CSE 331 Software Design & Implementation

Spring 2022 Section 1 – Code Reasoning

Administrivia

- HW1 due next Wednesday.
 - Last two questions have loops: covered tomorrow in lecture.
- Any questions before we dive in?
 - What are the most interesting/confusing/puzzling things so far in the course?

Agenda

- Introductions?
- Review logical reasoning about code with Floyd Logic
- Practice both forward and backward modes
 - Just assignment, conditional ("if-then-else"), and sequence
 - Logical rules from yesterday's lecture/notes
- Review logical strength of assertions (weaker vs. stronger)
- Practice determining stronger/weaker assertions

Why reason about code?

- Prove that code is correct
- Understand why code is correct
- Diagnose why/how code is not correct
- Specify code behavior

Logical reasoning about code

- Determine facts that hold of program state between statements
 - "Fact" ~ assertion (logical formula over program state, informally "value(s) of some/all program variables)
 - Driven by assumption (precondition) or goal (postconditon)
- Forward reasoning
 - What facts follow from initial assumptions?
 - Go from <u>pre</u>condition to <u>post</u>condition
- Backward reasoning
 - What facts need to be true to reach a goal?
 - Go from <u>post</u>condition to <u>pre</u>condition

Hoare Logic: Validity by Reasoning

- Checking validity of {{P}} s {{Q}}
 - Valid iff, starting from any state satisfying P, executing S results in a state satisfying Q
- Forward reasoning:
 - Reason from P to strongest postcondition {{P}} S {{R}}
 - Check that R implies Q (i.e., Q is weaker)
- Backward reasoning:
 - Reason from Q to get weakest precondition {{R}} S {{Q}}
 - Check that P implies R (i.e., P is stronger)

Implication (=>)

- Logic formulas with and (&, &&, or ∧), or (|, ||, or ∨) and not
 (! or ¬) have the same meaning they do in programs
- Implication might be a bit new, but the basic idea is pretty simple. Implication p=>q is true as long as q is always true whenever p is

р	q	p => q
Т	Т	Т
Т	F	F
F	Т	Т
F	F	Т

Assignment Statements

- Reasoning about x = y;
- Forward reasoning:
 - add "x = y" as a new fact
 - (also rewrite any existing references to "x" to use new value)
- Backward reasoning:
 - replace all instances of "x" in the postcondition with "y"

Conditionals, more closely

Forward reasoning

```
\{\{P\}\}\}
if (b)
\{\{P \land b\}\}\}
S_1
\{\{Q_1\}\}\}
else
\{\{P \land !b\}\}\}
S_2
\{\{Q_2\}\}\}
\{\{Q_1 \lor Q_2\}\}
```

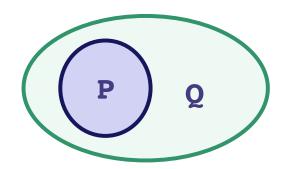
Backward reasoning

```
 \left\{ \left\{ \left( b \land P_{1} \right) \lor \left( ! b \land P_{2} \right) \right\} \right\}  if (b)  \left\{ \left\{ P_{1} \right\} \right\}   S_{1}   \left\{ \left\{ \mathcal{Q} \right\} \right\}  else  \left\{ \left\{ P_{2} \right\} \right\}   S_{2}   \left\{ \left\{ \mathcal{Q} \right\} \right\}   \left\{ \left\{ \mathcal{Q} \right\} \right\}
```

Weaker vs. stronger

Formal definition:

- If $P \Rightarrow Q$, then
 - Q is weaker than P
 - P is stronger than Q



Intuitive definition:

- "Weak" means unrestrictive; a weaker assertion has a larger set of possible program states (e.g., x != 0)
- "Strong" means restrictive; a stronger assertion has a smaller set of possible program states (e.g., x = 1 or x > 0 are both stronger than x != 0).

Worksheet

- Take ~10 minutes to get where you can
- Find a partner and work with them
- Let me know if you feel stuck
- We'll walk through some solutions afterwards

```
{{ true }}
if (x>0) {
  \{\{ x > 0 \}\}
  y = 2*x;
  \{\{ \mathbf{x} > 0 \land \mathbf{y} = 2\mathbf{x} \}\}
} else {
  \{\{ x \le 0 \}\}
  y = -2*x;
  \{\{ x \le 0 \land y = -2x \}\}
\{\{(x > 0 \land y = 2x) \lor (x \le 0 \land y = -2x) \}\}
\Rightarrow \{\{y = 2|x|\}\}
```

```
\{\{ y > 15 \lor (y \le 5 \land y + z > 17) \}\}
if (y > 5) {
  {\{\{y > 15\}\}}
  x = y + 2
  \{\{ x > 17 \}\}
} else {
  \{\{y + z > 17\}\}
  x = y + z;
  \{\{ x > 17 \}\}
\{\{ x > 17 \}\}
```

Worksheet – problem 6 (forward)

```
{{ true }}
if (x < y) {
  \{\{ true \land x < y \} \}
  m = x;
  \{\{ x < y \land m = x \}\}
} else {
  \{\{ true \land x >= y \}\}
  m = y;
  \{\{ x >= y \land m = y \}\}
\{\{ (x < y \land m = x) \lor (x >= y \land m = y) \}\}
\Rightarrow \{\{ m = \min(x, y) \}\}
```

Worksheet – problem 6 (backward)

```
{{ true }} ⇔
\{\{ (x \le y \land x \le y) \lor (y \le x \land x \ge y) \}\}
if (x < y) {
   \{\{ \mathbf{x} = \min(\mathbf{x}, \mathbf{y}) \}\} \Leftrightarrow \{\{ \mathbf{x} \leq \mathbf{y} \}\}
   m = x;
   \{\{ m = min(x, y) \}\}
} else {
   \{\{ y = \min(x, y) \}\} \Leftrightarrow \{\{ x >= y \}\}
   m = y;
   \{\{ m = \min(x, y) \}\}
\{\{ m = \min(x, y) \}\}
```

```
{{ y > 23 }} is stronger than {{ y >= 23 }}

{{ y = 23 }} is stronger than {{ y >= 23 }}

{{ y < 0.23 }} is weaker than {{ y < 0.00023 }}

{{ x = y * z }} is incomparable with {{ y = x / z }}

{{ is_prime(y) }} is incomparable with {{ is_odd(y) }}</pre>
```

Questions?

- What is the most surprising thing about this?
- What is the most confusing thing?
- What will need a bit more thinking to digest?