

predicate
 ↓
 $\frac{x \text{ is even}}{\text{Subject} \quad \text{Predicate}}$ proposition
 + function
 → $E(x)$

is composite $C(x)$

is greater than $G(x,y) \quad x > y$

Universe or Domain of Discourse

Quantifiers

Existential \exists "there exists"

Universal \forall "for all"

$\exists x E(x)$ "there is some even nat"

$\forall x E(x)$ "all nats are even"

if universe is finite x_1, x_2, \dots, x_n

$$\exists x E(x) \equiv E(x_1) \vee E(x_2) \vee \dots \vee E(x_n)$$

$$\forall x E(x) \equiv E(x_1) \wedge E(x_2) \wedge \dots \wedge E(x_n)$$

To Show " $\exists x P(x)$ "

find one example x_0 st $P(x_0)$

to show " $\forall x P(x)$ "

to prove that $P(x)$ holds
for every x in domain

to show " $\neg \forall x P(x)$ "

find x_0 in domain st $\neg P(x_0)$

Counterexample

E even \mathbb{N}
 C composite $G(x, y) : x > y$

$\forall x \ E(x)$

$\forall x (E(x) \rightarrow C(x))$

$\forall x ((G(x, 2) \wedge E(x)) \rightarrow C(x))$

Dilemma
why should $\left\{ \begin{array}{l} F \rightarrow T \\ F \rightarrow F \end{array} \right\}$ be true? \times

A common case: a rule with exceptions
"all even numbers are composite, except 0, 2"
"all numbers are composite, except odd or ≤ 2 "

convenient:

$\forall x (\text{Unexceptional}(x) \rightarrow \text{Rule}(x))$

false for exceptions

true for exceptions by \times
(indp. of Rule(x))

Nested Quantifiers

$\left[\forall x (\forall y P(x, y)) \right]$
 $\forall x Q(x)$
 when $Q(x) = \underline{\forall y P(x, y)}$
 $\exists x \exists y P(x, y)$

free variables
↓
bound
variables
↑

For all pairs x, y $P(x, y)$ is true

$$\forall x \forall y P(x, y) \equiv \forall y \forall x P(x, y)$$

$$\exists x (\forall y P(x, y))$$

$$\neq \forall y (\exists x P(x, y))$$

Domain \mathbb{N}

$$\exists x \forall y y < x \leftarrow F$$

$$\forall y \exists x y < x \leftarrow T \quad \text{if } x = y + 1$$

\equiv
two formulas are logically
equivalent if they have the
same truth value for all
domains and all predicates

$$\begin{aligned}
 & \forall x [P(x) \wedge \neg Q(x)] \\
 \equiv & \forall x P(x) \vee \forall x \neg Q(x) \\
 \equiv & \forall x P(x) \vee \underline{\forall y} (\neg Q(y))
 \end{aligned}$$

\neq

$$\begin{aligned}
 & \forall x (P(x) \vee Q(x)) \\
 & \forall x (P(x)) \vee \forall x (Q(x))
 \end{aligned}$$

N
 $P = \text{even}$

$Q = \text{odd}$