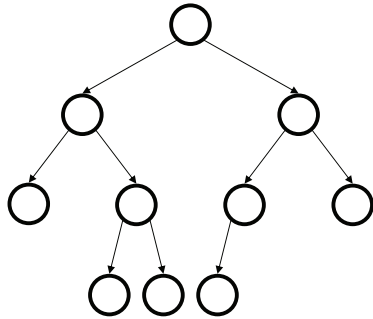


Definition: Null Path Length

Another useful definition:

$npl(x)$ is the height of the largest perfect binary tree that is both itself rooted at x and contained within the subtree rooted at x .



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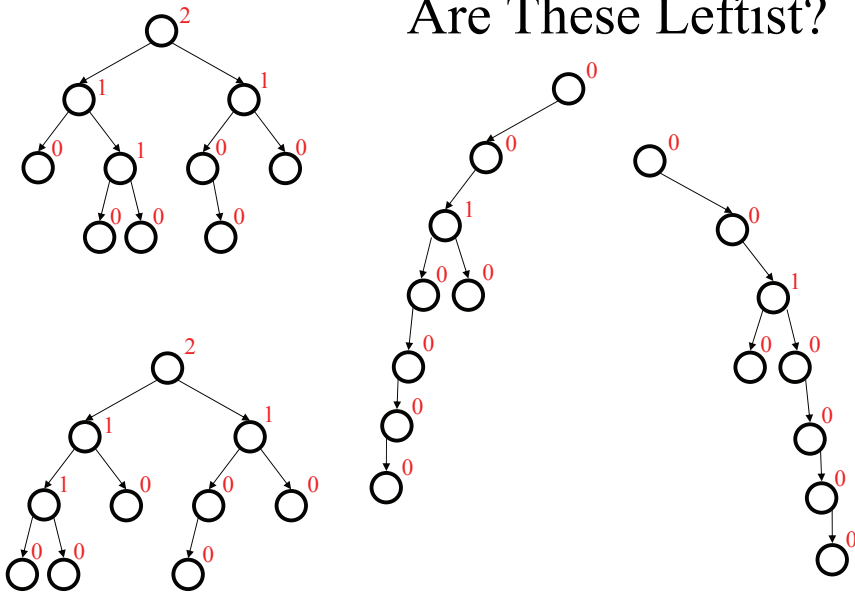
Leftist Heap Properties

- Order property
 - parent's priority value is \leq to children's priority values
 - result: minimum element is at the root
 - (Same as binary heap)
- Structure property
 - For every node x , $npl(\text{left}(x)) \geq npl(\text{right}(x))$
 - result: tree is at least as "heavy" on the left as the right

(Terminology: we will say a leftist heap's tree is a leftist tree.)

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Are These Leftist?



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Observations

Are leftist trees always...

- complete?
- balanced?

Consider a subtree of a leftist tree...

- is it leftist?

Right Path in a Leftist Tree is Short (#1)

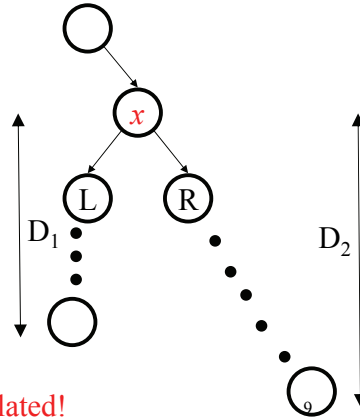
Claim: The right path (path from root to rightmost leaf) is as short as *any* in the tree.

Proof: (By contradiction)

Pick a shorter path: $D_1 < D_2$
 Say it diverges from right path at x

$npl(L) \leq D_1 - 1$ because of the path of length $D_1 - 1$ to null

$npl(R) \geq D_2 - 1$ because every node on right path is leftist



Leftist property at x violated!

Right Path in a Leftist Tree is Short (#2)

Claim: If the right path has r nodes, then the tree has at least $2^r - 1$ nodes.

Proof: (By induction)

Base case : $r=1$. Tree has at least $2^1 - 1 = 1$ node

Inductive step : assume true for $r-1$. Prove for tree with right path at least r .

1. Right subtree: right path of $r-1$ nodes
 $\Rightarrow 2^{r-1} - 1$ right subtree nodes (by induction)
2. Left subtree: also right path of length at least $r-1$ (prev. slide)
 $\Rightarrow 2^{r-1} - 1$ left subtree nodes (by induction)

\Rightarrow Total tree size: $(2^{r-1} - 1) + (2^{r-1} - 1) + 1 = 2^r - 1$

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Why do we have the leftist property?

Because it guarantees that:

- the *right path is really short* compared to the number of nodes in the tree
- A leftist tree of N nodes, has a right path of at most $\log_2(N+1)$ nodes

Idea – perform all work on the right path

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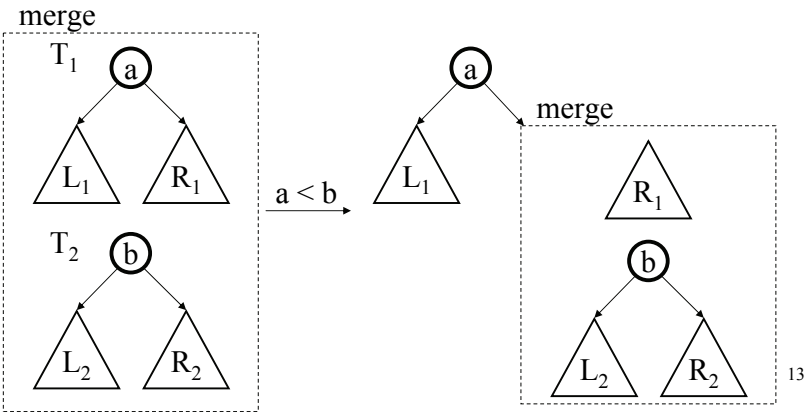
Merge two heaps (basic idea)

- Put the root with smaller value as the new root.
- Hang its left subtree on the left.
- Recursively merge its right subtree and the other tree.
- Before returning from recursion:
 - Update npl of merged root.
 - Swap left and right subtrees just below root, if needed, to keep leftist property of merged result.

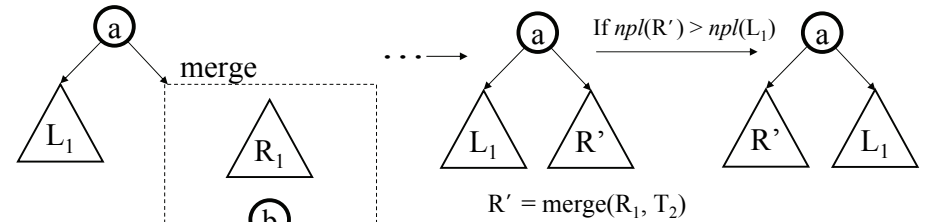
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Merging Two Leftist Heaps

Recursive calls to $\text{merge}(T_1, T_2)$: returns one leftist heap containing all elements of the two (distinct) leftist heaps T_1 and T_2



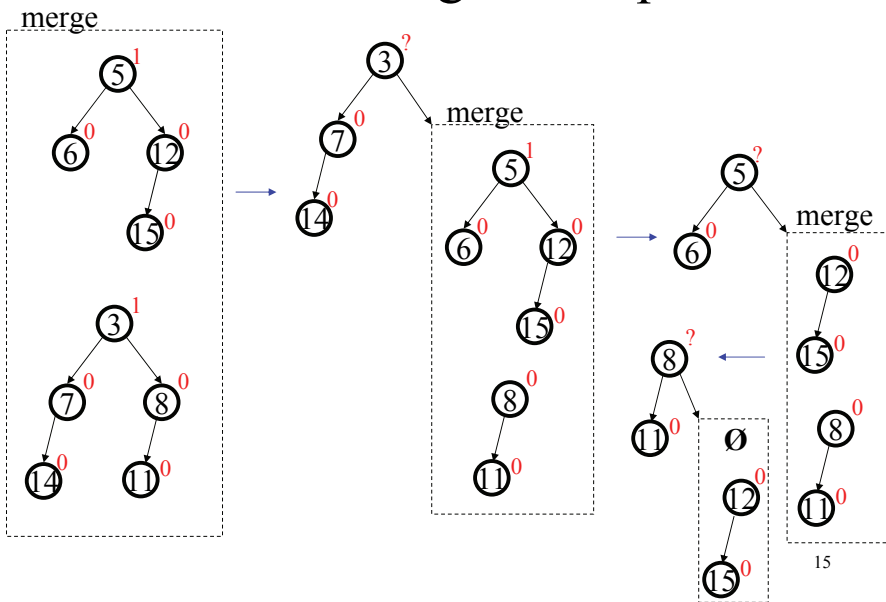
Merge Continued



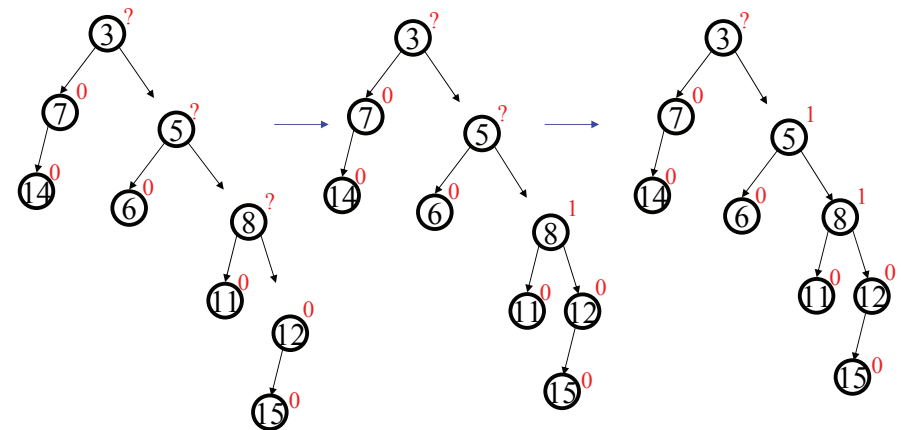
Note special case: $\text{merge}(\text{null}, T) = \text{merge}(T, \text{null}) = T$

runtime:

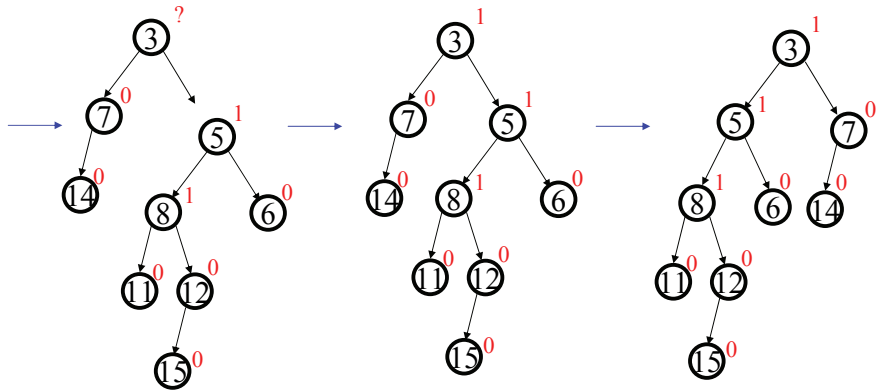
Leftest Merge Example



Sewing Up the Example



Sewing Up the Example



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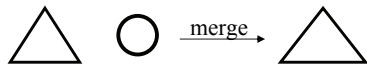
Other Heap Operations

- insert ?
- deleteMin ?

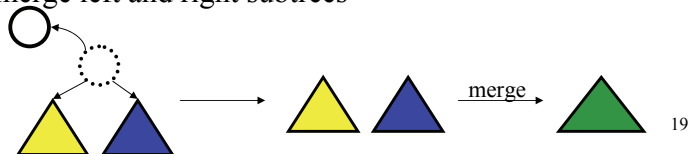
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Operations on Leftist Heaps

- merge with two trees of total size n : $O(\log n)$
- insert with heap size n : $O(\log n)$
 - pretend node is a size 1 leftist heap
 - insert by merging original heap with one node heap



- deleteMin with heap size n : $O(\log n)$
 - remove and return root
 - merge left and right subtrees



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Leftist Heaps: Summary

Good

-
-

Bad

-
-

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