CSE 331 Software Design & Implementation

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Lecture 2 – Reasoning About Straight-Line Code

Outline

- Adminstrivia
- Recap (highlights only)
- Q&A
- Exercises
- More Examples (bit more complex)

Administrivia: HW

- HW0 may have been a struggle
 - will show you how to make this easy
- HW1 posted shortly
 - worksheet
 - practice applying these ideas
 - verifying correctness of short, non-loop code
 - due on Monday by 11pm

Administrivia: Section Splits

Sections

Each section will split into two sub-sections. For example, section AA on the calendar becomes AA-1 and AA-2 that both meet at 8:30am. The table below shows which students should go to which of the two subsections based on one of the digits in their **UW Student Number**.

See the Zoom page to find the link to the meeting for that section (e.g., "Section AA-1").

Time	Name	Split	Value	TA
8:30	AA-1	last digit	odd	Yihang
	AA-2	last digit	even	Chloe
9:30	AB-1	last digit	odd	Alexey
	AB-2	last digit	even	Rachel
10:30	AC-1	last digit	odd	Andrew
	AC-2	last digit	even	Manchen
11:30	AD-1	last digit	odd	Dmitriy
	AD-2	last digit	even	Chanwut
12:30	AE-1	second digit	odd	Frank
	AE-2	second digit	even	Jasmine

https://canvas.uw.edu/courses/1370605/pages/sections

Administrivia: Section

- Each section has 16-20 students
 - hopefully, you will get to know the other students
- Section plan: Q&A, review, worksheet
 - may want to print the worksheet beforehand (if you can)
 - worksheet is similar to HW1

Quick Recap (10 min)

Correctness Toolkit

Hoare Logic

A Hoare triple is two assertions and one piece of code:

- P the precondition
- S the code
- Q the postcondition



specification method body

- A Hoare triple { P } S { Q } is called valid if:
 - in any state where P holds, executing S produces a state where Q holds
 - i.e., if P is true before S, then Q must be true after it
 - otherwise the triple is called invalid
 - code is correct iff triple is valid

Reasoning Forward & Backward

- Forward:
 - start with the given precondition
 - fill in the strongest postcondition



- Backward
 - start with the required postcondition
 - fill in the weakest precondition

Finds the "best" assertion that makes the triple valid

Reasoning: Assignments

 $x = \dots$

- Forward
 - add the fact "x = ..." to what is known
 - BUT you must fix any existing references to "x"
- Backward
 - just replace "x" with "..." in the postcondition (substitution)

Reasoning: If Statements

Forward reasoning

```
{{ P}}
 if (cond)
 → {{ P and cond }}
   S1
   {{ P1 }}
 else
 {{ P and not cond }}
  - {{ P2 }}
{{ P1 or P2 }}
```

```
Backward reasoning
  {{ cond and Q1 or
    not cond and Q2 }}
  if (cond)
   - {{ Q1 }}
     S1
   → {{ Q }}
  else
    - {{ Q2 }}
     S2
   → {{ Q }}
  {{ Q }}
```

Validity with Fwd & Back Reasoning

Reasoning in either direction gives valid assertions

Just need to check adjacent assertions:

top assertion must imply bottom one

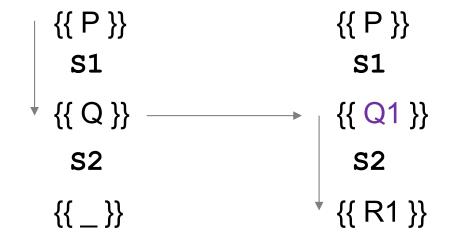
```
{{ P }}
s1
s2
{{ P1 }}
{{ Q }}
```

```
{{ P }}
{{ Q1 }}
s1
s2
{{ Q }}
```

Q & A

Dropping Irrelevant Facts

- Forward reasoning often adds many irrelevant facts
- Dropping them is usually okay



- Result is still a valid triple (ok to weaken postcondition)
- BUT no longer the strongest postcondition
- May get a final postcondition that doesn't imply the given one
- In that case, put them back and try again...

Exercises

Hoare Triples

Valid or invalid?

(Assume all variables are integers without overflow)

```
{x != 0} y = x*x; {y > 0} valid
{z != 1} y = z*z; {y != z} invalid
{x >= 0} y = 2*x; {y > x} invalid
{} if(x > 7) {y=4;} else {y=3;} {y < 5} valid</li>
{} x = y; z = x; {y=z} valid
{x=7 \lambda y=5} tmp=x; x=tmp; y=x; invalid
{y=7 \lambda x=5}
```

Forward Reasoning

```
{{ x >= 0 }}
if (x != 0) {
   z = x;
} else {
   z = x + 1;
}
```

Forward Reasoning

```
\{\{ x >= 0 \}\}
  if (x != 0) {
      \{\{ x > 0 \}\}
        z = x;
      \{\{ x > 0 \text{ and } z = x \}\}
  } else {
      \{\{ x = 0 \}\}
        z = x + 1;
      \{\{ x = 0 \text{ and } z = 1 \}\}
\{\{(x > 0 \text{ and } z = x) \text{ or } (x = 0 \text{ and } z = 1)\}\} \Rightarrow \{\{(z > 0)\}\} \text{ but strictly weaker } \}
```

Backward Reasoning

```
{{
    if (x > 7) {
        y = x;
    } else {
        y = 20;
    }
{{ y > 5 }}
```

Backward Reasoning

```
\{\{(x > 7 \text{ and } x > 5) \text{ or } (x <= 7)\}\}\ \Leftrightarrow \{\{(x > 7) \text{ or } (x <= 7)\}\}\
                                                      \Leftrightarrow {{ }}
  if (x > 7) {
      \{\{x > 5\}\}
        y = x;
      \{\{y > 5\}\}
  } else {
      \{\{\ 20 > 5\}\} \Leftrightarrow \{\{\}\}\}
        y = 20;
      \{\{y > 5\}\}
\{\{y > 5\}\}
```

More Examples

Compute x/2 rounded toward minus infinity.

```
{{}}
if (x >= 0)
y = x/2;
else
y = -((-x+1)/2);
{{ 2y = x \text{ or } 2y = x - 1}}
```

Note that, in Java, a/b rounds toward zero.

```
{{}}
if (x >= 0)
   \{\{ x >= 0 \}\}
    y = x/2;
 \rightarrow {{ 2y = x or 2y = x - 1 }}
else
   \{\{ x < 0 \} \}
   y = -((-x+1)/2);
\rightarrow {{ 2y = x or 2y = x - 1 }}
\{\{2y = x \text{ or } 2y = x - 1\}\}
```

```
{{ }}
if (x >= 0)
   {{ x >= 0 }}
   y = x/2;
   {{ 2y = x or 2y = x - 1 }}
else
   {{ x < 0 }}
   y = -((-x+1)/2);
   {{ 2y = x or 2y = x - 1 }}
{{ 2y = x or 2y = x - 1 }}</pre>
```

```
{{}}

if (x >= 0)

{{ x >= 0}}

y = x/2;

since x >= 0, "/" rounds down so this is valid

{{ 2y = x \text{ or } 2y = x - 1}}

else

{{ x < 0}}

y = -((-x+1)/2);

{{ 2y = x \text{ or } 2y = x - 1}}

{{ 2y = x \text{ or } 2y = x - 1}
```

```
{{ }}
if (x >= 0)
...
else
{{ x < 0 }}
y = (x+1)/2; // was y = -((-x+1)/2);
y = -y;
{{ 2y = x \text{ or } 2y = x - 1 }}
{{ 2y = x \text{ or } 2y = x - 1 }}
```

```
{{ }}
if (x >= 0)
...
else
   {{ x < 0 }}
   y = (-x+1)/2;
   {{ 2y = -x or 2y = -x + 1 }}
   y = -y;
   {{ 2y = x or 2y = x - 1 }}</pre>
```

```
{{ }}
if (x >= 0)
...
else
  {{ x < 0 }}
  y = (-x+1)/2;
  {{ 2y = (-x + 1) - 1 or 2y = -x + 1 }}
  y = -y;
  {{ 2y = x or 2y = x - 1 }}</pre>
```

```
{{ }}
if (x >= 0)
...
else
{{ x < 0 }}
y = (-x+1)/2;
\{ 2y = (-x+1) - 1 \text{ or } 2y = -x+1 \}}
y = -y;
{{ 2y = x \text{ or } 2y = x-1 \}}
```

Useful Subscripts Example: swap

Consider code for a swapping x and y

```
{{ }}

tmp = x;

{{ tmp = x}}

x = y;

{{ tmp = x_0 \text{ and } x = y}}

y = tmp;

{{ tmp = x_0 \text{ and } x = y_0 \text{ and } y = tmp}}
```

- Post condition implies $x = y_0$ and $y = x_0$
- I.e., their final values are equal to the original values swapped