

CSE 331: Software Design & Implementation

Section 1 – Code Reasoning

For these problems, assume that all numbers are integers, that integer overflow will never occur, and that division is truncating integer division (like in Java).

1. Fill in the blanks using **forward** reasoning.

```
{ x >= 0 ∧ y >= 0 }  
y = 16;  
  
{ x >= 0 ∧ y = 16 }  
x = x + y;  
  
{ x >= 16 ∧ y = 16 }  
x = sqrt(x);  
  
{ x >= 4 ∧ y = 16 }  
y = y - x;  
  
{ x >= 4 ∧ y = 16 - x }  
  
⇒ { x >= 4 ∧ y <= 12 }
```

2. Fill in the blanks using **forward** reasoning.

```
{ true }  
if (x > 0) {  
    { x > 0 }  
    abs = x;  
    { x > 0 ∧ abs = x }  
} else {  
    { x <= 0 }  
    abs = -x;  
    { x <= 0 ∧ abs = -x }  
}  
  
{ (x > 0 ∧ abs = x) ∨ (x <= 0 ∧ abs = -x) }  
  
⇒ { abs = |x| }
```

3. Fill in the blanks using **backward** reasoning.

```
{  $x + 3 * b - 4 > 0$  }  
a = x + b;  
  
{  $a + 2 * b - 4 > 0$  }  
c = 2 * b - 4;  
  
{  $a + c > 0$  }  
x = a + c;  
  
{  $x > 0$  }
```

4. Fill in the blanks using **backward** reasoning.

```
{  $y > 15 \vee (y \leq 5 \wedge y + z > 17)$  }  
if (y > 5) {  
    {  $y > 15$  }  
    x = y + 2  
    {  $x > 17$  }  
} else {  
    {  $y + z > 17$  }  
    x = y + z;  
    {  $x > 17$  }  
}  
{  $x > 17$  }
```

5. Fill in the blanks using **backward** reasoning (extra practice problems).

a.

```
{  $x - 1 \neq 0$  }  
x = x - 2;  
  
{  $x + 1 \neq 0$  }  
z = x + 1;  
  
{  $z \neq 0$  }
```

b.

```
{  $3 * y > 0$  }  
x = 2 * y;  
  
{  $x + y > 0$  }  
z = x + y;  
  
{ z > 0 }
```

c.

```
{  $v + 1 < 0 \vee -2 * w < 0$  }  
w = 2 * w;  
  
{  $v + 1 < 0 \vee -w < 0$  }  
z = -w;  
  
{  $v + 1 < 0 \vee z < 0$  }  
y = v + 1;  
  
{  $y < 0 \vee z < 0$  }  
x = min(y, z);  
  
{ x < 0 }
```

6. Prove that the following code computes the minimum.

This can be done using either forward or backward reasoning. Or a combination!

Forward reasoning:

```
{ true }  
if (x < y) {  
  
    {  $true \wedge x < y$  }  
    m = x;  
  
    {  $x < y \wedge m = x$  }  
} else {  
  
    {  $true \wedge x \geq y$  }  
    m = y;  
  
    {  $x \geq y \wedge m = y$  }  
}  
{  $(x < y \wedge m = x) \vee (x \geq y \wedge m = y)$  }  
  
{ m = min(x, y) }
```

Backward reasoning:

```
{ (x <= y /\ x < y) \/ (y <= x /\ x >= y) } ⇒ { true }
if (x < y) {

    { x <= y }
    m = x;

    { m = min(x, y) }
} else {

    { y <= x }
    m = y;

    { m = min(x, y) }
}
{ same }

{ m = min(x, y) }
```

7. For each pair of logical assertions, circle the **stronger** logical assertion (or indicate that the pair is **incomparable**).

- | | |
|---------------------------|---------------------------------------|
| a. { y > 23 } | { y >= 23 } |
| b. { y = 23 } | { y >= 23 } |
| c. { y < 0.23 } | { y < 0.00023 } |
| d. { x = y * z } | { y = x / z } (incomparable) |
| e. { is_prime(y) } | { is_odd(y) } (incomparable) |

7d is fairly non-trivial, requiring close examination of possible program states. Also, remember that variables are always integers, and a forward slash is always truncating integer division unless otherwise stated.

To see that the assertions in 7d are incomparable, consider the following program states: x is 5, z is 2, and y is 2; and x is 0, z is 0, and y is 0.