poly equal to this * q
```

```
public Poly mul(Poly q)
```

```
// returns: -this
```

```
public Poly negate()
```

Notes on producers

- Operations on a type that create other objects of the type
- Common in immutable types like `java.lang.String`
 - `String substring(int offset, int len)`
- No side effects
 - Cannot change the abstract value of existing objects

IntSet, a mutable datatype: overview and creator

```
// Overview: An IntSet is a mutable,  
// unbounded set of integers.  A typical  
// IntSet is { x1, ..., xn }.  
class IntSet {  
  
// effects: makes a new IntSet = {}  
public IntSet()  
  
}
```


IntSet: observers

```
// returns: true if and only if  $x \in$  this
```

```
public boolean contains(int x)
```

```
// returns: the cardinality of this
```

```
public int size()
```

```
// returns: some element of this
```

```
// throws: EmptyException when size()==0
```

```
public int choose()
```

IntSet: mutators

```
// modifies: this  
// effects:  thispost = thispre ∪ {x}  
public void add(int x)
```

```
// modifies: this  
// effects:  thispost = thispre - {x}  
public void remove(int x)
```

Notes on mutators

- Operations that modify an element of the type
- Rarely modify anything (available to clients) other than `this`
 - List `this` in modifies clause (if appropriate)
- Typically have no return value
 - “Do one thing and do it well”
 - (Sometimes return “old” value that was replaced)
- Mutable ADTs may have producers too, but that is less common

Perspective

- Manipulating and presenting data is pervasive
 - choosing how to organize that data is key design problem
 - inventing and describing algorithms is less common
- Data abstractions (ADTs) are a fundamental design idea
 - Client perspective: a set of operations – a specification
- Is everything an ADT? No – there are classes/modules that are simply bundles of data, or just collections of procedures
 - But data abstractions are everywhere and a key design and problem-solving principle

Next time

- Implementing ADTs
 - Picking concrete representations for data abstractions (“the rep” – instance variables)
 - Reasoning about implementations: rep invariants and abstraction functions