

Representation invariants and abstraction functions

CSE 331

Autum 2010

ADTs and specifications

- An ADT is more than just a data structure
 - data structure + a set of conventions
- Specification: only in terms of the abstraction
 - Never mentions the representation
- *Representation invariant*: $\text{Object} \rightarrow \text{boolean}$
 - Indicates whether a data structure is well-formed
 - Defines the set of valid values of the data structure
- *Abstraction function*: $\text{Object} \rightarrow \text{abstract value}$
 - What the data structure means (as an abstract value)
 - How the data structure is to be interpreted
 - How do you compute the inverse, $\text{abstract value} \rightarrow \text{Object}$?

A data abstraction is defined by a specification

- A collection of procedural abstractions
 - Not a collection of procedures
- Together, these procedural abstractions provide
 - A set of values
 - All the ways of directly using that set of values: creating, manipulating, observing
- Creators and producers make new values
- Mutators change the value (but don't affect ==)
- Observers allow one to tell values apart, which is the key to understanding

Implementation of an ADT is provided by a class

- To implement a data abstraction
 - Select the representation of instances, the *rep*
 - Implement operations in terms of that rep
- Choose a representation so that
 - It is possible (preferably easy) to implement operations
 - The most frequently used operations are efficient
 - But which will these be?
 - Abstraction allows changes to rep late in the game

CharSet Abstraction

// Overview: A CharSet is a finite mutable set of Characters

// effects: creates a fresh, empty CharSet

public CharSet ()

// modifies: this

// effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$

public void insert (Character c);

// modifies: this

// effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} - \{c\}$

public void delete (Character c);

// returns: ($c \in \text{this}$)

public boolean member (Character c);

// returns: cardinality of this

public int size ();

A CharSet implementation.

Where is the error?

```
class CharSet {
    private List<Character> elts
        = new ArrayList<Character>();
    public void insert(Character c)
        elts.add(c);
    }
    public void delete(Character c)
        elts.remove(c);
    }
    public boolean member(Character c)
        return elts.contains(c);
    }
    public int size() {
        return elts.size();
    }
}
```

```
CharSet s = new CharSet();
Character a
    = new Character('a');
s.insert(a);
s.insert(a);
s.delete(a);
if (s.member(a))
    // print "wrong";
else
    // print "right";
```

Where Is the Error?

- The answer to this question tells you what needs to be fixed
- *Perhaps* **delete** is wrong
 - It should remove all occurrences
- *Perhaps* **insert** is wrong
 - It should not insert a character that is already there
- How can we know?
 - The representation invariant tells us

The representation invariant

- States data structure well-formedness
- Captures information that must be shared across implementations of multiple operations
- Write it this way

```
class CharSet {  
    // Rep invariant: elts has no nulls and no  
    // duplicates  
    private List<Character> elts;  
    ...  
}
```

- Or, if you are the pedantic sort

$\forall$ indices i of `elts` . `elts.elementAt(i)  $\neq$  null`

$\forall$ indices i, j of `elts` . $i \neq j \Rightarrow$

$\neg \text{elts.elementAt}(i).\text{equals}(\text{elts.elementAt}(j))$

Now, we can locate the error

```
// Rep invariant:  
// elts has no nulls and no duplicates  
  
public void insert(Character c) {  
    elts.add(c);  
}  
  
public void delete(Character c) {  
    elts.remove(c);  
}
```

Another rep invariant example

```
class Account {  
    private int balance;  
    // history of all transactions  
    private List<Transaction> transactions;  
    ...  
}
```

// real-world constraints:

balance ≥ 0

balance = $\sum_i$ transactions.get(i).amount

// implementation-related constraints:

transactions $\neq$ null

no nulls in transactions

Listing the elements of a CharSet

- Consider adding the following method to CharSet

```
// returns: a List containing the members of this
public List<Character> getElts ( );
```
- Consider this implementation

```
// Rep invariant: elts has no nulls & no duplicates
public List<Character> getElts() { return elts; }
```
- Does the implementation of `getElts` preserve the rep invariant?

Well, sort of:

Representation exposure

- Consider the client code

```
CharSet s = new CharSet();  
Character a = new Character('a');  
s.insert(a);  
s.getElts().add(a);  
s.delete(a);  
if (s.member(a)) ...
```

- Representation exposure is external access to the rep
- Representation exposure is almost always **evil – why?**
- If you do it, document why and how
 - And feel guilty about it!

Ways to avoid rep exposure

- Make a copy

```
List<Character> getElts() {  
    return new ArrayList<Character>(elts);  
    // or: return (ArrayList<Character>) elts.clone();  
}
```

Mutating a copy doesn't affect the original

Don't forget to make a copy on the way in!

- Make an immutable copy

```
List<Character> getElts() {  
    return Collections.unmodifiableList<Character>(elts);  
}
```

Client cannot mutate

Still need to make a copy on the way in

Checking rep invariants

- Should code check that the rep invariant holds?
 - Yes, if it's inexpensive
 - Yes, for debugging (even when it's expensive)
 - It's quite hard to justify turning the checking off
 - Some private methods need not check (Why?)

Checking the rep invariant

```
/** Verify that elts contains no duplicates. */  
private void checkRep() {  
    for (int i = 0; i < elts.size(); i++) {  
        assert elts.indexOf(elts.elementAt(i)) == i);  
    }  
}
```

- An alternative implementation
 - `repOK()` returns a boolean
 - callers of `repOK` check its return value

Check on entry *and* on exit

- As a rule of thumb... but why?

```
public void delete(Character c) {  
    checkRep();  
    elts.remove(c);  
    // Is this guaranteed to get called?  
    // See handouts for a less error-prone way  
    // to check at exit.  
    checkRep();  
}
```

Practice defensive programming

- Assume that you will make mistakes
- Write and incorporate code designed to catch them
 - On entry
 - Check rep invariant
 - Check preconditions (requires clause)
 - On exit
 - Check rep invariant
 - Check postconditions
- Checking the rep invariant helps you *discover* errors
- Reasoning about the rep invariant helps you *avoid* errors
 - Or prove that they do not exist!
 - We will discuss such reasoning later on

The rep invariant constrains structure, not meaning

- New implementation of insert that *preserves the rep invariant*

```
public void insert(Character c) {  
    Character cc = new Character(encrypt(c));  
    if (!elts.contains(cc))  
        elts.addElement(cc);  
}  
  
public boolean member(Character c) {  
    return elts.contains(c);  
}
```

- The program is still wrong
 - Clients observe incorrect behavior
 - Where is the error?
 - We must consider the **meaning**
 - The ***abstraction function*** helps us

```
CharSet s = new CharSet();  
Character a = new  
Character('a');  
s.insert(a);  
if (s.member(a))  
    // print "right";  
else  
    // print "wrong";
```

Abstraction function: $\text{rep} \rightarrow \text{abstract value}$

- The abstraction function maps the concrete representation to the abstract value it represents
 - AF: $\text{Object} \rightarrow \text{abstract value}$
 - $\text{AF}(\text{CharSet this}) = \{ c \mid c \text{ is contained in this.elts} \}$
“set of Characters contained in this.elts”
Typically not executable
- The abstraction function lets us reason about behavior from the client perspective
- Our real goal is to satisfy the specification of insert
 - // modifies: this
 - // effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$
 - public void **insert** (Character c);
- Once again we can identify the problem
 - Applying the abstraction function to the result of the call to insert yields $\text{AF}(\text{elts}) \cup \{\text{encrypt}('a')\}$
 - What if we used this abstraction function?
 - $\text{AF}(\text{this}) = \{ c \mid \text{encrypt}(c) \text{ is contained in this.elts} \}$
 - $\text{AF}(\text{this}) = \{ \text{decrypt}(c) \mid c \text{ is contained in this.elts} \}$

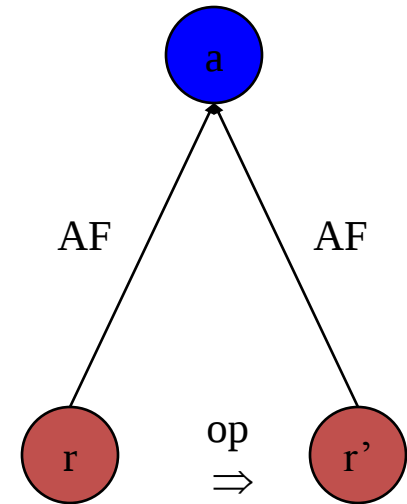
Placing the blame

- Our real goal is to satisfy the **specification of insert**:
 // modifies: this
 // effects: $\text{this}_{\text{post}} = \text{this}_{\text{pre}} \cup \{c\}$
 public void **insert** (Character c);
- The **AF** tells us what the rep means (and lets us place the blame)
 $\text{AF}(\text{CharSet this}) = \{ c \mid c \text{ is contained in this.elts} \}$
- Consider a call to insert:
 On **entry**, the meaning is $\text{AF}(\text{this}_{\text{pre}}) \approx \text{elts}_{\text{pre}}$
 On **exit**, the meaning is $\text{AF}(\text{this}_{\text{post}}) = \text{AF}(\text{this}_{\text{pre}}) \cup \{\text{encrypt('a')}\}$
- What if we used this abstraction function?
 $\text{AF}(\text{this}) = \{ c \mid \text{encrypt}(c) \text{ is contained in this.elts} \}$
 $= \{ \text{decrypt}(c) \mid c \text{ is contained in this.elts} \}$

Benevolent side effects

- Different implementation of member:

```
boolean member(Character c1) {  
    int i = elts.indexOf(c1);  
    if (i == -1)  
        return false;  
    // move-to-front optimization  
    Character c2 = elts.elementAt(0);  
    elts.set(0, c1);  
    elts.set(i, c2);  
    return true;  
}
```



The abstraction function is a function

- Q: Why do we map concrete to abstract rather than vice versa?
- It's not a function in the other direction.
 - Ex: lists $[a,b]$ and $[b,a]$ each represent the set $\{a, b\}$
- It's not as useful in the other direction.
 - Can construct objects via the provided operators

Writing an abstraction function

- The domain: all representations that satisfy the rep invariant
- The range: can be tricky to denote
 - For mathematical entities like sets: easy
 - For more complex abstractions: give them fields
 - AF defines the value of each “specification field”
- The overview section of the specification should provide a way of writing abstract values
 - A printed representation is valuable for debugging

Summary

- Rep invariant
 - Which concrete values represent abstract values
- Abstraction function
 - Which abstract value each concrete value represents
- Together, they modularize the implementation
 - Can examine operators one at a time
 - Neither one is part of the abstraction (the ADT)
- In practice
 - Always write a representation invariant
 - Write an abstraction function when you need it
 - Write an informal one for most non-trivial classes
 - A formal one is harder to write and usually less useful