Graphs IV
Chapter 9, 10.3 in Weiss

CSE 326
Data Structures
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Today's Outline

- **Announcements**
  - **Last Homework! Written Homework #8 due Fri March 1**
- **Today's Topics:**
  - **Graphs**
    - **All-Pairs Shortest Paths**

Graphs

- Representations
- Topological Sort
- Finding paths
  - DFS
  - BFS
  - Dijkstra's
- MST
  - Prim's
  - Kruskal's

Single-Source Shortest Path

- Given a graph **G** = (V, E) and a single distinguished vertex **s**, find the shortest weighted path from **s** to every other vertex in **G**.

All-Pairs Shortest Path:

- Find the shortest paths between all pairs of vertices in the graph.
- How?

Analysis

- Total running time for Dijkstra's:

  $O(|V|^2 + |E|)$ (linear scan)

  $O(|V| \log |V| + |E| \log |V|)$ (heaps)

What if we want to find the shortest path from each point to ALL other points?

Dynamic Programming

Algorithmic technique that systematically records the answers to sub-problems in a table and re-uses those recorded results (rather than re-computing them).

Simple Example: Calculating the Nth Fibonacci number.

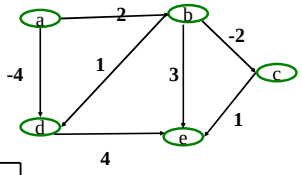
Fib(N) = Fib(N-1) + Fib(N-2)

Floyd-Warshall

```
for (int k = 1; k =< V; k++)
 for (int i = 1; i =< V; i++)
  for (int j = 1; j =< V; j++)
   if ( ( M[i][k]+ M[k][j] ) < M[i][j] )
     M[i][j] =  M[i][k]+ M[k][j]
```

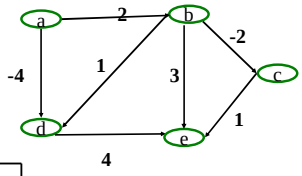
Invariant: After the kth iteration, the matrix includes the shortest paths for all pairs of vertices (i,j) containing only vertices 1..k as intermediate vertices

Floyd-Warshall



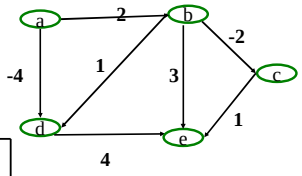
	a	b	c	d	e
a					
b					
c					
d					
e					

Floyd-Warshall



	a	b	c	d	e
a	0				
b		0			
c			0		
d				0	
e					0

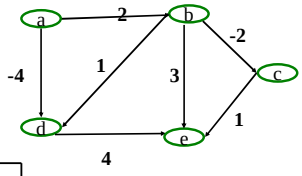
Initial state of the matrix:



	a	b	c	d	e
a	0	2	∞	-4	∞
b	∞	0	-2	1	3
c	∞	∞	0	∞	1
d	∞	∞	∞	0	4
e	∞	∞	∞	∞	0

M[i][j] = min(M[i][j], M[i][k]+ M[k][j])

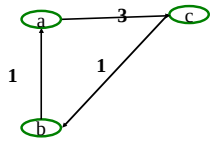
Floyd-Warshall - for All-pairs shortest path



	a	b	c	d	e
a	0	2	0	-4	0
b	-	0	-2	1	-1
c	-	-	0	-	1
d	-	-	-	0	4
e	-	-	-	-	0

Final Matrix Contents

Perform Floyd-Warshall



	a	**b**	**c**
a			
b			
c			

Activity

Transitive Closure

The transitive closure of a graph G=(V,E)
Is the graph G* = (V, E*) where

E* = { (i,j) : there is a path from vertex i to vertex j in G}

"All-pairs reachability"