

Representation invariants and abstraction functions

CSE 331

UW CSE

ADTs and specifications

Specification: only in terms of the abstraction

Never mentions the representation

An ADT is more than just a data structure

data structure + a set of conventions

Why do we need to relate the specification to the representation?

Connecting specifications and implementations

Representation invariant: $\text{Object} \rightarrow \text{boolean}$

Indicates whether a data structure is **well-formed**

Defines the set of valid values of the data structure

Abstraction function: $\text{Object} \rightarrow \text{abstract value}$

What the data structure **means** (as an abstract value)

How the data structure is to be interpreted

How do you compute the inverse, abstract value \rightarrow Object ?

A data abstraction is defined by a specification

A collection of procedural abstractions

Not a collection of procedures

Together, these procedural abstractions provide

A set of values

All the ways of directly using that set of values

Creating

Manipulating

Observing

Creators and producers: make new values

Mutators: change the value (but don't affect ==)

Observers: allow one to tell values apart

Implementation of an ADT is provided by a class

To implement a data abstraction

- Select the representation of instances, the *rep*
- Implement operations in terms of that rep

Choose a representation so that

- It is possible to implement operations
- The most frequently used operations are efficient

But which will these be?

Abstraction allows the rep to change later

CharSet Abstraction

// Overview: A CharSet is a finite mutable set of Characters

// effects: creates a fresh, empty CharSet

public CharSet ()

// modifies: this

// effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$

public void insert (Character c);

// modifies: this

// effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} - \{c\}$

public void delete (Character c);

// returns: ($c \in \text{this}$)

public boolean member (Character c);

// returns: cardinality of this

public int size ();

A CharSet implementation.

What client code will expose the error?

```
class CharSet {  
    private List<Character> elts  
        = new ArrayList<Character>();  
    public void insert(Character c)    {  
        elts.add(c);  
    }  
    public void delete(Character c)    {  
        elts.remove(c);  
    }  
    public boolean member(Character c)  
        return elts.contains(c);  
    }  
    public int size() {  
        return elts.size();  
    }  
}
```

```
CharSet s = new CharSet();  
Character a  
    = new Character('a');  
s.insert(a);  
s.insert(a);  
s.delete(a);  
if (s.member(a))  
    // print "wrong";  
else  
    // print "right";
```

Where Is the Error?

The answer to this question tells you what needs to be fixed

*Perhaps **delete** is wrong*

It should remove all occurrences

*Perhaps **insert** is wrong*

It should not insert a character that is already there

How can we know?

The **representation invariant** tells us

The representation invariant

- States data structure well-formedness
- Captures information that must be shared across implementations of multiple operations

Write it this way:

```
class CharSet {  
    // Rep invariant: elts has no nulls and no duplicates  
    private List<Character> elts;  
    ...  
}
```

Or, if you are the pedantic sort:

\forall indices i of `elts` . `elts.elementAt(i) \neq null`

\forall indices i, j of `elts` .

$i \neq j \Rightarrow \neg \text{elts.elementAt}(i).\text{equals}(\text{elts.elementAt}(j))$

Now, we can locate the error

```
// Rep invariant:  
// elts has no nulls and no duplicates  
  
public void insert(Character c) {  
    elts.add(c);  
}  
  
public void delete(Character c) {  
    elts.remove(c);  
}
```

Another rep invariant example

```
class Account {  
    private int balance;  
    // history of all transactions  
    private List<Transaction> transactions;  
    ...  
}
```

// real-world constraints:

balance ≥ 0

balance = \sum_i transactions.get(i).amount

// implementation-related constraints:

transactions \neq null

no nulls in transactions

Listing the elements of a CharSet

Consider adding the following method to CharSet

```
// returns: a List containing the members of this  
public List<Character> getElts();
```

Consider this implementation:

```
// Rep invariant: elts has no nulls and no duplicates  
public List<Character> getElts() { return elts; }
```

Does the implementation of getElts preserve the rep invariant?

... sort of

Representation exposure

Consider the client code (outside the CharSet implementation)

```
CharSet s = new CharSet();  
Character a = new Character('a');  
s.insert(a);  
s.getElts().add(a);  
s.delete(a);  
if (s.member(a)) ...
```

Representation exposure is external access to the rep

Representation exposure is almost always **evil**

If you do it, document why and how

And feel guilty about it!

Ways to avoid rep exposure

1. Exploit immutability

```
Character choose() {  
    return elts.elementAt(0);  
}
```

Character is immutable.

2. Make a copy

```
List<Character> getElts() {  
    return new ArrayList<Character>(elts);  
    // or: return (ArrayList<Character>) elts.clone();  
}
```

Mutating a copy doesn't affect the original.

Don't forget to make a copy on the way in!

3. Make an immutable copy

```
List<Character> getElts() {  
    return Collections.unmodifiableList<Character>(elts);  
}
```

Client cannot mutate

Still need to make a copy on the way in

Checking rep invariants

Should code check that the rep invariant holds?

- Yes, if it's inexpensive
- Yes, for debugging (even when it's expensive)
- It's quite hard to justify turning the checking off
- Some private methods need not check (Why?)

Checking the rep invariant

Rule of thumb: check on entry *and* on exit (why?)

```
public void delete(Character c) {
    checkRep();
    elts.remove(c)
    // Is this guaranteed to get called?
    // See handouts for a less error-prone way to check at exit.
    checkRep();
}
...
/** Verify that elts contains no duplicates. */
private void checkRep() {
    for (int i = 0; i < elts.size(); i++) {
        assert elts.indexOf(elts.elementAt(i)) == i;
    }
}
```


Practice defensive programming

Assume that you will make mistakes

Write and incorporate code designed to catch them

On entry:

- Check rep invariant

- Check preconditions (requires clause)

On exit:

- Check rep invariant

- Check postconditions

Checking the rep invariant helps you **discover** errors

Reasoning about the rep invariant helps you **avoid** errors

Or prove that they do not exist!

We will discuss such reasoning, later in the term

The rep invariant constrains structure, not meaning

New implementation of insert that preserves the rep invariant:

```
public void insert(Character c) {  
    Character cc = new Character(encrypt(c));  
    if (!elts.contains(cc))  
        elts.addElement(cc);  
}  
public boolean member(Character c) {  
    return elts.contains(c);  
}
```

The program is still wrong

Clients observe incorrect behavior

What client code exposes the error?

Where is the error?

We must consider the meaning

The *abstraction function* helps us

```
CharSet s = new CharSet();  
Character a = new  
Character('a');  
s.insert(a);  
if (s.member(a))  
    // print "right";  
else  
    // print "wrong";
```

Abstraction function: $\text{rep} \rightarrow \text{abstract value}$

The **abstraction function** maps the concrete representation to the abstract value it represents

AF: Object \rightarrow abstract value

AF(CharSet this) = { c | c is contained in this.elts }

“set of Characters contained in this.elts”

Typically *not* executable

The abstraction function lets us reason about behavior **from the client perspective**

Our real goal is to satisfy the specification of insert:

// modifies: this

// effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$

public void **insert** (Character c);

Once again we can place the blame

Applying the abstraction function to the result of the call to insert yields $\text{AF}(\text{elts}) \cup \{\text{encrypt('a')}\}$

What if we used this abstraction function?

$\text{AF}(\text{this}) = \{ c \mid \text{encrypt}(c) \text{ is contained in } \text{this.elts} \}$

$\text{AF}(\text{this}) = \{ \text{decrypt}(c) \mid c \text{ is contained in } \text{this.elts} \}$

Placing the blame

Our real goal is to satisfy the **specification of insert**:

```
// modifies: this  
// effects:  $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$   
public void insert (Character c);
```

The **AF** tells us what the rep means (and lets us place the blame)

$$\text{AF}(\text{CharSet this}) = \{ c \mid c \text{ is contained in } \text{this.elts} \}$$

Consider a call to insert:

On **entry**, the meaning is $\text{AF}(\text{this}_{\text{pre}}) \approx \text{elts}_{\text{pre}}$

On **exit**, the meaning is $\text{AF}(\text{this}_{\text{post}}) = \text{AF}(\text{this}_{\text{pre}}) \cup \{\text{encrypt}('a')\}$

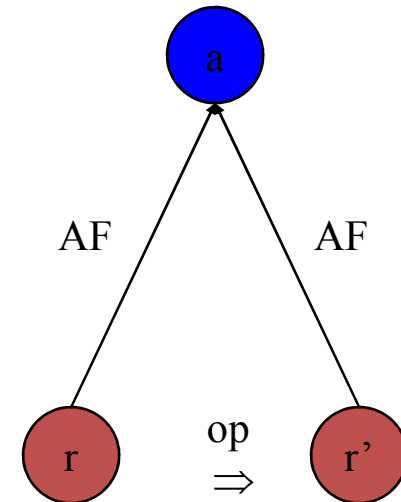
What if we used this abstraction function?

$$\begin{aligned} \text{AF}(\text{this}) &= \{ c \mid \text{encrypt}(c) \text{ is contained in } \text{this.elts} \} \\ &= \{ \text{decrypt}(c) \mid c \text{ is contained in } \text{this.elts} \} \end{aligned}$$

Benevolent side effects

Different implementation of member:

```
boolean member(Character c1) {  
    int i = elts.indexOf(c1);  
    if (i == -1)  
        return false;  
    // move-to-front optimization  
    Character c2 = elts.elementAt(0);  
    elts.set(0, c1);  
    elts.set(i, c2);  
    return true;  
}
```



Move-to-front speeds up repeated membership tests
Mutates rep, but does not change *abstract* value
AF maps both reps to the same abstract value

The abstraction function is a function

Q: Why do we map concrete to abstract rather than vice versa?

1. It's not a function in the other direction.

E.g., lists $[a,b]$ and $[b,a]$ each represent the set $\{a, b\}$

2. It's not as useful in the other direction.

Can construct objects via the provided operators

Writing an abstraction function

The domain: all representations that satisfy the rep invariant

The range: can be tricky to denote

- For mathematical entities like sets: easy

- For more complex abstractions: give them fields

 - AF defines the value of each “specification field”

 - For “derived specification fields”, see the handouts

The overview section of the specification should provide a way of writing abstract values

- A printed representation is valuable for debugging

Summary

Rep invariant

Which concrete values represent abstract values

Abstraction function

For each concrete value, which abstract value it represents

Together, they modularize the implementation

Can examine operators one at a time

Neither one is part of the abstraction (the ADT)

In practice

Always write a representation invariant

Write an abstraction function when you need it

Write an informal one for most non-trivial classes

A formal one is harder to write and usually less useful