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# CSE 331

# Software Design & Implementation

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Lecture 2 – Reasoning About Straight-Line Code

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# Motivation for Reasoning

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- Want a way to determine correctness without running the code
- Most important part of the correctness techniques
  - tools, **inspection**, testing
- You need a way to do this in interviews
  - key reason why coding interviews are done without computers
- This is not easy (see HW0)

# Our Approach

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- We will learn a set of **formal tools** for proving correctness
  - (later, this will also allow us to generate the code)
- Most professionals can do reasoning like this in their head
  - most do an *informal* version of what we will see
  - eventually, it will be the same for you
- Formal version has key advantages
  - teachable
  - mechanical (no intuition or creativity required)
  - necessary for hard problems
    - we turn to formal tools when problems get too hard

# Formal Reasoning

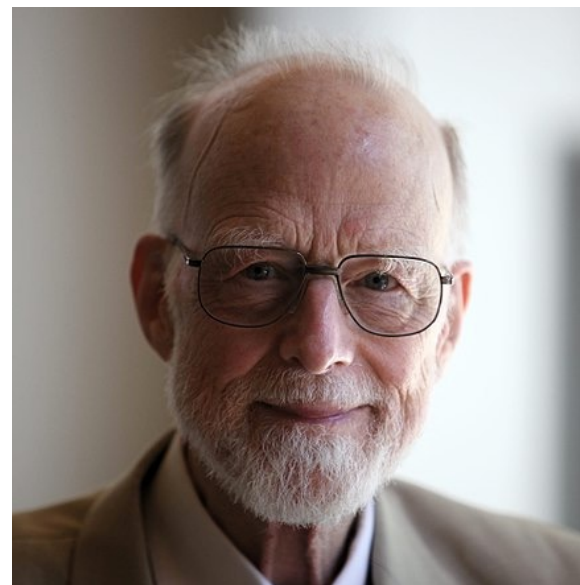
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- Invented by Robert Floyd and Sir Anthony Hoare
  - Floyd won the Turing award in 1978
  - Hoare won the Turing award in 1980



Robert Floyd

picture from [Wikipedia](#)



Tony Hoare

# Terminology of Floyd Logic

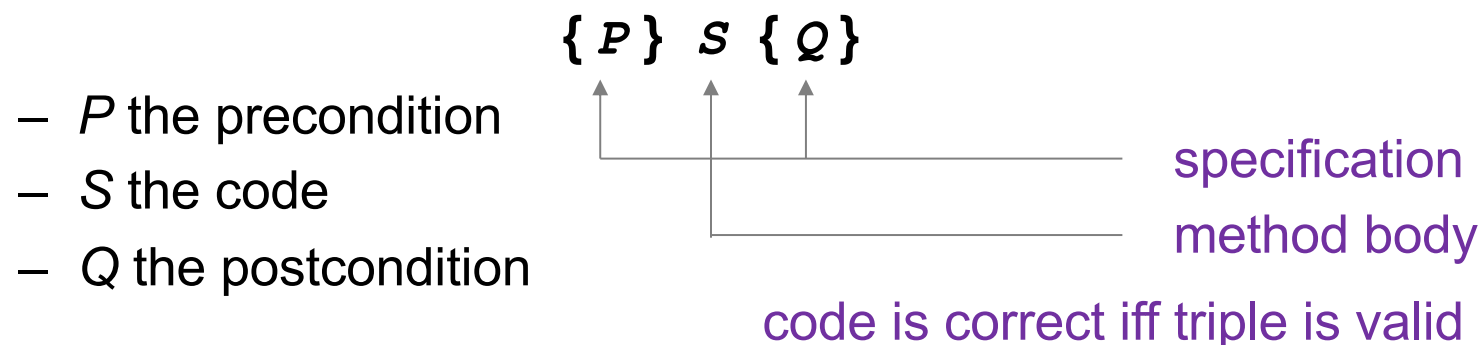
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- The *program state* is the values of all the (relevant) variables
- An *assertion* is a true / false claim (proposition) about the state at a given point during execution (e.g., on line 39)
- An assertion *holds* for a program state if the claim is true when the variables have those values
- An assertion before the code is a *precondition*
  - these represent assumptions about when that code is used
- An assertion after the code is a *postcondition*
  - these represent what we want the code to accomplish

# Hoare Triples

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- A **Hoare triple** is two assertions and one piece of code:



- A Hoare triple  $\{P\} S \{Q\}$  is called **valid** if:
  - in any state where  $P$  holds, executing  $S$  produces a state where  $Q$  holds
  - i.e., if  $P$  is true before  $S$ , then  $Q$  must be true after it
  - otherwise, the triple is called **invalid**

# Notation

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- Floyd logic writes assertions in `{..}`
  - since Java code also has `{..}`, I will use `{{...}}`
  - e.g., `{{ w >= 1 }} x = 2 * w; {{ x >= 2 }}`
- Assertions are math / logic not Java
  - you can use the usual math notation
    - (e.g., `=` instead of `==` for equals)
  - purpose is communication with other humans (not computers)
  - we will need **and**, **or**, **not** as well
    - can also write use  $\wedge$  (and)  $\vee$  (or) etc.
- The Java language also has assertions (**assert** statements)
  - throws an exception if the condition does not evaluate true
  - we will discuss these more later in the course

# Example 1

---

Is the following Hoare triple valid or invalid?

- assume all variables are integers and there is no overflow

$\{\{ x \neq 0 \}\} y = x * x; \{\{ y > 0 \}\}$



# Example 1

---

Is the following Hoare triple valid or invalid?

- assume all variables are integers and there is no overflow

$\{\{ \mathbf{x} \neq 0 \}\} \mathbf{y} = \mathbf{x} * \mathbf{x}; \{\{ \mathbf{y} > 0 \}\}$

Valid

- $\mathbf{y}$  could only be zero if  $\mathbf{x}$  were zero (which it isn't)

## Example 2

---

Is the following Hoare triple valid or invalid?

- assume all variables are integers and there is no overflow

$\{\{ z \neq 1 \}\} y = z * z; \{\{ y \neq z \}\}$

## Example 2

---

Is the following Hoare triple valid or invalid?

- assume all variables are integers and there is no overflow

$\{\{ z \neq 1 \}\} y = z * z; \{\{ y \neq z \}\}$

Invalid

- counterexample:  $z = 0$

# Checking Validity

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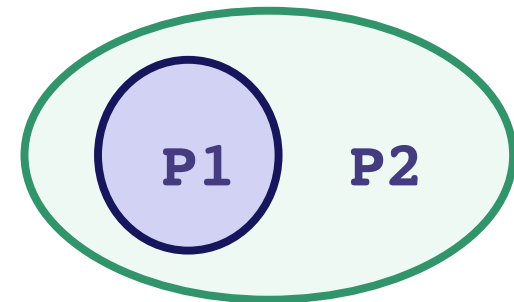
- So far: decided if a Hoare triple is valid by ... **hard** thinking
- Soon: mechanical process for reasoning about
  - assignment statements
  - conditionals
  - [next lecture] loops
  - (all code can be understood in terms of those 3 elements)
- Can use those to check correctness in a “turn the crank” manner
- Next: a way to compare different assertions
  - useful, e.g., to compare possible preconditions

# Weaker vs. Stronger Assertions

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If P1 implies P2 (written  $P1 \Rightarrow P2$ ), then:

- P1 is **stronger** than P2
- P2 is **weaker** than P1



Whenever P1 holds, P2 also holds

- So it is more (or at least as) “difficult” to satisfy P1
  - the program states where P1 holds are a subset of the program states where P2 holds
- So P1 puts more constraints on program states
- So it is a stronger set of requirements on the program state
  - P1 gives you more information about the state than P2

# Examples

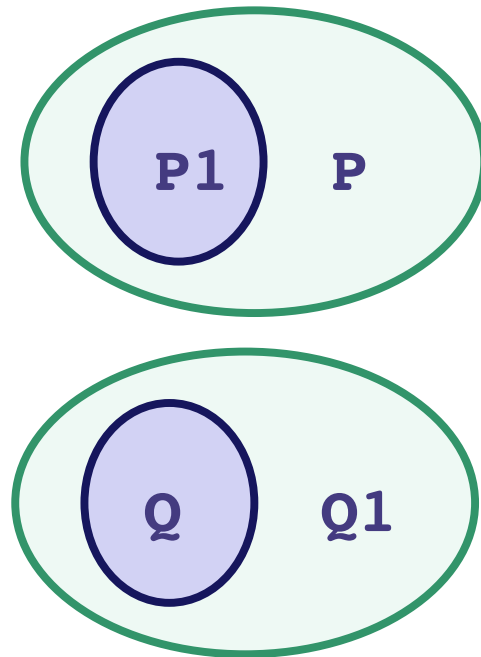
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- $x = 17$  is stronger than  $x > 0$
- $x$  is prime is neither stronger nor weaker than  $x$  is odd
- $x$  is prime and  $x > 2$  is stronger than  $x$  is odd

# Floyd Logic Facts

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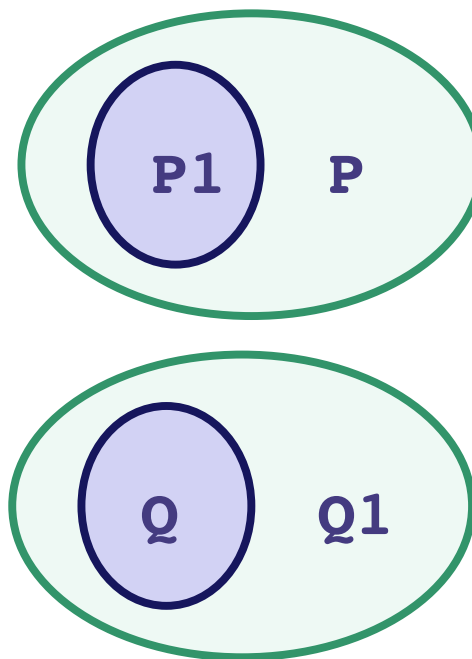
- Suppose  $\{P\} S \{Q\}$  is valid.
- If  $P1$  is stronger than  $P$ , then  $\{P1\} S \{Q\}$  is valid.
- If  $Q1$  is weaker than  $Q$ , then  $\{P\} S \{Q1\}$  is valid.
- Example:
  - Suppose  $P$  is  $x \geq 0$  and  $P1$  is  $x > 0$
  - Suppose  $Q$  is  $y > 0$  and  $Q1$  is  $y \geq 0$
  - Since  $\{x \geq 0\} y = x+1 \{y > 0\}$  is valid,  $\{x > 0\} y = x+1 \{y \geq 0\}$  is also valid



# Floyd Logic Facts

---

- Suppose  $\{P\} S \{Q\}$  is valid.
- If  $P1$  is stronger than  $P$ , then  $\{P1\} S \{Q\}$  is valid.
- If  $Q1$  is weaker than  $Q$ , then  $\{P\} S \{Q1\}$  is valid.
- **Key points:**
  - always okay to **strengthen** a **precondition**
  - always okay to **weaken** a **postcondition**

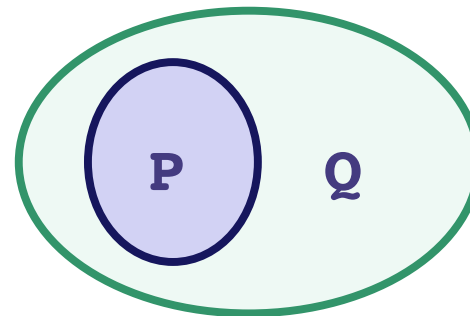




# Floyd Logic Facts

---

- When is  $\{P\} ; \{Q\}$  is valid?
  - with no code in between
- Valid if any state satisfying P also satisfies Q
- I.e., if P is **stronger** than Q



# Forward & Backward Reasoning

# Example of Forward Reasoning

---

Work forward from the precondition

`{{ w > 0 }}`

`x = 17;`

`{{ _____ }}`

`y = 42;`

`{{ _____ }}`

`z = w + x + y;`

`{{ _____ }}`

# Example of Forward Reasoning

---

Work forward from the precondition

$\{\{ w > 0 \}\}$

$x = 17;$

$\{\{ w > 0 \text{ and } x = 17 \}\}$

$y = 42;$

$\{\{ \text{_____} \}\}$

$z = w + x + y;$

$\{\{ \text{_____} \}\}$

# Example of Forward Reasoning

---

Work forward from the precondition

$\{\{ w > 0 \}\}$

$\mathbf{x} = 17;$

$\{\{ w > 0 \text{ and } x = 17 \}\}$

$\mathbf{y} = 42;$

$\{\{ w > 0 \text{ and } x = 17 \text{ and } y = 42 \}\}$

$\mathbf{z} = \mathbf{w} + \mathbf{x} + \mathbf{y};$

$\{\{ \underline{\hspace{10em}} \}\}$

# Example of Forward Reasoning

---

Work forward from the precondition

$\{\{ w > 0 \}\}$

$\mathbf{x} = 17;$

$\{\{ w > 0 \text{ and } x = 17 \}\}$

$\mathbf{y} = 42;$

$\{\{ w > 0 \text{ and } x = 17 \text{ and } y = 42 \}\}$

$\mathbf{z} = \mathbf{w} + \mathbf{x} + \mathbf{y};$

$\{\{ w > 0 \text{ and } x = 17 \text{ and } y = 42 \text{ and } z = w + x + y \}\}$



# Example of Forward Reasoning

---

Work forward from the precondition

$\{\{ w > 0 \}\}$

$\mathbf{x = 17;}$

$\{\{ w > 0 \text{ and } x = 17 \}\}$

$\mathbf{y = 42;}$

$\{\{ w > 0 \text{ and } x = 17 \text{ and } y = 42 \}\}$

$\mathbf{z = w + x + y;}$

$\{\{ w > 0 \text{ and } x = 17 \text{ and } y = 42 \text{ and } z = w + 59 \}\}$

# Forward Reasoning

---

- Start with the **given** precondition
- Fill in the **strongest** postcondition
- For an assignment,  $\mathbf{x} = \mathbf{y}$ ...
  - add the fact “ $x = y$ ” to what is known
  - important subtleties here... (more on those later)
- Later: if statements and loops...



# Example of Backward Reasoning

---

Work backward from the desired postcondition

`{{ _____ }}`

`x = 17;`

`{{ _____ }}`

`y = 42;`

`{{ _____ }}`

`z = w + x + y;`

`{{ z < 0 }}`

# Example of Backward Reasoning

---

Work backward from the desired postcondition

`{{ _____ }}`

`x = 17;`

`{{ _____ }}`

`y = 42;`

`{{ w + x + y < 0 }}`

`z = w + x + y;`

`{{ z < 0 }}`



# Example of Backward Reasoning

---

Work backward from the desired postcondition

$\{ \underline{\hspace{10em}} \}$

$\mathbf{x} = 17;$

$\{ \{ w + x + 42 < 0 \} \}$

$\mathbf{y} = 42;$

$\{ \{ w + x + y < 0 \} \}$

$\mathbf{z} = \mathbf{w} + \mathbf{x} + \mathbf{y};$

$\{ \{ z < 0 \} \}$



# Example of Backward Reasoning

---

Work backward from the desired postcondition

$\{\{ w + 17 + 42 < 0 \}\}$

$\mathbf{x} = 17;$

$\{\{ w + x + 42 < 0 \}\}$

$\mathbf{y} = 42;$

$\{\{ w + x + y < 0 \}\}$

$\mathbf{z} = \mathbf{w} + \mathbf{x} + \mathbf{y};$

$\{\{ z < 0 \}\}$

# Backward Reasoning

---

- Start with the **required** postcondition
- Fill in the **weakest** precondition
- For an assignment,  $\mathbf{x = y}$ :
  - just replace “x” with “y” in the postcondition
  - if the condition using “y” holds beforehand, then the condition with “x” will afterward since  $x = y$  then
- Later: if statements and loops...

# Correctness by Forward Reasoning

---

Use forward reasoning to determine if this code is correct:

`{{ w > 0 }}`

`x = 17;`

`y = 42;`

`z = w + x + y;`

`{{ z > 50 }}`

# Example of Forward Reasoning

---

$\{\{ w > 0 \}\}$

$\mathbf{x = 17;}$

$\{\{ w > 0 \text{ and } x=17 \}\}$

$\mathbf{y = 42;}$

$\{\{ w > 0 \text{ and } x=17 \text{ and } y=42 \}\}$

$\mathbf{z = w + x + y;}$

$\{\{ w > 0 \text{ and } x=17 \text{ and } y=42 \text{ and } z = w + 59 \}\}$

$\{\{ z > 50 \}\}$

Do the facts that are always true  
imply the facts we need?

I.e., is the bottom statement  
**weaker** than the top one?

(Recall that weakening the postcondition is always okay.)

# Correctness by Backward Reasoning

---

Use backward reasoning to determine if this code is correct:

```
{{ w < -60 }}
```

```
  x = 17;
```

```
  y = 42;
```

```
  z = w + x + y;
```

```
{{ z < 0 }}
```



# Correctness by Backward Reasoning

---

Use backward reasoning to determine if this code is correct:

$\{\{ w < -60 \}\}$

$\{\{ w + 17 + 42 < 0 \}\} \Leftrightarrow \{\{ w < -59 \}\}$

$\mathbf{x = 17;}$

$\{\{ w + x + 42 < 0 \}\}$

$\mathbf{y = 42;}$

$\{\{ w + x + y < 0 \}\}$

$\mathbf{z = w + x + y;}$

$\{\{ z < 0 \}\}$

Do the facts that are always true imply the facts we need?

I.e., is the top statement **stronger** than the bottom one?

(Recall that strengthening the precondition is always okay.)

# Combining Forward & Backward

---

It is okay to use both types of reasoning

- Reason forward from precondition
- Reason backward from postcondition

Will meet in the middle:

**{{ P }}**

**S1**

**S2**

**{{ Q }}**

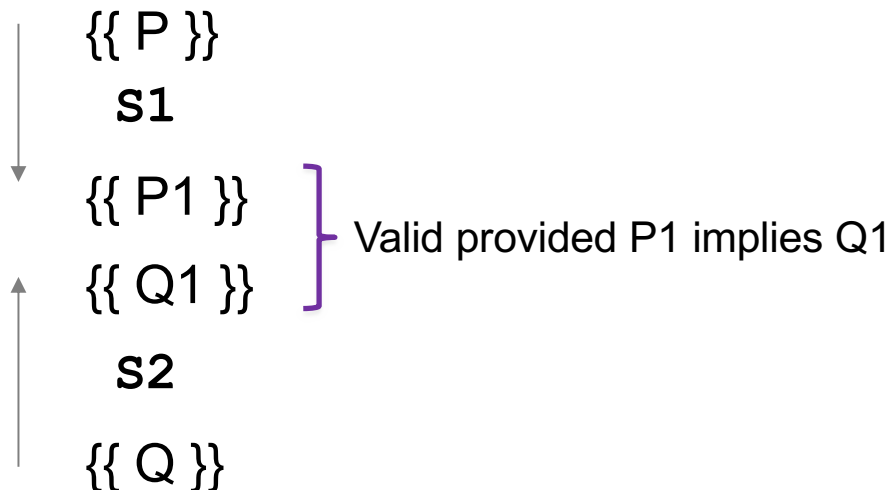
# Combining Forward & Backward

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It is okay to use both types of reasoning

- Reason forward from precondition
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Will meet in the middle:



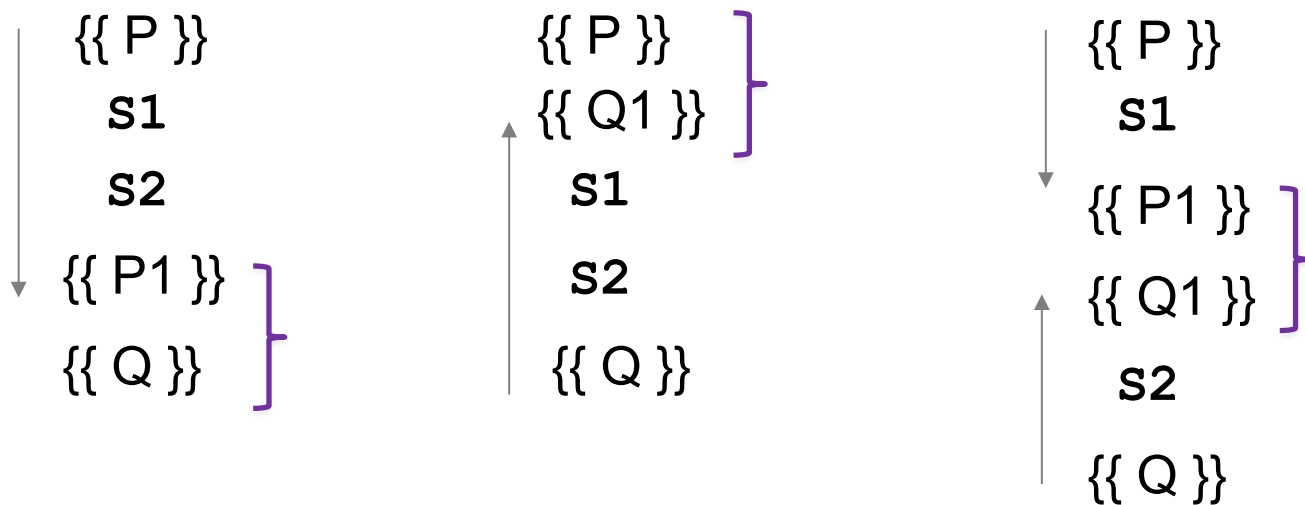
# Combining Forward & Backward

---

Reasoning in either direction gives valid assertions

Just need to check adjacent assertions:

- top assertion must imply bottom one



# Subtleties in Forward Reasoning...

---

- Forward reasoning can **fail** if applied blindly...

```

  {{ }}
  w = x + y;
  {{ w = x + y }}
  x = 4;
  {{ w = x + y and x = 4 }}
  y = 3;
  {{ w = x + y and x = 4 and y = 3 }}

```

This implies that  $w = 7$ , but that is not true!

- $w$  equals whatever  $x + y$  was **before** they were changed

# Fix 1

---

- Use **subscripts** to refer to old values of the variables
- Un-subscripted variables should always mean **current** value

`{{}}`

`w = x + y;`

`{{ w = x + y }}`

`x = 4;`

`{{ w = x1 + y and x = 4 }}`

`y = 3;`

`{{ w = x1 + y1 and x = 4 and y = 3 }}`

## Fix 2 (better)

---

- Express prior values in terms of the current value

$\{\{\}$

$w = x + y;$

$\{\{ w = x + y \}\}$

$x = x + 4;$

$\{\{ w = x_1 + y \text{ and } x = x_1 + 4 \}\}$  Now,  $x_1 = x - 4$

$\Rightarrow \{\{ w = x - 4 + y \}\}$  So  $w = x_1 + y \Leftrightarrow w = x - 4 + y$

Note for updating variables, e.g.,  $x = x + 4$ :

- Backward reasoning just substitutes new value (no change)
- Forward reasoning requires you to invert the “+” operation

# Forward vs. Backward

---

- Forward reasoning:
  - Find strongest postcondition
  - Intuitive: “simulate” the code in your head
    - BUT you need to change facts to refer to *prior values*
  - Inefficient: Introduces many irrelevant facts
    - usually need to weaken as you go to keep things sane
- Backward reasoning
  - Find weakest precondition
  - Formally simpler
  - Efficient
  - (Initially) unintuitive



# If Statements

# If Statements

---

Forward reasoning

```
{{ P }}  
if (cond)  
  S1  
else  
  S2  
{{ ? }}
```

# If Statements

---

Forward reasoning

```
  {{ P }}  
  if (cond)  
  → {{ P and cond }}  
    S1  
  else  
  → {{ P and not cond }}  
    S2  
  {{ ? }}
```




# If Statements

---

Forward reasoning

```
  {{ P }}  
  if (cond)  
    {{ P and cond }}  
    S1  
  {{ P1 }}  
  else  
    {{ P and not cond }}  
    S2  
  {{ P2 }}  
  {{ P1 or P2 }}
```



# If Statements

---

Backward reasoning

```
{{ ? }}  
if (cond)  
  S1  
else  
  S2  
{{ Q }}
```

# If Statements

---

Backward reasoning

```
  {{ ? }}  
  if (cond)  
    S1  
  → {{ Q }}  
  else  
    S2  
  → {{ Q }}  
  {{ Q }}
```

# If Statements

---

Backward reasoning

```
  {{ ? }}  
  if (cond)  
    ↑ {{ Q1 }}  
    S1  
    ↑ {{ Q }}  
  else  
    ↑ {{ Q2 }}  
    S2  
    ↑ {{ Q }}  
  {{ Q }}
```



# If Statements

---

Backward reasoning

`{{ cond and Q1 or`

`not cond and Q2 }}`

`if (cond)`

`{{ Q1 }}`

`S1`

`{{ Q }}`

`else`

`{{ Q2 }}`

`S2`

`{{ Q }}`

`{{ Q }}`

# If-Statement Example

---

Forward reasoning

```
{  
}  
if (x >= 0)  
    y = x;  
else  
    y = -x;  
{ ? }
```

# If-Statement Example

---

Forward reasoning

```
  {{ }}  
  if (x >= 0)  
  → {{ x >= 0 }}  
    y = x;  
  else  
  → {{ x < 0 }}  
    y = -x;  
  {{ ? }}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {} x >= 0 {}  
  y = x;  
  ↓ {} x >= 0 and y = x {}  
else  
  {} x < 0 {}  
  y = -x;  
  ↓ {} x < 0 and y = -x {}  
{} ? {}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{(x >= 0 and y = x) or  
(x < 0 and y = -x)}
```

# If-Statement Example

---

Forward reasoning

```
{  
}  
if (x >= 0)  
  {  
    x >= 0  
    y = x;  
    {  
      x >= 0 and y = x  
    }  
  }  
else  
  {  
    x < 0  
    y = -x;  
    {  
      x < 0 and y = -x  
    }  
  }  
{  
  y = |x|  
}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {} x >= 0 {}  
  y = x;  
  {} x >= 0 and y = x {}  
else  
  {} x < 0 {}  
  y = -x;  
  {} x < 0 and y = -x {}  
{} y = |x| {}
```

**Warning:** many write `{} y >= 0 {}` here

That is true but it is *strictly* weaker.  
(It includes cases where  $y \neq x$ )

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{?}  
if (x >= 0)  
  y = x;  
else  
  y = -x;  
{y = |x|}
```



# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{?}  
if (x >= 0)  
  y = x;  
  → {y = |x|}  
else  
  y = -x;  
  → {y = |x|}  
{y = |x|}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{?}  
if (x >= 0)  
  ↑ {x = |x|}  
  y = x;  
  {y = |x|}  
else  
  ↑ {-x = |x|}  
  y = -x;  
  {y = |x|}  
{y = |x|}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{?}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {y = |x|}  
else  
  {x <= 0}  
  y = -x;  
  {y = |x|}  
{y = |x|}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{(x >= 0 and x >= 0) or  
(x < 0 and x <= 0)}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {y = |x|}  
else  
  {x <= 0}  
  y = -x;  
  {y = |x|}  
{y = |x|}
```

# If-Statement Example

---

Forward reasoning

```
{}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{x >= 0 or x < 0}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {y = |x|}  
else  
  {x <= 0}  
  y = -x;  
  {y = |x|}  
{y = |x|}
```

# If-Statement Example

---

Forward reasoning

```
{  
}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {x >= 0 and y = x}  
else  
  {x < 0}  
  y = -x;  
  {x < 0 and y = -x}  
{y = |x|}
```

Backward reasoning

```
{  
}  
if (x >= 0)  
  {x >= 0}  
  y = x;  
  {y = |x|}  
else  
  {x <= 0}  
  y = -x;  
  {y = |x|}  
{y = |x|}
```

Next time: Loops...