## CSE 331 Fall 2017 Midterm Exam

Name $\qquad$

The exam is closed book and closed electronics. One page of notes is allowed. Please wait to turn the page until everyone is told to begin.

Score: $\qquad$ / 62

1. $\qquad$ / 12
2. $\qquad$ / 12
3. $\qquad$ / 12
4. $\qquad$ 110
5. $\qquad$ / 8
6. $\qquad$ / 8

## Problem 1 (Specifications)

a. Alice is writing a function to find the minimum of an array of numbers. She intends to implement it by sorting the array, but she does not want clients to depend on that fact. Write a specification for her function:

```
/**
    * @requires vals != null
    * @modifies vals
    * @effects
    * @returns minimum number in vals
    * @throws
    */
int findMin(int[] vals) { ...
```

b. Suppose that Alice decides to change her implementation to no longer sort the array. How should she change the specification above?

Remove vals from @modifies.
c. This new specification would be (circle one):
weaker incomparable stronger
d. Suppose that Alice decides instead to stick with the version that sorts the array but will now allow clients to depend on that behavior. How should she change the specification above?

Add "vals is sorted" to @effects.
e. This new specification would be (circle one):
weaker incomparable stronger

## Problem 2 (Reasoning)

Fill in an implementation of the method runLengthEncode. It takes as input a string, str, an array of characters, chars, and an array of ints, lens. You can assume the string and both arrays are of length at least $n$. You can assume that str is non-empty and that it does not contain the character ' 10 '.

Your method will write its output into the arrays chars and lens, and it should return a number $t$ such that (after returning) str $=$ chars[0] * lens[0] + ... + chars[t-1] * lens[t-1], where a char * int means a string containing that many copies of the char. For example, if $s t r=$ "aaabbccccaaddd", it would return $t=5$ and leave chars[0..4] = $[a, b, c, a, d]$ and lens $[0 . .4]=[3,2,4,2,3]$.

The invariant for the loop is already provided. Do not add any additional loops.

You do not need to turn in a complete proof of correctness, but you should complete one since your code will be graded on correctness.
$\{\{\mathrm{P}: 0<\mathrm{n}<=$ str.length chars.length, lens.length \}\}

```
int runLengthEncode(String str, int n, char[] chars, int[] lens) {
    int i = 0;
    int j = -1;
    char cur = '\0';
```

    \(\{\{\operatorname{lnv}: P\) and \(\operatorname{str}[0 . . i-1]=\) chars[0] * lens[0] + ... + chars[j] * lens[j] and
                chars[0] != chars[1], ..., chars[j-1] != chars[j] and
                ( \(\mathrm{i}=0\) or cur \(=\operatorname{str}[i-1]\) ) \(\}\}\)
    while (i != n) \{
        if (str.charAt(i) == cur) \{
            lens[j] = lens[j] + 1;
        \} else \{
            \(j=j+1\);
            cur = str.charAt(i);
            chars[j] = cur;
            lens[j] = 1;
        \}
        \(i=i+1 ;\)
    \}
    \{\{ str[0..n-1] = chars[0] * lens[0] + ... + chars[j] * lens[j] and
        chars[0] != chars[1], ..., chars[j-1] != chars[j]\}\}
    return j+1;
    \}

## Problem 3 (Reasoning II)

Fill in another implementation for runLengthEncode below.

For this version, you will use nested loops. The invariants for both loops are already provided. Do not add any additional loops.

You do not need to turn in a complete proof of correctness, but you should complete one since your code will be graded on correctness.

```
{{P:0<n<= str.length, chars.length, lens.length }}
int runLengthEncode(String str, int n, char[] chars, int[] lens) {
    int i = -1;
    int j = -1;
    {{ Inv: P and str[0..i] = chars[0] * lens[0] + ... + chars[j] * lens[j] and
                chars[0] != chars[1], .., chars[j-1] != chars[j] and
            (i < 0 or n <= i+1 or str[i] != str[i+1]) }}
    while (i != n-1) {
        int k = i;
        {{ Inv: Inv and str[i+1] = str[i+2] = ... = str[k+1] }}
        while (k+1 != n-1 && str.charAt(i+1) == str.charAt(k+2)) {
            k = k + 1;
        }
        j = j + 1;
        chars[j] = str.charAt(i+1);
        lens[j] = k + 1 - i;
        i = k + 1;
    }
    {{ str[0..n-1] = chars[0] * lens[0] + ... + chars[j] * lens[j] and
        chars[0] != chars[1], .., chars[j-1] != chars[j] }}
    return j+1;
```

Problem 4 (Testing)
Describe three test cases for the runLengthEncode method on the previous pages.
The three tests should fall into different subdomains, i.e., they should be from subsets of the input where the expected or actual behavior is fundamentally different.

1. Input: $\qquad$ "abc" $\qquad$ and $n=$ $\qquad$ 3 $\qquad$
Output: returns $\qquad$ 3 $\qquad$ chars starts with $\qquad$ ['a', 'b', 'c'] $\qquad$ lens starts with $\qquad$ $[1,1,1]$ $\qquad$
2. Input:
$\operatorname{str}=$ $\qquad$ "aabbbcc" $\qquad$ and $n=$ $\qquad$
Output: returns $\qquad$ 3 $\qquad$
chars starts with $\qquad$ ['a', 'b', 'c'] $\qquad$
lens starts with $\qquad$ [2, 3, 2]

If it's not obvious, why is this testing a different behavior ${ }^{1}$ from the case above?
The method should return a number smaller than str's length.
3. Input:
str $=$ $\qquad$ "abc" $\qquad$ and $\mathrm{n}=$ $\qquad$
Output: returns $\qquad$ 2 $\qquad$
chars starts with $\qquad$ ['a', 'b'] $\qquad$
lens starts with $\qquad$ [1, 1] $\qquad$
If it's not obvious, why is this testing a different behavior ${ }^{1}$ from the cases above?
The method should ignore the last character of the input.

[^0]
## Problem 5 (ADTs)

Suppose that we created a CharList ADT whose abstract value is a string but whose concrete representation was the run-length encoding used in the previous problems:

```
/** Represents an immutable sequence of characters like "abc" or "". */
class CharList {
    private char[] chars;
    private int[] lens;
    private int count; // number of entries used in above arrays
```

...
(Note: count corresponds to the return value of runLengthEncode.)
What would the representation invariant ${ }^{2}$ be for this ADT?
chars != null and lens != null and $0<=$ count <= chars.length, lens.length

What would the abstraction function ${ }^{2}$ be for this ADT?
chars[0] * lens[0] + ... + chars[count-1] * lens[count-1]

Fill in the implementation of the following method:

```
@Override // (returns the abstract value, which is a string)
public String toString() {
    StringBuilder buf = new StringBuilder();
    for (int i = 0; i < count; i++) {
        for (int j = 0; j < lens[i]; j++)
            buf.append(chars[i]);
    }
    return buf.toString();
}
```

[^1]
## Problem 6 (Miscellaneous)

a. Suppose that Alice has written a method and has a postcondition that correctly describes what she wants it to do but does not yet have a precondition. How could she come up with one that is certain to make her code correct?

## Use backward reasoning

b. When is it safe to mutate an object being used as a key in a HashMap?
i. when the key and value are the same object
ii. when the associated value is immutable
iii. never
c. Which of the following types of operations are NOT usually included in the results of a requirements analysis? (Circle one.)
i. operations explicitly mentioned in relevant use cases
ii. operations we can infer will be necessary to complete use cases
iii. operations on which our chosen representation is most efficient
iv. operations clients would strongly expect to see based on conventions in the language or its standard libraries
d. Which of the following is the most important reason that it is difficult to write arbitrarily large software programs?
i. software tends to become out-of-date with hardware changes as it grows
ii. software tends to become too complex to understand as it grows
iii. compilers are too asymptotically slow to use on large amounts of code
iv. computers would not have enough memory to load the bytecode
e. Which of the following is a most often a symptom of writing poor quality code?
i. lots of time spent adding assertions
ii. lots of time spent testing
iii. lots of time spent debugging


[^0]:    ${ }^{1}$ You can define behavior, e.g., in terms of expected (black box) or actual (clear box) execution equivalence using either implementation of runLengthEncode.

[^1]:    ${ }^{2}$ While CharList uses the same representation as the runLengthEncode methods from before, you cannot use those methods to define your RI or AF here. You should define both directly in terms of the fields, as usual.

