

# Implementing an ADT: Representation invariants

CSE 331

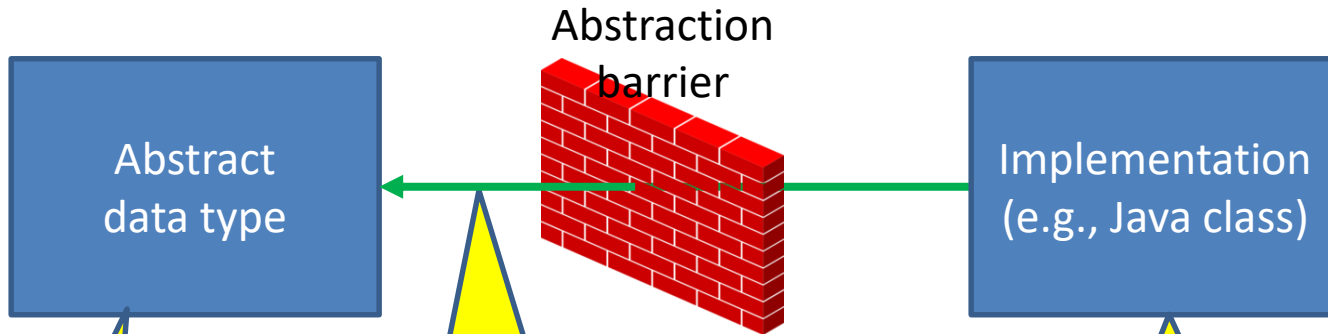
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# Outline of data abstraction lectures

ADT  
specification

ADT  
implementation



ADT  
represents  
something in  
the world

Abstraction function (AF):  
Relationship between ADT  
specification and  
implementation

**Today:**  
Representation invariant (RI):  
Relationship among  
implementation fields

# Review: A data abstraction is defined by a specification

An ADT is a collection of **procedural abstractions**

*Not* a collection of procedures

Together, these procedural abstractions provide:

A set of values

**All** the ways of directly using that set of values

Creating

Manipulating

Observing

Creators and producers: make new values

Mutators: change the value (affect `equals ()` but not `==`)

Observers: allow the client to distinguish different values

# ADTs and specifications

Specification: only in terms of the abstraction

Never mentions the representation

How should we implement an ADT?

How can we ensure the implementation satisfies the specification?

# Connecting specifications and implementations

**Representation invariant:** Object  $\rightarrow$  boolean

Is a concrete rep (an instance) **well-formed**?

**Abstraction function:** Object  $\rightarrow$  abstract value

What the instance **means** as an abstract value

Example: Rectangle (`getHeight`, `getWidth`, `getArea`)

3, 4, 12

Rep invariant says which of these are valid

0, 7, 0

2, -10, -20

Abstraction function says which of these the instance represents

# Review:

## Implementing a data abstraction (ADT)

To implement a data abstraction:

- Select the representation of instances, the *rep*
- Implement operations in terms of that rep

Choose a representation so that

- It is possible to implement operations
- The most frequently used operations are efficient
  - You don't know which these will be
  - Abstraction allows the rep to change later

# CharSet Abstraction

// Overview: A CharSet is a finite mutable set of Characters

// effects: creates a fresh, empty CharSet

public CharSet()

// modifies: this

// effects:  $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$

public void insert(Character c);

// modifies: this

// effects:  $\text{this}_{\text{post}} = \text{this}_{\text{pre}} - \{c\}$

public void delete(Character c);

// returns: ( $c \in \text{this}$ )

public boolean member(Character c);

// returns: cardinality of this

public int size();

# A CharSet implementation

```
class CharSet {
    private List<Character> elts
        = new ArrayList<Character>();

    public void insert(Character c) {
        elts.add(c);
    }
    public void delete(Character c) {
        elts.remove(c);
    }
    public boolean member(Character c) {
        return elts.contains(c);
    }
    public int size() {
        return elts.size();
    }
}
```



# A buggy CharSet implementation.

## What client code will expose the error?

```
class CharSet {  
    private List<Character> elts  
        = new ArrayList<Character>();  
  
    public void insert(Character c) {  
        elts.add(c);  
    }  
    public void delete(Character c) {  
        elts.remove(c);  
    }  
    public boolean member(Character c) {  
        return elts.contains(c);  
    }  
    public int size() {  
        return elts.size();  
    }  
}
```

Where is the defect?

```
CharSet s = new CharSet();  
s.insert('a');  
s.insert('a');  
s.delete('a');  
if (s.member('a'))  
    // print "wrong";  
else  
    // print "right";
```

# Where is the defect?

The answer to this question tells you what code needs to be fixed

*Perhaps **delete** is wrong (and so is **size**)*

It should remove all occurrences

*Perhaps **insert** is wrong*

It should not insert a character that is already there

How can we know?

The **representation invariant** tells us

# The representation invariant

- Defines data structure well-formedness
  - Which instances/refs are valid?
- Holds before and after every **CharSet** operation
- Operation implementations (methods) may depend on it

Write it this way:

```
class CharSet {  
    // Rep invariant: elts has no nulls and no duplicates  
    private List<Character> elts;  
    ...  
}
```

Or, if you are the pedantic sort:

$\forall$  indices  $i$  of `elts` . `elts.elementAt(i)  $\neq$  null`

$\forall$  indices  $i, j$  of `elts` .

$i \neq j \Rightarrow \neg \text{elts.elementAt}(i).\text{equals}(\text{elts.elementAt}(j))$

# Now, we can locate the error

```
// Rep invariant:  
// elts has no nulls and no duplicates  
  
public void insert(Character c) {  
    elts.add(c);  
}  
  
public void delete(Character c) {  
    elts.remove(c);  
}
```

# Another rep invariant example

```
class Account {  
    private int balance;  
    // history of all transactions  
    private List<Transaction> transactions;  
    ...  
}
```

Real-world constraints:

- $\text{balance} \geq 0$
- $\text{balance} = \sum_i \text{transactions.get}(i).\text{amount}$

Implementation-related constraints:

- $\text{transactions} \neq \text{null}$
- no nulls in transactions

# Listing the elements of a CharSet

Consider adding the following method to CharSet:

```
// returns: a List containing the members of this  
public List<Character> getElts();
```

Consider this implementation:

```
// Rep invariant: elts has no nulls and no duplicates  
public List<Character> getElts() { return elts; }
```

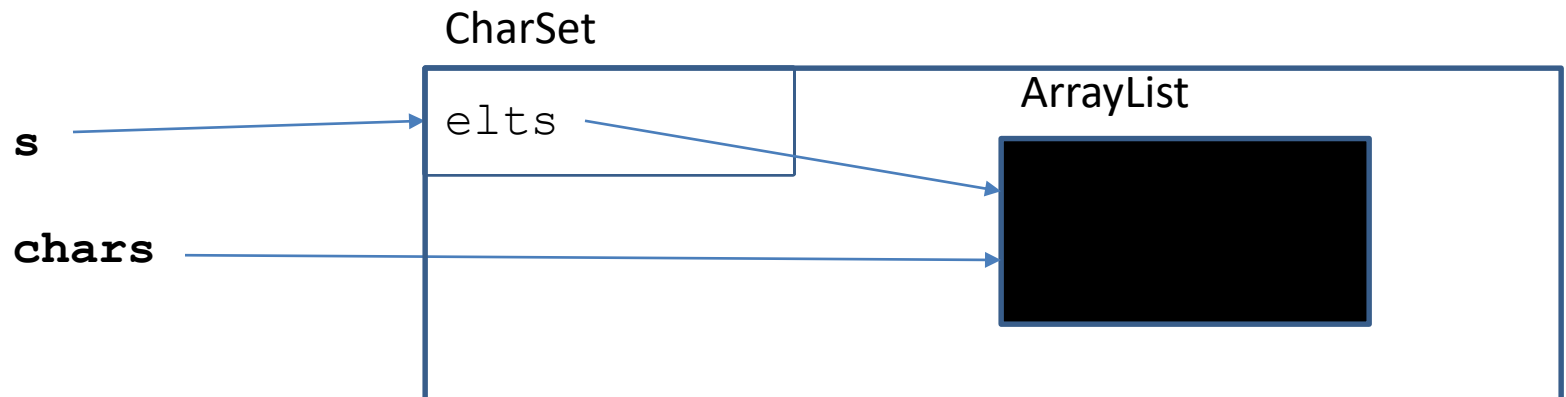
Does the implementation of getElts preserve the rep invariant?

... sort of

# Representation exposure

Consider this client code (outside the CharSet implementation):

```
CharSet s = new CharSet();  
s.insert('a');  
List<Character> chars = s.getElts();  
chars.add('a');  
s.delete('a');  
if (s.member('a')) ...
```



# Representation exposure

Consider this client code (outside the CharSet implementation):

```
CharSet s = new CharSet();  
s.insert('a');  
List<character> chars = s.getElts();  
chars.add('a');  
s.delete('a');  
if (s.member('a')) ...
```

**Representation exposure** is external access to the rep

A **big deal**, a **common bug**. Now you have a name for it.

Representation exposure is almost always **EVIL**

Enables violation of abstraction boundaries and the rep invariant

If you do it, document why and how

And feel guilty about it!

How to avoid/prevent rep exposure: **immutability** or **copying**



# Avoid rep exposure #1: Immutability

Aliasing is no problem if the client cannot change the data

Assume `Point` is an *immutable* ADT:

```
class Line {
    private Point start;
    private Point end;
    public Line(Point start, Point end) {
        this.start = start;
        this.end = end;
    }
    public Point getStart() {
        return this.start;
    }
    ...
}
```

# Pros and cons of immutability

Immutability greatly simplifies reasoning

- Aliasing does not matter
- No need to make copies with identical contents
- Rep invariants cannot be broken

Can be less efficient (new objects for every modification)

Can be more efficient (no need for redundant copies)

Does require different designs.

Suppose **Point** is immutable but **Line** is mutable:

```
class Line {  
    ...  
    void raiseLine(double deltaY) {  
        this.start = new Point(start.x, start.y + deltaY);  
        this.end = new Point(end.x, end.y + deltaY);  
    }  
}
```

Immutable Java classes include **Character**, **Color**, **File** (path), **Font**, **Integer**, **Locale**, **String**, **URL**, ...

# Are `private` and `final` enough?

## Making fields `private`

- Is necessary to prevent rep exposure (why?)
- Is insufficient to prevent rep exposure (see CharSet example)
- The real issue is **aliasing of mutable data**

## Making fields `final`

- Is neither necessary nor sufficient to achieve immutability
- A `final` field cannot be reassigned
  - But it can be mutated (its fields can be reassigned and/or mutated)

```
class Line {
    private final Point start;
    private final Point end;
    ...
    public void translate(int deltaX, int deltaY) {
        start.x += deltaX;
        start.y += deltaY;
        end.x += deltaX;
        end.y += deltaY;
    }
    ...
}
```

# Avoiding rep exposure #2: Copying

Copy data that crosses the abstraction barrier

Example (assume Point is a mutable ADT):

```
class Line {  
    private Point start;  
    private Point end;  
  
    public Line(Point start, Point end) {  
        this.start = new Point(start.x, start.y);  
        this.end = new Point(end.x, end.y);  
    }  
  
    public Point getStart() {  
        return new Point(this.start.x, this.start.y);  
    }  
    ...  
}
```

Copy in: parameters that become part of the implementation

Copy out: results that are part of the implementation

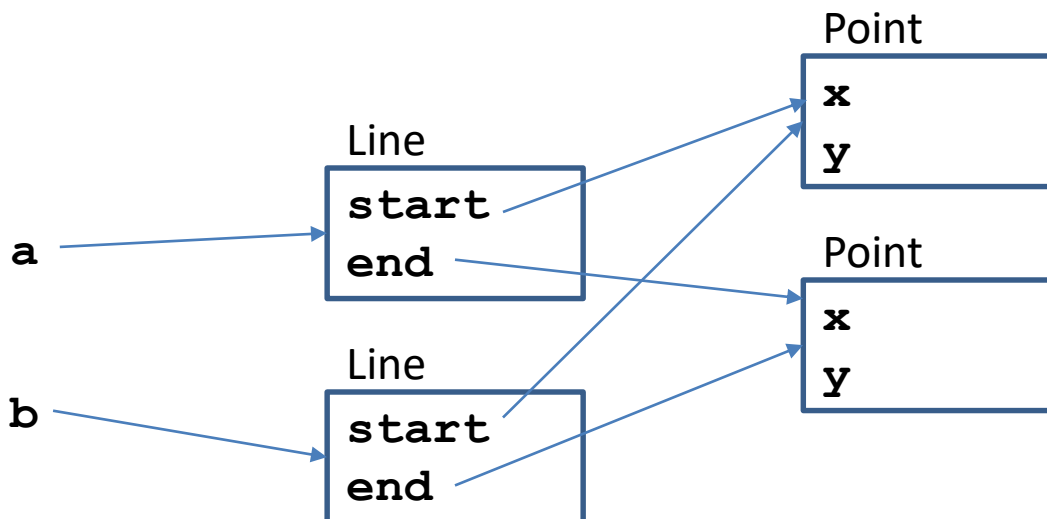
# Shallow copying is not enough

Example (assume Line and Point are mutable ADTs):

```
class Line {  
    private Point start;  
    private Point end;  
  
    public Line(Line other) {  
        this.start = other.start;  
        this.end = other.end;  
    }  
}
```

Client code:

```
Line a = ...;  
Line b = new Line(a);  
a.translate(3, 4);
```



# Deep copying is not necessary

Must copy-in and copy-out “all the way down”  
to immutable parts

This combines our two ways to avoid rep  
exposure: **immutability** and **copying**.

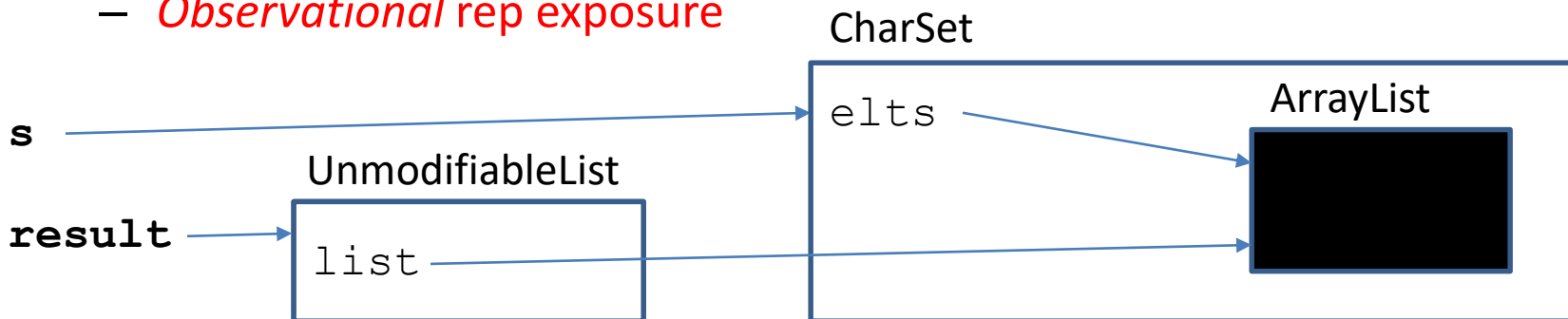
# Avoiding rep exposure #3: Readonly wrapper (“immutable copy”)

```
class CharSet {  
    private List<Character> elts = ...;  
  
    public List<Character> getElts() { // copy  
        return new ArrayList(elts);  
    }  
  
    public List<Character> getElts() { // readonly wrapper  
        return Collections.unmodifiableList(elts);  
    }  
}
```

```
CharSet s;  
result = s.getElts();  
s.add('a');
```

`unmodifiableList()`: result can be read but not modified

- Doesn't make a copy: its rep is aliased to its input (efficient!)
- Attempts to modify throw `UnsupportedOperationException`
- Still need to copy on the way in
- *Observational rep exposure*



# The specification “Returns a list containing the elements”

Could mean any of these things:

1. Returns a fresh mutable list containing the elements in the set *at the time of the call*.
  - Difficult to implement efficiently
2. Returns read-only view that is *always up to date* with the current elements in the set.
  - Makes it hard to change the rep later
3. Returns a list containing the current set elements. *Behavior is unspecified* if client attempts to mutate the list or to access the list after mutating the set.
  - Weaker than #1 and #2
  - Less simple, harder to use, but sufficient for some purposes

Lesson: A seemingly simple spec may be *ambiguous and subtle*!



# Avoiding representation exposure

*Understand* what representation exposure is

*Design* ADT implementations to prevent it

*Prove* that your ADT is free of representation exposure

*Test* for it with adversarial clients:

- Pass values to methods and then mutate them
- Mutate values returned from methods
- Check the rep invariant (in addition to client behavior)

*Fix* any rep exposure bugs

# Checking rep invariants

Should code check that the rep invariant holds?

- Yes, if it's inexpensive
- Yes, for debugging (even when it's expensive)
- It's quite hard to justify turning the checking off
- Some private methods need not check (Why?)
- Some private methods should not check (Why?)

# Checking the rep invariant

Rule of thumb: check on entry *and* on exit (why?)

```
public void delete(Character c) {
    checkRep();
    elts.remove(c)
    // Is this guaranteed to get called?
    // See handouts for a less error-prone way to check at exit.
    checkRep();
}
...
/** Verify that elts contains no duplicates. */
private void checkRep() {
    for (int i = 0; i < elts.size(); i++) {
        assert elts.indexOf(elts.elementAt(i)) == i;
    }
}
```

# Practice defensive programming

Assume that you will make mistakes

Write and incorporate code designed to catch them

On entry:

- Check rep invariant

- Check preconditions (requires clause)

On exit:

- Check rep invariant

- Check postconditions

Checking the rep invariant helps you **discover** errors

Reasoning about the rep invariant helps you **avoid** errors

Or prove that they do not exist!

# The rep invariant constrains structure, not meaning

New implementation of insert that preserves the rep invariant:

```
public void insert(Character c) {
    Character cc = new Character(encrypt(c));
    if (!elts.contains(cc))
        elts.addElement(cc);
}
public boolean member(Character c) {
    return elts.contains(c);
}
```

The program is still wrong

Clients observe incorrect behavior

What client code exposes the error?

Where is the error?

We must consider the **meaning**

The *abstraction function* helps us

```
CharSet s = new CharSet();
s.insert('a');
if (s.member('a'))
    // print "right";
else
    // print "wrong";
```