CSE 312 Foundations of Computing II

Lecture 19: More Joint Distributions Tail Bounds part I

Midterm

Average: **82.48** Standard Deviation: **17.82** (Median: **87.25**)

- Solutions available on Canvas Pages
- Regrade requests only until Wednesday
	- Look at solutions and then check if a regrade requests is necessary
	- Graders have bounded resources

Review Conditional Expectation

Definition. Let X be a discrete random variable then the **conditional expectation** of X given event A is

$$
\mathbb{E}[X \mid A] = \sum_{x \in \Omega_X} x \cdot P(X = x \mid A)
$$

Notes:

• Can be phrased as a "random variable version"

 $E[X|Y=y]$

• Linearity of expectation still applies here $\mathbb{E}[aX + bY + c | A] = a \mathbb{E}[X | A] + b \mathbb{E}[Y | A] + c$

Review Law of Total Expectation

Law of Total Expectation (event version). Let *X* be a random variable and let events $A_1, ..., A_n$ partition the sample space. Then, $\mathbb{E}[X] = \sum$ \overline{n} $\mathbb{E}[X \mid A_i] \cdot P(A_i)$

Law of Total Expectation (random variable version). Let *X* be a random variable and Y be a discrete random variable. Then,

 $\dot{i}=1$

$$
\mathbb{E}[X] = \sum_{y \in \Omega_Y} \mathbb{E}[X \mid Y = y] \cdot P(Y = y)
$$

- Law of Total Expectation (LTE)
	- Another LTE example
	- Conditional expectation and LTE for continuous RVs
- Tail Bounds
	- Markov's Inequality

Example – Computer Failures (a familiar example)

Suppose your computer operates in a sequence of steps, and that at each step i your computer will fail with probability q (independently of other steps).

Let X be the number of steps it takes your computer to fail.

What is $\mathbb{E}[X]$?

What kind of RV is X ?

Review Geometric RV

- Example: Biased coin: $P(H) = q > 0$ $P(T) = 1 - q$
- $X = #$ of coin flips until first head

$$
P(X=i)=q(1-q)^{i-1}
$$

$$
P(X = i) = q (1 - q)^{i-1}
$$

\n
$$
\mathbb{E}[X] = \sum_{i=1}^{\infty} i \cdot P(X = i) = \sum_{i=1}^{\infty} i \cdot q(1 - q)^{i-1}
$$

\nConverges, so E[X] is finite

 \overline{q}

 $- q$

 \boldsymbol{Q}

 \boldsymbol{q}

 $1 - q$

 \overline{q}

 $- q$) q

 $- q$

 q

 $(1 - q)^2 q$

Can calculate this directly …

 $(1 - q)^3 q$

Direct Analysis – Expectation of Geometric RV

$$
\mathbb{E}[X] = \sum_{i=1}^{\infty} i \cdot q(1-q)^{i-1} = q \sum_{i=1}^{\infty} i(1-q)^{i-1} \quad \text{Converges, so } \mathbb{E}[X] \text{ is finite}
$$
\n
$$
\text{So} \qquad \mathbb{E}[X] = q \left[1 + 2(1-q) + 3(1-q)^2 + \dots + i(1-q)^{i-1} + \dots \right]
$$
\n
$$
\text{Then } (1-q)\mathbb{E}[X] = q \left[(1-q) + 2(1-q)^2 + \dots + (i-1)(1-q)^{i-1} + \dots \right]
$$
\n
$$
\text{Subtracting gives}
$$

$$
q \mathbb{E}[X] = q[1 + (1 - q) + (1 - q)^2 + \dots + (1 - q)^{i-1} + \dots]
$$

$$
q \mathbb{E}[X] = q\left[\frac{1}{1 - (1 - q)}\right] = 1 \qquad \text{since for } 0 < r < 1, \sum_{i=0}^{\infty} r^i = \frac{1}{1 - r}
$$

Therefore $\mathbb{E}[X] = 1/q$

Same examples with the LTE

Suppose your computer operates in a sequence of steps, and that at each step i your computer will fail with probability q (independently of other steps).

Let X be the number of steps it takes your computer to fail. What is $\mathbb{E}[X]$?

Let Y be the indicator random variable for the event of failure (heads) in step 1

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Then by LTE, \mathbb{E}[X] = \mathbb{E}[X | Y = 1] \cdot P(Y = 1) + \mathbb{E}[X | Y = 0] \cdot P(Y = 0)= 1 \cdot q + \mathbb{E}[X | Y = 0] \cdot (1 - q)= q + (1 + \mathbb{E}[X]) \cdot (1 - q) since if Y = 0 experiment
                                                       starting at step 2 looks like 
                                                       original experiment
```
Solving we get $\mathbb{E}[X] = 1/q$

Conditional Expectation again…

Definition. Let X be a discrete random variable; then the **conditional expectation** of X given event A is

$$
\mathbb{E}[X \mid A] = \sum_{x \in \Omega_X} x \cdot P(X = x \mid A)
$$

Therefore for X and Y discrete random variables, the conditional expectation of X given $Y = y$ is

$$
\mathbb{E}[X \mid Y = y] = \sum_{x \in \Omega_X} x \cdot P(X = x \mid Y = y) = \sum_{x \in \Omega_X} x \cdot p_{X \mid Y}(x \mid y)
$$

 $p_{X|Y}(x|y) = P(X = x | Y = y) =$ $p_{X,Y}(x, y)$ where we **define** $p_{X|Y}(x|y) = P(X = x | Y = y) = \frac{1}{p_Y(y)}$

Conditional Expectation – Discrete & Continuous

Discrete: Conditional PMF:
$$
p_{X|Y}(x|y) = \frac{p_{X,Y}(x,y)}{p_Y(y)}
$$

Conditional Expectation: $\mathbb{E}[X | Y = y] = \sum$ $x \in \Omega_X$ $x \cdot p_{X|Y}(x|y)$

Continuous: Conditional PDF:
$$
f_{X|Y}(x|y) = \frac{f_{X,Y}(x,y)}{f_Y(y)}
$$

Conditional Expectation:

$$
\mathbb{E}[X \mid Y = y] = \int_{-\infty}^{\infty} x \cdot f_{X|Y}(x|y) \, dx
$$

Law of Total Expectation - continuous

Law of Total Expectation (event version). Let X be a random variable and let events $A_1, ..., A_n$ partition the sample space. Then, $\mathbb{E}[X] = \sum$ \overline{n} $\mathbb{E}[X \mid A_i] \cdot P(A_i)$

Law of Total Expectation (random variable version). Let *X* and *Y* be continuous random variables. Then,

 $\dot{i}=1$

$$
\mathbb{E}[X] = \int_{-\infty}^{\infty} \mathbb{E}[X \mid Y = y] \cdot f_Y(y) \, dy
$$

Using LTE for Continuous RVs

PDF for
$$
\text{Exp}(\lambda)
$$
 is
$$
\begin{cases} \lambda e^{-\lambda x} & x \ge 0 \\ 0 & \text{o.w.} \end{cases}
$$

Expectation is $1/\lambda$

Suppose that we first choose $Y \sim \text{Exp}(1/2)$ and then choose $X \sim \text{Exp} \left(\frac{1}{\nu} \right)$ $\frac{1}{Y}$. What is $\mathbb{E}[X]$?

$$
f_{X|Y}(x|y) = \begin{cases} (1/y) e^{-(x/y)} & x \ge 0 & y \text{ is fixed here} \\ 0 & 0. W. \end{cases}
$$

$$
\mathbb{E}[X \mid Y = y] = \int_{-\infty}^{\infty} x \cdot f_{X|Y}(x|y) dx = \int_{0}^{\infty} x \cdot (1/y) e^{-(x/y)} dx = y
$$

$$
\mathbb{E}[X] = \int_{-\infty}^{\infty} \mathbb{E}[X \mid Y = y] f_{Y}(y) dy = \int_{0}^{\infty} y \cdot \frac{1}{2} e^{-y/2} dy = 2
$$

Reference Sheet (with continuous RVs)

Brain Break

Agenda

- Law of Total Expectation (LTE)
	- Another LTE example
	- Conditional expectation and LTE for continuous RVs
- Tail Bounds
	- Markov's Inequality

Tail Bounds (Idea)

Bounding the probability that a random variable is far from its mean. Usually statements of the form:

> $P(X \geq k \cdot \mathbb{E}[X]) \leq b$ $P(|X - \mathbb{E}[X]| \ge a) \le b$

Useful tool when

- An approximation that is easy to compute is sufficient
- The process is too complex to analyze exactly

Markov's Inequality

Theorem. Let X be a random variable taking only *non-negative* values. Then, for any $t > 0$,

$$
P(X \ge t) \le \frac{\mathbb{E}[X]}{t}.
$$

(Alternative form) For any $k \geq 1$, $P(X \geq k \cdot \mathbb{E}[X]) \leq \frac{1}{k}$ \boldsymbol{k}

Incredibly simplistic – only requires that the random variable is non-negative and only needs you to know expectation. You don't need to know **anything else** about the distribution of X .

Markov's Inequality – Proof I

 $x \geq t$

Theorem. Let X be a (discrete) random variable taking only non-negative values. Then, for any $t > 0$,

 $\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}.$

$$
\mathbb{E}[X] = \sum_{x} x \cdot P(X = x)
$$

=
$$
\sum_{x \ge t} x \cdot P(X = x) \left(\sum_{x < t} x \cdot P(X = x) \right)
$$

$$
\ge \sum_{x \ge t} x \cdot P(X = x)
$$

$$
\ge \sum t \cdot P(X = x) = t \cdot P(X \ge t)
$$

 ≥ 0 because $x \geq 0$ whenever $P(X = x) \ge 0$ $(X$ takes only non-negative values)

Follows by re-arranging terms

…

Markov's Inequality – Proof II

Theorem. Let X be a (continuous) random variable taking only non-negative values. Then, for any $t > 0$,

 $\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}.$

$$
\mathbb{E}[X] = \int_0^\infty x \cdot f_X(x) dx
$$

=
$$
\int_t^\infty x \cdot f_X(x) dx + \int_0^t x \cdot f_X(x) dx
$$

$$
\geq \int_t^\infty x \cdot f_X(x) dx
$$

$$
\geq \int_t^\infty t \cdot f_X(x) dx = t \cdot \int_t^\infty f_X(x) dx = t \cdot P(X \geq t)
$$

so $P(X \ge t) \le \mathbb{E}[X]/t$ as before

Example – Geometric Random Variable

Let X be geometric RV with parameter p

$$
P(X = i) = (1 - p)^{i-1}p \qquad \mathbb{E}[X] = \frac{1}{p}
$$

"X is the number of times Alice needs to flip a biased coin until she sees heads, if *heads occurs with probability ?*

What is the probability that $X \geq 2\mathbb{E}[X] = 2/p$?

Markov's inequality: $P(X \ge 2 \mathbb{E}[X]) \le \frac{1}{2}$ *

Example

Suppose that the average number of ads you will see on a website is 25. Give an upper bound p on the probability of seeing a website with 75 or more ads.

Example

Suppose that the average number of ads you will see on a website is 25. Give an upper bound p on the probability of seeing a website with 20 or more ads.

Example – Geometric Random Variable

Let X be geometric RV with parameter p

$$
P(X = i) = (1 - p)^{i-1}p \qquad \mathbb{E}[X] = \frac{1}{p}
$$

 X is Next time we will see that we can get better X is sees heads, if **Next time we will see that we can get better tail bounds using variance**

h probability *p*?

What is the probability that $X \geq 2\mathbb{E}[X] = 2/p$?

Markov's inequality: $P(X \ge 2 \mathbb{E}[X]) \le \frac{1}{2}$ *