CSE 331 Software Design & Implementation

Kevin Zatloukal Spring 2020 Lecture 4½ – An Interview Question

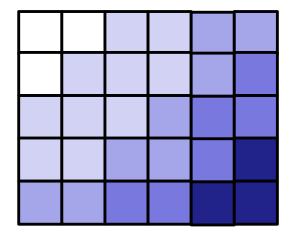
Sorted Matrix Search

Problem Description

Given a matrix M (of size m x n), where every row and every column is sorted, find out whether a given number x is in the matrix.

Sorted Matrix Search

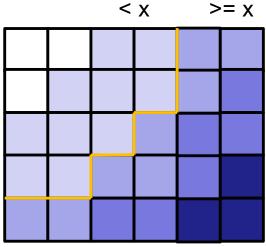
Given a sorted matrix M (of size $m \ge n$), where every row and every column is sorted, find out whether a given number x is in the matrix.



(darker color means larger)

Sorted Matrix Search

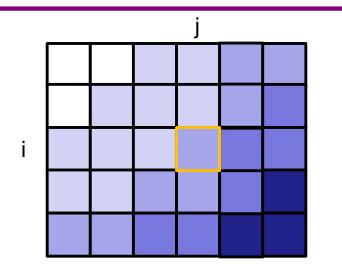
Given a sorted matrix M (of size m x n), where every row and every column is sorted, find out whether a given number x is in the matrix.





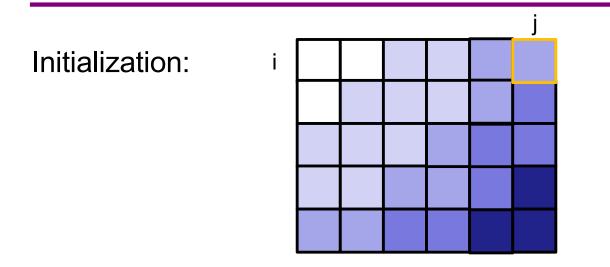
(darker color means larger)

(One) **Idea**: Trace the contour between the numbers $\leq x$ and > xin each row to see if x appears.



Partial Invariant: M[i,0], ..., M[i,j-1] < x ≤ M[i,j], ..., M[i,n-1]

- for each i, holds for exactly one j
- holds when we are in the right spot in row i

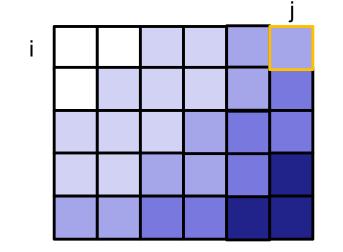


Partial Invariant: M[i,0], ..., M[i,j-1] < x ≤ M[i,j], ..., M[i,n-1]

How do we get the invariant to hold with i = 0?

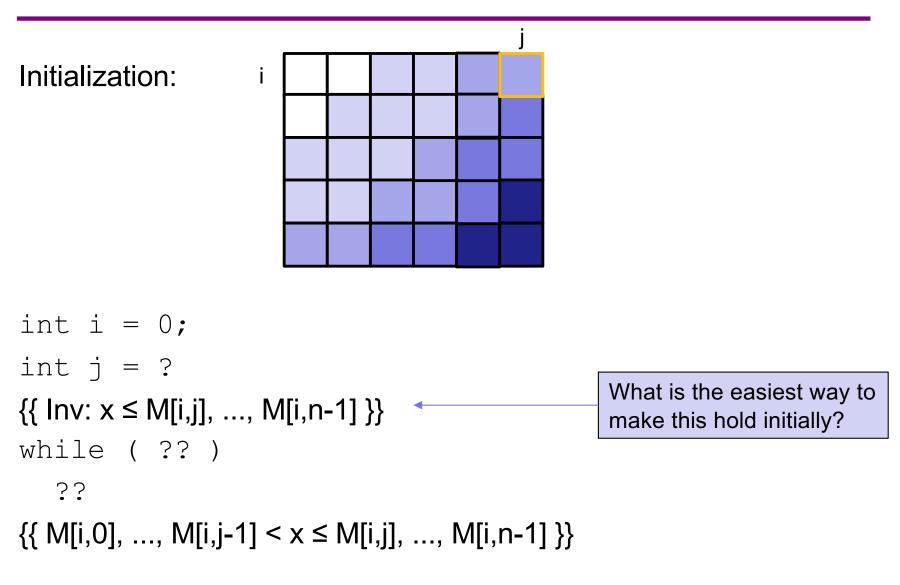
- no easy way to initialize it so the invariant holds
- we need to search...

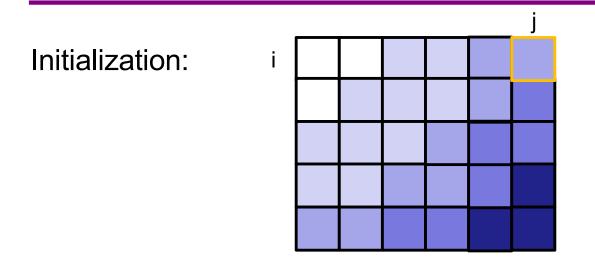
Initialization:

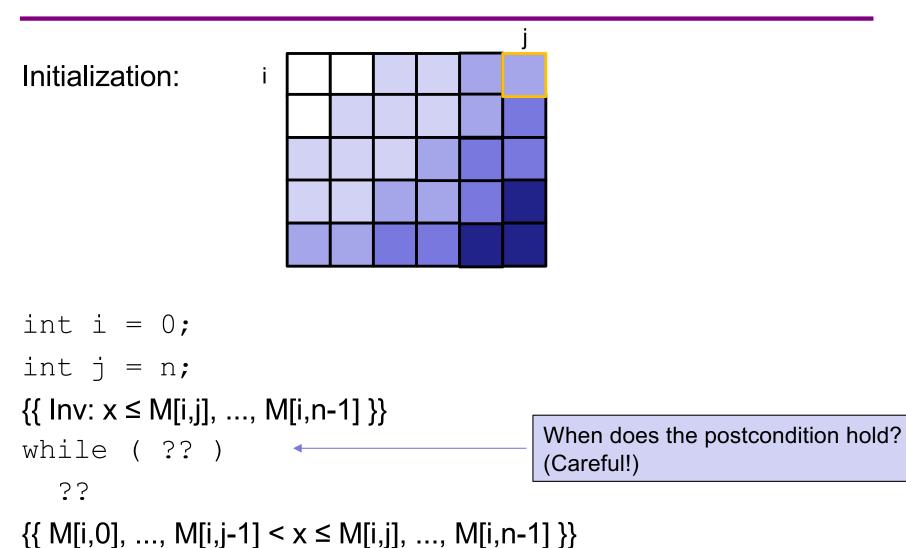


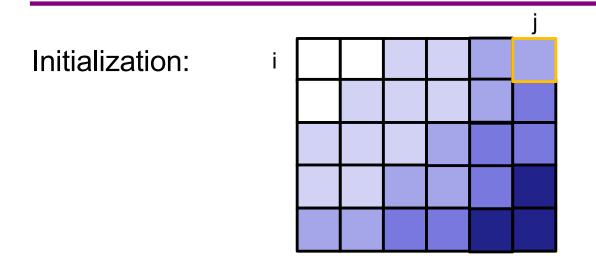
New goal: M[0,0], ..., M[0,j-1] < $x \le M[0,j]$, ..., M[0,n-1]

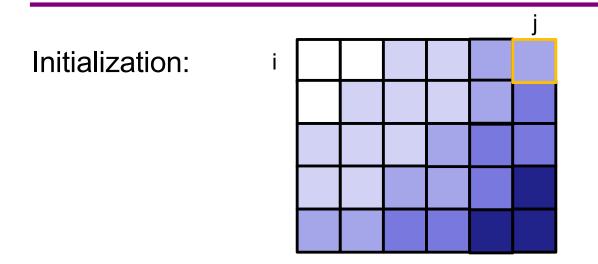
- will need a loop to find j
- Loop invariant: x ≤ M[0,j], ..., M[0,n-1]
 - weakening of the new goal
 - decrease j until we get M[0,j-1] to also hold



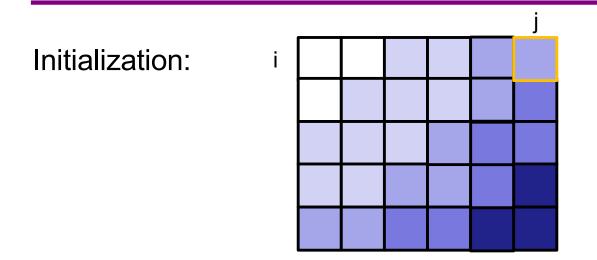




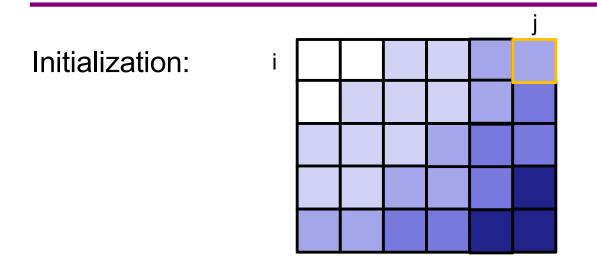




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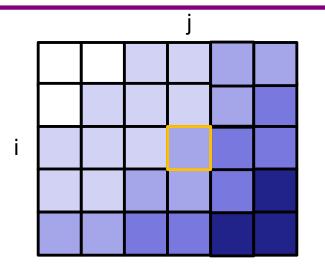


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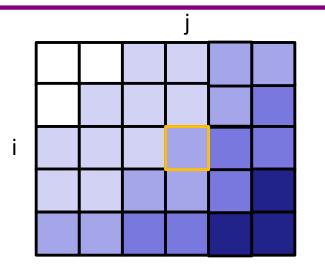
Initialization: i



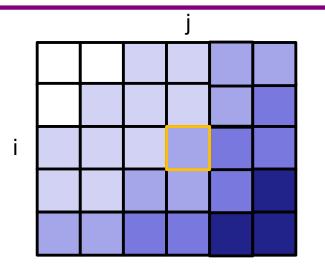
That finds the right column in row 0

- can now check M[0,j] = x (if j < n)
- if not, we can move onto the next row
 - x cannot be anywhere in the row if it's not at M[i,j]
 - set i = i + 1

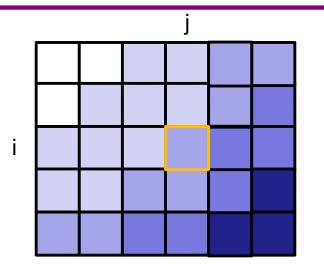
Process continues in each row thereafter...



- Make progress by setting i = i + 1
- When i increases, the invariant may be broken
 - we have $x \le M[i,j] \le M[i+1,j]$ since columns are sorted
 - and M[i+1,j] ≤ M[i +1,j+1], .., M[i +1,n-1] since rows are sorted
 - so we get $x \le M[i + 1, j], ..., M[i + 1, n-1]$



- Make progress by setting i = i + 1
- When i increases, the invariant may be broken
 - we have x <= M[i +1,j], ..., M[i +1,n-1]</p>
 - may need to restore invariant for M[i,0], ..., M[i,j-1] < x
 - decrease j until it holds again...
 - when have we seen this before?
 - initialization



- Make progress by setting i = i + 1
- When i increases, the invariant may be broken
 - we have x <= M[i +1,j], ..., M[i +1,n-1]</p>
 - may need to restore invariant for M[i,0], ..., M[i,j-1] < x
 - could copy and paste the same loop
 - or you can do it with one copy

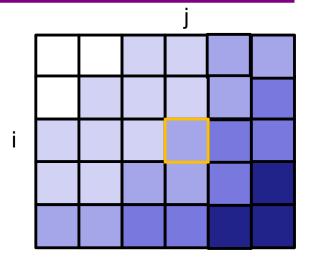
Don't try this at home!

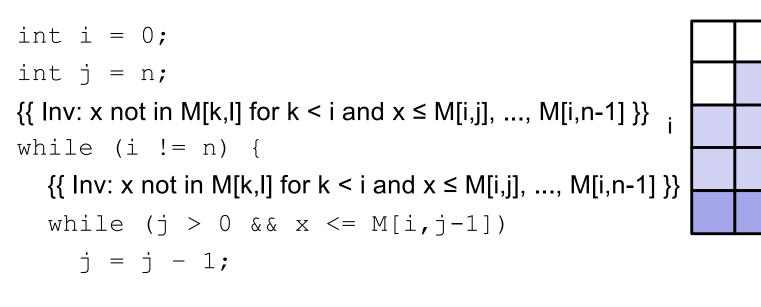
instead of

we can write

```
int i = 0, j = n;
[move j left]
{{ Inv: M[i,0], ..., M[i,j-1] < x \le M[i,j], ..., M[i,n-1] }}
while (i != n) {
  i = i + 1;
  [move j left]
}
int i = 0, j = n;
while (i != n) \{
  [move j left]
  {{ M[i,0], ..., M[i,j-1] < x \le M[i,j], ..., M[i,n-1] }
  i = i + 1;
}
```

```
int i = 0;
int j = n;
while (i != n) {
  {{ Inv: x \le M[i,j], ..., M[i,n-1] }}
  while (j > 0 \&\& x \le M[i, j-1])
     j = j - 1;
  {{ M[i,0], ..., M[i,j-1] < x \le M[i,j], ..., M[i,n-1] }}
  if (j < n \&\& x == M[i,j])
     return true;
  i = i + 1;
}
return false;
```





```
{{ x not in M[k,I] for k < i and M[i,0], ..., M[i,j-1] < x ≤ M[i,j], ..., M[i,n-1] }}
if (j < n && x == M[i,j])
return true;
i = i + 1;
}
return false;</pre>
```