CSE 331

Software Design & Implementation Topic: Rep. Exposure; Abstraction Functions

Discussion: How was your long weekend?

Reminders

- Make sure to check Gitlab when submitting
 - must commit, tag, and pass the Gitlab pipeline
- Uploaded replacement recording for Specifications

Upcoming Deadlines

• HW3 due Thursday (7/7)

Last Time...

Today's Agenda

- Abstract Data Types
- ADTs in Java
 - overview
 - abstract state
 - creators
 - observers
 - producers
 - mutators
- Representation Invariants

- Representation Exposure
- Abstraction Functions
- Intro to Testing

Abstract Data Type (ADT)

ADT abstracts from the *organization* to *meaning* of data

- details of data structures are hidden from the client
- client see only the operations that provided

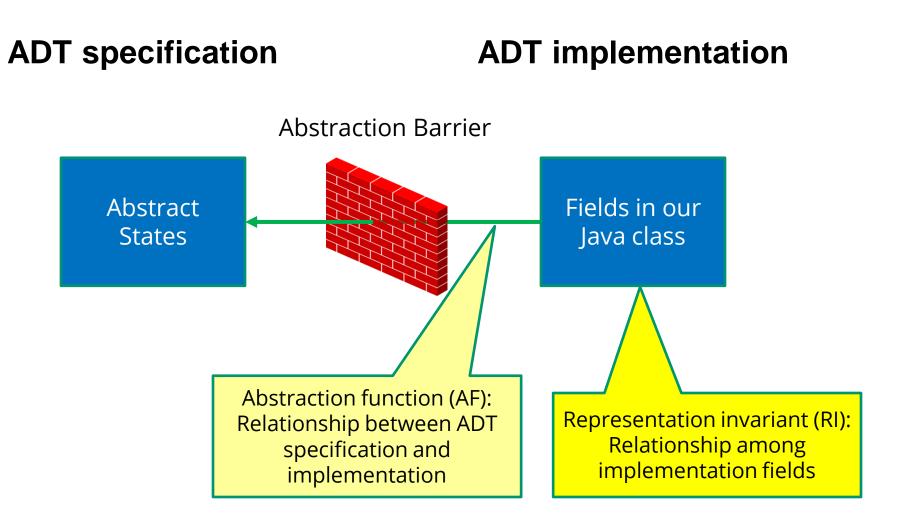
Choose a representation so that:

- it is possible to implement required operations
- the most frequently used operations are efficient / simple / ...
 - abstraction allows the rep to change later
 - almost always better to start simple

Then use **reasoning** to verify the operations are correct

- two intellectual tools are helpful for this...

Data abstraction outline



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Connecting implementations to specs

For implementers / debuggers / maintainers of the implementation:

Representation Invariant: maps Object → boolean

- defines the set of valid concrete values
- must hold before and after any public method is called
- no object should *ever* violate the rep invariant
 - such an object has no useful meaning

Abstraction Function: maps Object \rightarrow abstract state

- we'll discuss this later!

/** Represents a mutable circle in the plane. For example, * it can be a circle with center (0,0) and radius 1. */ public class Circle {

```
// Rep invariant: center != null and edge != null
// and !center.equals(edge)
private Point center, edge;
```

```
// Abstraction function:
// AF(this) = a circle with center at this.center
// and radius this.center.distanceTo(this.edge)
```

```
// ...
}
```

/** An immutable polynomial with integer coefficients.

* Examples include 0, 2x, and x + 3x² + 5x. */
public class IntPoly {

// Rep invariant: terms != null and // no two terms have the same degree and // terms is sorted in descending order by degree private final LinkedList<IntTerm> terms;

// Abstraction function:
// AF(this) = sum of monomials in this.terms

Defensive Programming with ADTs

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Remember that representation invariants should hold <u>before</u> and <u>after</u> each method in the public specification.

Should you write code to check that the rep invariant holds?

- Yes, if it's inexpensive [depends on the invariant]
- Yes, for debugging [even when it's expensive]
- Often hard to justify turning the checking off
 - better argument is removing clutter (improve understandability)

A great debugging technique:

Catch bugs by implementing and using a function to check the rep-invariant

Example: CharSet ADT

```
// Overview: A CharSet is a finite mutable set of Characters
// @effects: creates a fresh, empty CharSet
public CharSet() {...}
// @modifies: this
// @effects: this changed to this + {c}
public void insert(Character c) {...}
// @modifies: this
// @effects: this changed to this - {c}
public void delete(Character c) {...}
// @return: true iff c is in this set
public boolean member(Character c) {...}
// @return: cardinality of this set
public int size() {...}
```

Example: CharSet ADT

// Rep invariant: elts != null and // elts has no nulls and no dups // AF(this) = list of chars in elts private List<Character> elts;

Checking the rep invariant

How do we check whether this invariant holds?

```
public void delete(Character c) {
```

}

elts.remove(c); // removes 0 or 1 copies of c

Checking the rep invariant

Rule of thumb: check on entry and on exit (why?)

```
public void delete(Character c) {
    checkRep();
    elts.remove(c); // removes 0 or 1 copies of c
    checkRep();
}
```

```
// Verify that elts contains no nulls or dups
private void checkRep() {
   assert elts != null;
   for (int i = 0; i < elts.size(); i++) {
      assert elts.get(i) != null;
      assert elts.indexOf(elts.get(i)) == i;
   }
}</pre>
```

Practice *defensive* programming

- Question is not: will you make mistakes? You will.
- Question is: will you **catch** those mistakes before users do?
- Write and incorporate code designed to catch the errors you make
 - check rep invariant on entry and exit (of mutators)
 - check preconditions (don't trust other programmers)
 - check postconditions (don't trust yourself either)
- Checking the rep invariant helps *discover* errors while testing
- Reasoning about the rep invariant helps *discover* errors while coding

Practice *defensive* programming

- Checking pre- and post-conditions and rep invariants is one tip
- More of these in Effective Java
 - first required reading (see calendar for items)
- Focus on defensive programming against **subtle bugs**
 - obvious bugs (e.g., crashing every time) will be caught in testing
 - subtle bugs that only occasionally cause problems can sneak out
 - be especially defensive against (and scared of) these

Listing the elements of a CharSet

Consider adding the following method to **CharSet**

```
// returns: a List containing the members of this
public List<Character> getElts();
```

Consider this implementation:

```
public List<Character> getElts() { return elts; }
```

Does this implementation preserve the rep invariant? *Can't say!*

Representation exposure

Consider this client code (outside the **CharSet** implementation):

```
CharSet s = new CharSet();
Character a = new Character('a');
s.insert(a);
s.getElts().add(a);
s.delete(a);
if (s.member(a)) ...
```

- Representation exposure is external access to the rep
- Representation exposure is almost always bad
 - can cause bugs that will be **very hard to detect**
- Rule #1: Don't do it!
- Rule #2: If you do it, document it clearly and then feel guilty about it!

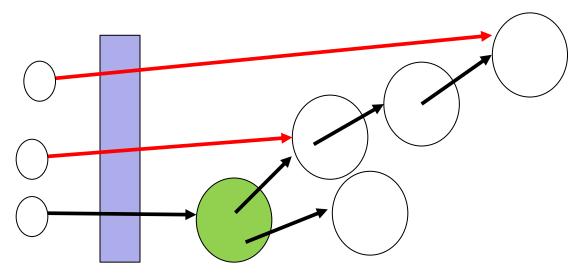
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Avoiding representation exposure

- *Understand* what representation exposure is
- *Design* ADT implementations to make sure it doesn't happen
- Treat rep exposure as a bug: *fix* your bugs
 - absolutely must avoid in libraries with many clients
 - can allow (but feel guilty) for code with few clients
- *Test* for it with *adversarial clients:*
 - pass values to methods and then mutate them
 - mutate values returned from methods

private is not enough

- Making fields **private** does *not* suffice to prevent rep exposure
 - see our example
 - issue is *aliasing of mutable data outside the abstraction*



- So private is a hint to you: no aliases outside abstraction to references to mutable data reachable from private fields
- Three general ways to avoid representation exposure...

Avoiding rep exposure (way #1)

- One way to avoid rep exposure is to make copies of all data that cross the abstraction barrier
 - Copy in [parameters that become part of the implementation]
 - Copy out [results that are part of the implementation]
- Examples of copying (assume **Point** is a mutable ADT):

...

```
class Line {
   private Point s, e;
   public Line(Point s, Point e) {
     this.s = new Point(s.x,s.y);
     this.e = new Point(e.x,e.y);
   }
   public Point getStart() {
     return new Point(this.s.x,this.s.y);
   }
}
```

Avoiding rep exposure (way #2)

- One way to avoid rep exposure is to exploit the immutability of (other) ADTs the implementation uses
 - aliasing is no problem if nobody can change data
 - have to mutate the rep to break the rep invariant
- Examples (assuming **Point** is an *immutable* ADT):

...

```
class Line {
   private Point s, e;
   public Line(Point s, Point e) {
     this.s = s;
     this.e = e;
   }
   public Point getStart() {
     return this.s;
   }
```

Alternative #3

```
// returns: elts currently in the set
public List<Character> getElts() { // version 1
   return new ArrayList<Character>(elts);//copy out!
}
public List<Character> getElts() { // version 2
```

```
return Collections.unmodifiableList(elts);
```

From the JavaDoc for Collections.unmodifiableList:

Returns an unmodifiable view of the specified list. This method allows modules to provide users with "read-only" access to internal lists. Query operations on the returned list "read through" to the specified list, and attempts to modify the returned list... result in an UnsupportedOperationException.

The good news

```
public List<Character> getElts() { // version 2
  return Collections.unmodifiableList(elts);
}
```

- Clients cannot *modify (mutate)* the rep
 - cannot break the rep invariant
- (For long lists,) more efficient than copy out
- Uses standard libraries

The bad news

```
public List<Character> getElts() { // version 1
  return new ArrayList<Character>(elts);//copy out!
}
public List<Character> getElts() { // version 2
```

```
return Collections.unmodifiableList(elts);
```

The two implementations do not do the same thing!

- both avoid allowing clients to break the rep invariant
- both return a list containing the elements

```
But consider: xs = s.getElts();
s.insert('a');
xs.contains('a');
```

Version 2 is *observing* an exposed rep, leading to different behavior

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Ambiguity of "returns a list containing the current set elements"

- "returns a fresh mutable list containing the elements in the set *at the time of the call*" vs.
- "returns read-only access to a list that the ADT *continues to update to hold the current elements in the set*"

A third spec weaker than both [but less simple and useful!]

"returns a list containing the current set elements. *Behavior is unspecified (!) if* client attempts to mutate the list or to access the list after the set's elements are changed"

Also note: Version 2's spec also makes changing the rep later harder

- only "simple" to implement with rep as a List

Suggestions

Best options for implementing getElts()

- if O(n) time is acceptable for relevant use cases, copy the list
 - safest option
 - best option for changeability
- if O(1) time is required, then return an unmodifiable list
 - prevents breaking rep invariant
 - clearly document that behavior is unspecified after mutation
 - ideally, write your own unmodifiable view of the list that throws an exception on all operations after mutation
- if O(1) time is required and there is no unmodifiable version and you don't have time to write one, expose rep and feel guilty

Abstraction Functions

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Specifying an ADT

Different types of operations:

- 1. creators
- 2. observers
- 3. producers
- 4. mutators (if mutable)

Described in terms of how they change the **abstract state**

- abstract description of what the object means
 - difficult (unless concept is already familiar) but vital
- specs have no information about concrete representation
 - leaves us free to change those in the future

Connecting implementations to specs

For implementers / debuggers / maintainers of the implementation:

Representation Invariant: maps Object → boolean

- we saw this earlier!

Abstraction Function: maps Object \rightarrow abstract state

- says what the data structure *means* in vocabulary of the ADT
- maps the fields to the abstract state they represent
 - can check that the abstract value after each method meets the postcondition described in the specification

/** Represents a mutable circle in the plane. For example, * it can be a circle with center (0,0) and radius 1. */ public class Circle {

// Rep invariant: center != null and rad > 0
private Point center;
private double rad;

```
// Abstraction function:
// AF(this) = a circle with center at this.center
// and radius this.rad
```

// ... } /** Represents a mutable circle in the plane. For example, * it can be a circle with center (0,0) and radius 1. */ public class Circle {

```
// Rep invariant: center != null and edge != null
// and !center.equals(edge)
private Point center, edge;
```

```
// Abstraction function:
// AF(this) = a circle with center at this.center
// and radius this.center.distanceTo(this.edge)
```

```
// ...
}
```

/** An immutable polynomial with integer coefficients.

* Examples include 0, 2x, and x + 3x² + 5x. */
public class IntPoly {

```
// Rep invariant: coeffs != null
private final int[] coeffs;
```

```
// Abstraction function:
// AF(this) = sum of this.coeffs[i] * x^i
// for i = 0 .. this.coeffs.length
```

```
// ... coeff, degree, etc.
```

/** An immutable polynomial with integer coefficients.

* Examples include 0, 2x, and x + 3x² + 5x. */
public class IntPoly {

// Rep invariant: terms != null and // no two terms have the same degree and // terms is sorted in descending order by degree private final LinkedList<IntTerm> terms;

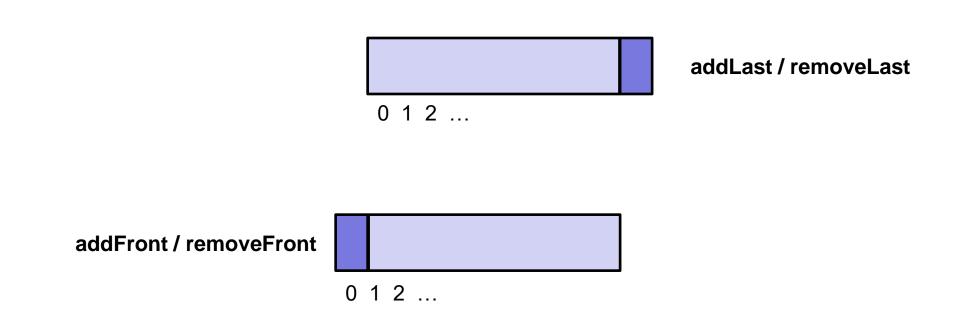
// Abstraction function:
// AF(this) = sum of monomials in this.terms

The abstraction function

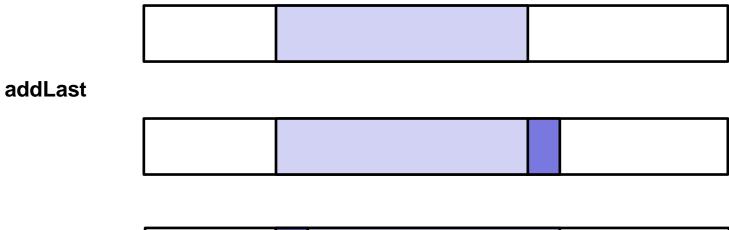
- Purely conceptual (not a Java function)
- Allows us to check correctness
 - use reasoning to show that the method leaves the abstract state such that it satisfies the postcondition

Example: IntDeque

// List that only allows insert/remove at ends.



// List that only allows insert/remove at ends.



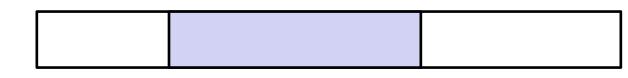


removeFront



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// List that only allows insert/remove at ends.



addLast + removeFront



addLast + removeFront

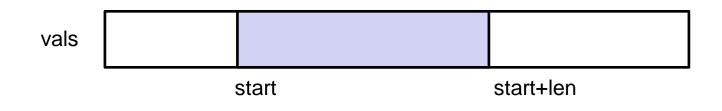


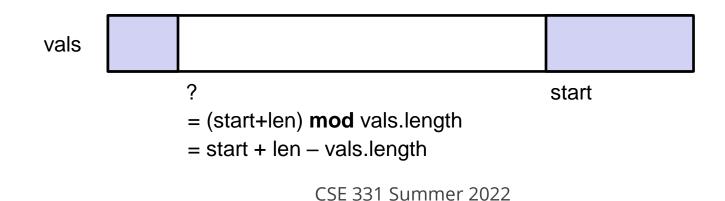
addLast + removeFront



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// List that only allows insert/remove at ends.





```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length</pre>
// vals[start..] + vals[0..?] otherwise
private int[] vals;
private int start, len;
// Creates an empty list.
public IntDeque() {
 vals = new int[3];
  start = len = 0;
                                           AF(this) = vals[0..-1] = []
}
```

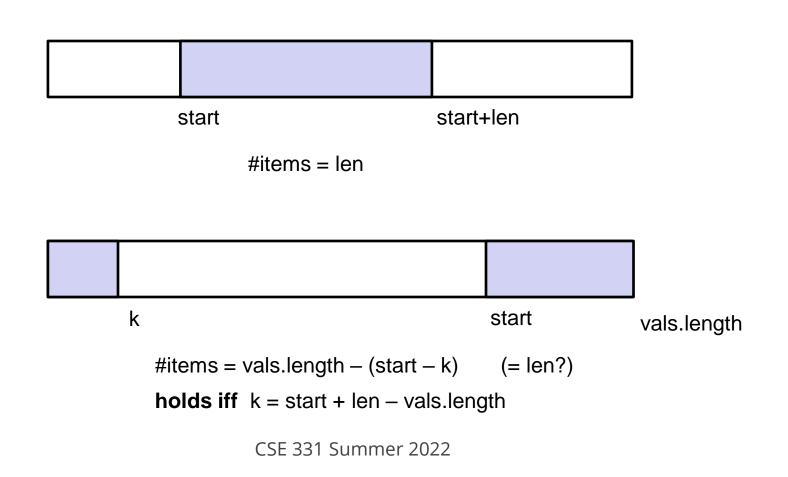
```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length
// vals[start..] + vals[0..?] otherwise
private int[] vals;
private int start, len;</pre>
```

// ...

```
// @returns length of the list
public int getLength() {
   return len;
}
```

// List that only allows insert/remove at ends.



```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length
// vals[start..] + vals[0..k] otherwise
private int[] vals;
private int start, len;</pre>
```

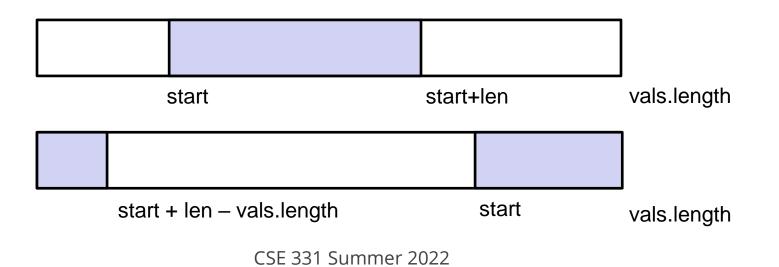
// ...

```
// @returns length of the list
public int getLength() {
   return len;
}
```

```
1 line of code
but 2 cases for reasoning
```

```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// @requires 0 <= i < length
// @returns this[i]
public int get(int i) { ... }</pre>
```



```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// @requires 0 <= i < length
// @returns this[i]
public int get(int i) {
    if (start + len <= vals.length) {
        return vals[start + i];
    } else {
        return vals[(start + i) % vals.length];
    }
}</pre>
```

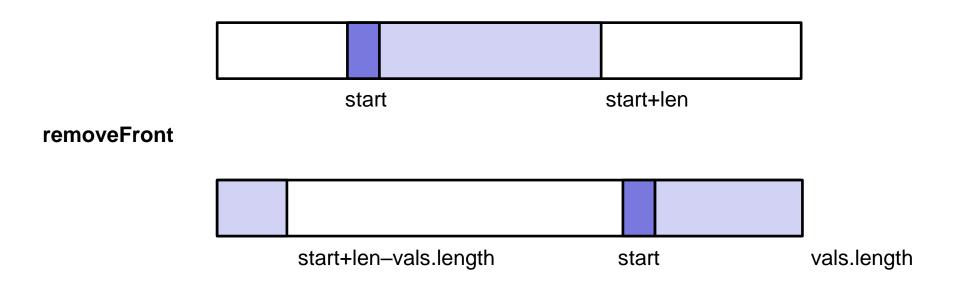
```
/** List that only allows insert/remove at ends. */
public class IntDeque {
```

```
// @requires 0 <= i < length
// @returns this[i]
public int get(int i) {
   return vals[(start + i) % vals.length];
}</pre>
```

/** List that only allows insert/remove at ends. */
public class IntDeque {

// @requires list length > 0
// @modifies this
// @effects first element of list is removed
// @returns value at the front of the list
public int removeFront() { ... }

// List that only allows insert/remove at ends.



```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length</pre>
// vals[start..] + vals[0..k]
                                   otherwise
// @requires list length > 0
// @modifies this
// @effects first element of list is removed
public void removeFront() {
  if (start + 1 < vals.length) {</pre>
    start += 1;
  } else {
    start = 0;
  }
  len -= 1;
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```

```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length</pre>
                                  otherwise
// vals[start..] + vals[0..k]
// @requires list length > 0
// @modifies this
// @effects first element of list is removed
public void removeFront() {
  start = (start + 1) % vals.length;
 len -= 1;
}
```

```
// AF(this) =
// vals[start..start+len-1] if start+len <= vals.length</pre>
// vals[start..] + vals[0..k]
                                  otherwise
// @requires list length > 0
// @modifies this
// @effects first element of list is removed
// @returns value at the front of the list
public int removeFront() {
  int val = get(0);
  start = (start + 1) % vals.length;
 len -= 1;
 return val;
}
```

Before next class...

- 1. Start on Prep. Quiz: HW4 as early as possible!
 - Reminds you about common set operations
 - E.g. union, intersection, complement
 - Think about some non-trivial cases needed for the homework
- 2. Section tomorrow will focus on HW4 preparation.

Extra: Abstract Interpretation

• Abstraction functions are good for much more (e.g. program analysis)

Extra: Testing

• What is testing? What makes something a good test case?