Sorting Lower Bound
Radix Sort

CSE 326
Data Structures
Lecture 16

Reading

- Reading
  - › Sections 7.8-7.11

How fast can we sort?

- Heapsort, Mergesort, and Quicksort all run in O(N log N) best case running time
- Can we do any better?
- No, if the basic action is a comparison.

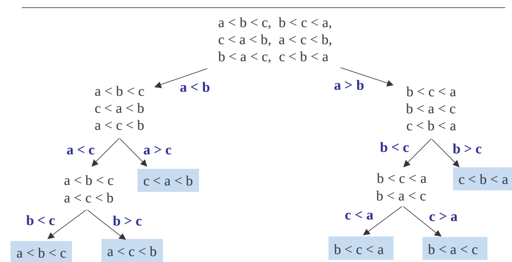
Sorting Model

- Recall our basic assumption: we can only compare two elements at a time
  - › we can only reduce the possible solution space by half each time we make a comparison
- Suppose you are given N elements
  - › Assume no duplicates
- How many possible orderings can you get?
  - › Example: a, b, c (N = 3)

Permutations

- How many possible orderings can you get?
  - › Example: a, b, c (N = 3)
  - › (a b c), (a c b), (b a c), (b c a), (c a b), (c b a)
  - › 6 orderings = 3•2•1 = 3! (ie, “3 factorial”)
  - › All the possible permutations of a set of 3 elements
- For N elements
  - › N choices for the first position, (N-1) choices for the second position, …, (2) choices, 1 choice
  - › N(N-1)(N-2)⋯(2)(1)= N! possible orderings

Decision Tree



The leaves contain all the possible orderings of a, b, c

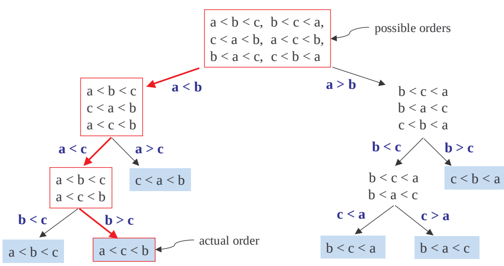
Decision Trees

- A Decision Tree is a Binary Tree such that:
  - › Each node = a set of orderings
    - ie, the remaining solution space
  - › Each edge = 1 comparison
  - › Each leaf = 1 unique ordering
  - › How many leaves for N distinct elements?
    - N!, ie, a leaf for each possible ordering
- Only 1 leaf has the ordering that is the desired correctly sorted arrangement

Decision Trees and Sorting

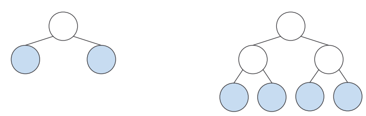
- Every sorting algorithm corresponds to a decision tree
  - › Finds correct leaf by choosing edges to follow
    - ie, by making comparisons
  - › Each decision reduces the possible solution space by one half
- Run time is ≥ maximum no. of comparisons
  - › maximum number of comparisons is the length of the longest path in the decision tree, i.e. the height of the tree

Decision Tree Example



How many leaves on a tree?

- Suppose you have a binary tree of height d . How many leaves can the tree have?
  - › d = 1 à at most 2 leaves,
  - › d = 2 à at most 4 leaves, etc.



Lower bound on Height

- A binary tree of height d has at most 2^d leaves
  - › depth d = 1 à 2 leaves, d = 2 à 4 leaves, etc.
  - › Can prove by induction
- Number of leaves, $L \leq 2^d$
- Height $d \geq \log_2 L$
- The decision tree has N! leaves
- So the decision tree has height $d \geq \log_2(N!)$

log(*N*!) is Ω(*N*log*N*)

$$\log(N!) = \log\big(N \cdot (N-1) \cdot (N-2) \cdots (2) \cdot (1)\big)$$

select just the first N/2 terms

$$= \log N + \log(N-1) + \log(N-2) + \cdots + \log 2 + \log 1$$

$$\geq \log N + \log(N-1) + \log(N-2) + \cdots + \log \frac{N}{2}$$

each of the selected terms is ≥ logN/2

$$\geq \frac{N}{2} \log \frac{N}{2}$$

$$\geq \frac{N}{2}(\log N - \log 2) = \frac{N}{2}\log N - \frac{N}{2}$$

$$= \Omega(N \log N)$$

Ω(N log N)

- Run time of any comparison-based sorting algorithm is **Ω(N log N)**
- Can we do better if we don't use comparisons?

Radix Sort: Sorting integers

- Historically goes back to the 1890 census.
- Radix sort = multi-pass bucket sort of integers in the range 0 to B^P-1
- Bucket-sort from least significant to most significant "digit" (base B)
- Requires P(B+N) operations where P is the number of passes (the number of base B digits in the largest possible input number).
- If P and B are constants then O(N) time to sort!

Radix Sort Example

Input data

478
537
9
721
3
38
123
67

Bucket sort by 1's digit

0	1	2	3	4	5	6	7	8	9
	721		3 123				537 67	478 38	9

After 1st pass

721
3
123
537
67
478
38
9

This example uses B=10 and base 10 digits for simplicity of demonstration. Larger bucket counts should be used in an actual implementation.

Radix Sort Example

After 1st pass

721
3
123
537
67
478
38
9

Bucket sort by 10's digit

0	1	2	3	4	5	6	7	8	9
03 09		721 123	537 38			67	478		

After 2nd pass

3
9
721
123
537
38
67
478

Radix Sort Example

After 2nd pass

3
9
721
123
537
38
67
478

Bucket sort by 100's digit

0	1	2	3	4	5	6	7	8	9
003 009 038 067	123			478	537		721		

After 3rd pass

3
9
38
67
123
478
537
721

Invariant: after k passes the low order k digits are sorted.

Implementation Options

- List
  - › List of data, bucket array of lists.
  - › Concatenate lists for each pass.
- Array / List
  - › Array of data, bucket array of lists.
- Array / Array
  - › Array of data, array for all buckets.
  - › Requires counting.

Array / Array

Index	Data Array	Count Array	Address Array	Target Array
0	478	0	0	721
1	537	1	0	3
2	9	0	1	123
3	721	2	1	537
4	3	0	3	67
5	38	0	3	478
6	123	0	3	38
7	67	2	3	9
8		2	5	
9		1	7	

Target Array buckets: 1 (721), 3 (3, 123), 7 (537, 67), 8 (478, 38), 9 (9)

add[0] := 0
add[i] := add[i-1] + count[i-1], i > 0

Bucket i ranges from add[i] to add[i+1]-1

Array / Array

- Pass 1 (over A)
  - Calculate counts and addresses for 1st "digit"
- Pass 2 (over T)
  - Move data from A to T
  - Calculate counts and addresses for 2nd "digit"
- Pass 3 (over A)
  - Move data from T to A
  - Calculate counts and addresses for 3rd "digit"
- …
- In the end an additional copy may be needed.

Choosing Parameters for Radix Sort

- N number of integers – given
- m bit numbers - given
- B number of buckets
  - $B = 2^r$ – calculations can be done by shifting.
  - N/B not too small, otherwise too many empty buckets.
  - P = m/r should be small.
- Example – 1 million 64 bit numbers. Choose $B = 2^{16} = 65{,}536$. 1 Million / B ≈ 15 numbers per bucket. P = 64/16 = 4 passes.

Properties of Radix Sort

- Not in-place
  - needs lots of auxiliary storage.
- Stable
  - equal keys always end up in same bucket in the same order.
- Fast
  - $B = 2^r$ buckets on m bit numbers

$$O\left(\frac{m}{r}(n+2^r)\right) \text{ time}$$

Internal versus External Sorting

- So far assumed that accessing A[i] is fast – Array A is stored in internal memory (RAM)
  - Algorithms so far are good for internal sorting
- What if A is so large that it doesn't fit in internal memory?
  - Data on disk or tape
  - Delay in accessing A[i] – e.g. need to spin disk and move head

Internal versus External Sorting

- Need sorting algorithms that minimize disk/tape access time
  - External sorting – Basic Idea:
    - Load chunk of data into RAM, sort, store this "run" on disk/tape
    - Use the Merge routine from Mergesort to merge runs
    - Repeat until you have only one run (one sorted chunk)
    - Text gives some examples

Summary of Sorting

- Sorting choices:
  - $O(N^2)$ – Bubblesort, Insertion Sort
  - O(N log N) average case running time:
    - Heapsort: In-place, not stable (read about it).
    - Mergesort: O(N) extra space, stable.
    - Quicksort: claimed fastest in practice but, $O(N^2)$ worst case. Needs extra storage for recursion. Not stable.
  - O(N) – Radix Sort: fast and stable. Not comparison based. Not in-place.