## Section 7:

## Dijkstra's Algorithm

Slides adapted from Alex Mariakakis with material by Kellen Donohue, David Mailhot, Dan Grossman, Mike Ernst, Michael Hart, and Jacob Murphy

## Review: Shortest Paths with BFS



From Node B

| Destination | Path | Cost |
| :---: | :---: | :---: |
| A | <B,A $>$ | 1 |
| B | <B $>$ | 0 |
| C | $<B, A, C>$ | 2 |
| D | $<B, D>$ | 1 |
| E | $<B, D, E>$ | 2 |

## Review: Shortest Paths with BFS



From Node B

| Destination | Path | Cost |
| :---: | :---: | :---: |
| A | $<B, A>$ | 1 |
| B | $<B>$ | 0 |
| C | $<B, A, C>$ | 2 |
| D | $<B, D>$ | 1 |
| E | $<B, D, E>$ | 2 |

## Shortest Paths with Weights



$$
B->E=106 ?
$$

## How can we find the shortest path with weights?



## Shortest Paths with Weights



Paths are not the same!

## BFS vs. Dijkstra's Algorithm



BFS can find the most direct path, but not necessarily the shortest path!
Note that Dijkstra's Algorithm only works if there is not a negative cycle

## Dijkstra's Algorithm

Named after its inventor Edsger Dijkstra (1930-2002)

- Among his contributions to the growing CS community was his work on Operating Systems, in which he motivated the design and structure of a large project, not just the code

The Algorithm: similar to BFS, but incorporating weights

- Create a set of nodes to examine next, but instead of using the node that was next in line, use the node with the shortest distance
- How can you find the node with the shortest distance so far?

Priority Queues (explained later)

## Dijkstra's Algorithm

Give a node two fields: cost and finished, cost gives an upper bound to the distance from the origin to that node, and finished describing if the cost of the node is the minimum cost of travelling to that node.

1. For each node $v$, set $v$. cost $=\infty$ andv.finished $=$ false
2. Set source.cost $=0$ (source is the starting node of our path)
3. While there are unknown nodes in the graph
a) Select the unknown node $v$ with lowest cost
b) Mark v as finalized
c) For each edge ( $v, u)$ with weight $w$,
```
c1 = v.cost + w // cost of best path through v to u
c2 = u.cost // cost of best path to u previously known
if (c1 < c2) // if the new path through v is better,update
            u.cost = c1
            u.path = v // add v to the nodes u has traversed
```


## Example \#1



Order Added to Known Set:

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B |  | $\infty$ |  |
| C |  | $\infty$ |  |
| D |  | $\infty$ |  |
| E |  | $\infty$ |  |
| F |  | $\infty$ |  |
| H |  |  |  |

## Example \#1



## Example \#1



## Example \#1



## Example \#1



## Example \#1



## Example \#1



## Example \#1



## Example \#1



Order Added to Known Set:
A, C, B, D, F

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E |  | $\leq 12$ | C |
| F | Y | 4 | B |
| G |  | $\infty$ |  |
| H |  | $\leq 7$ | F |

## Example \#1



Order Added to Known Set:
A, C, B, D, F, H

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E |  | $\leq 12$ | C |
| F | Y | 4 | B |
| G |  | $\infty$ |  |
| H | Y | 7 | F |

## Example \#1



Order Added to Known Set:
$A, C, B, D, F, H$

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E |  | $\leq 12$ | C |
| F | Y | 4 | B |
| G |  | $\leq 8$ | H |
| H | Y | 7 | F |

## Example \#1



## Example \#1



## Example \#1



## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results


(A)

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Interpreting the Results



| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

## Example \#2



Order Added to Known Set:

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B |  | $\infty$ |  |
| C |  | $\infty$ |  |
| D |  | $\infty$ |  |
| E |  | $\infty$ |  |
| F |  | $\infty$ |  |
| G |  | $\infty$ |  |

## Example \#2



Order Added to Known Set:
A, D, C, E, B, F, G

| vertex | known? | cost | path |
| :---: | :---: | :---: | :---: |
| A | Y | 0 |  |
| B | Y | 3 | E |
| C | Y | 2 | A |
| D | Y | 1 | A |
| E | Y | 2 | D |
| F | Y | 4 | C |
| G | $Y$ | 6 | D |

## Pseudocode

// pre-condition: start is the node to start at
// initialize things
active $=$ new empty priority queue of paths
from start to a given node
// A path's "priority" in the queue is the total // cost of that path.
finished $=$ new empty set of nodes
// Holds nodes for which we know the // minimum-cost path from start.
// We know path start->start has cost 0
Add a path from start to itself to active

## Pseudocode (cont.)

while active is non-empty:

```
minPath = active.removeMin()
minDest = destination node in minPath
```

if minDest is in finished:
continue
for each edge $e=\langle m i n D e s t, ~ c h i l d\rangle$ : if child is not in finished:
newPath $=$ minPath $+e$ add newPath to active
add minDest to finished

## Priority Queue

Given a set of weighted paths, find the shortest path
Increase efficiency by considering lowest cost unknown vertex with sorting instead of looking at all vertices

PriorityQueue is like a queue, but returns elements by lowest value instead of FIFO

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Two ways to implement:

1. Comparable
a) class Node implements Comparable<Node>
b) public int compareTo(other)
2. Comparator
a) class NodeComparator extends Comparator<Node>
b) new PriorityQueue(new NodeComparator())
