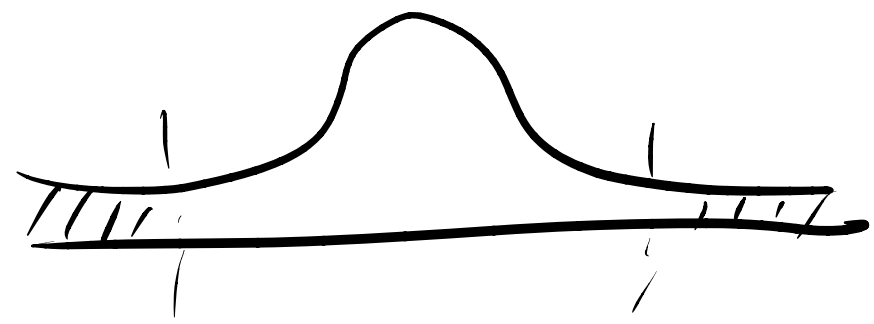


Tail Bounds

CSE 312 Su23
Lecture 18

What's a Tail Bound?



When we were finding our margin of error, we didn't need an exact calculation of the probability.

We needed an inequality: the probability of being outside the margin of error was **at most** 5%.

A tail bound (or concentration inequality) is a statement that bounds the probability in the “tails” of the distribution (says there's very little probability far from the center) or (equivalently) says that the probability is concentrated near the expectation.

Our First bound

one-sided.

Two statements are equivalent.
Left form is often easier to use.
Right form is more intuitive.

Markov's Inequality

Let X be a random variable supported (only) on non-negative numbers. For any $t > 0$

$$\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}$$

Markov's Inequality

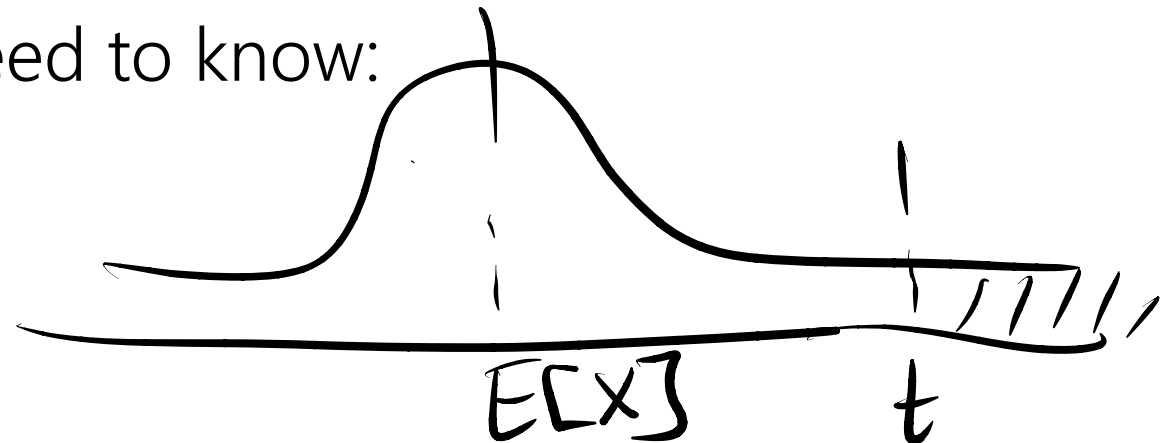
Let X be a random variable supported (only) on non-negative numbers. For any $k > 0$

$$\mathbb{P}(X \geq k\mathbb{E}[X]) \leq \frac{1}{k}$$

Handwritten annotations: an arrow points from k to $k\mathbb{E}[X]$, and a bracket with an arrow points from k to $\frac{1}{k}$.

To apply this bound you only need to know:

1. it's non-negative
2. Its expectation.



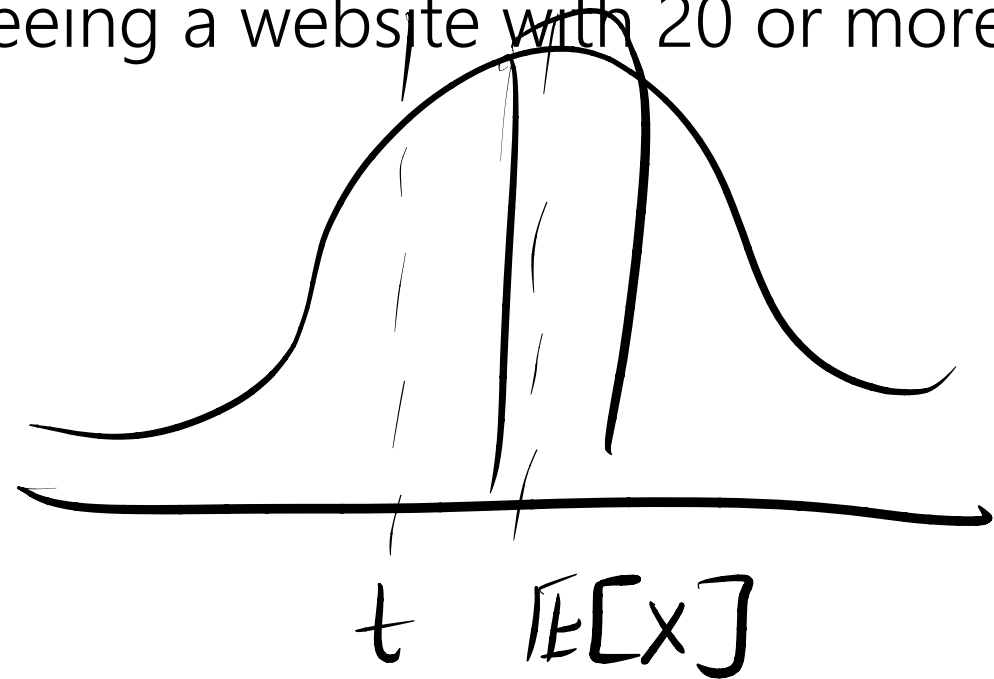
Useless Example

Suppose the average number of ads you see on a website is 25. Give an upper bound on the probability of seeing a website with 20 or more ads.

$$\mathbb{P}(X \geq 20) \leq \frac{\mathbb{E}[X]}{20} = \frac{25}{20} = 1.25$$

Well, that's...true. Technically.

But without more information we couldn't hope to do much better. What if every page gives exactly 25 ads? Then the probability really is 1.



So...what do we do?

A better inequality!

We're trying to bound the tails of the distribution.

What parameter of a random variable describes the tails?

The variance!

Chebyshev's Inequality

Two statements are equivalent.
Left form is often easier to use.
Right form is more intuitive.

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

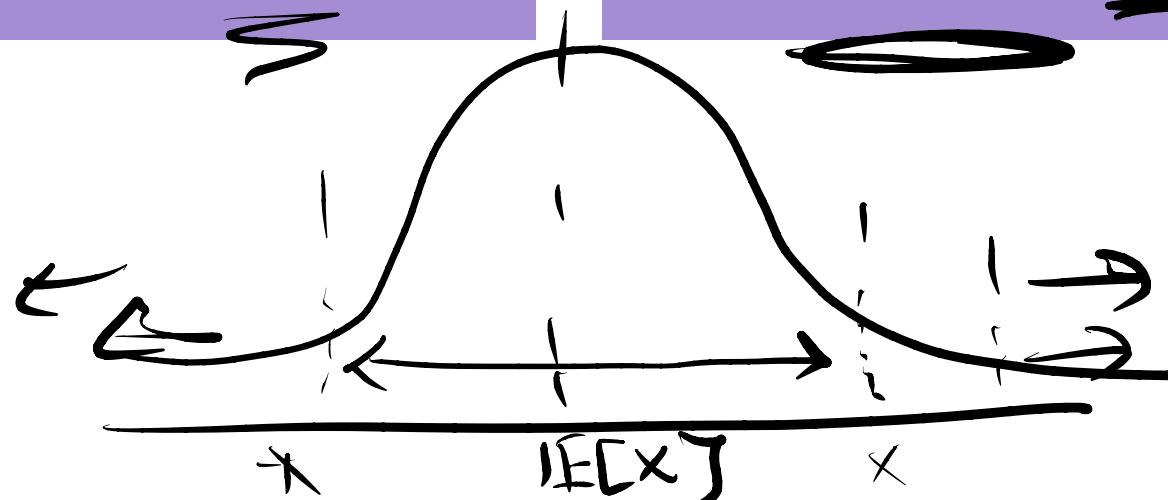
$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

$$t = k \sqrt{\text{Var}(X)}$$

Chebyshev's Inequality

Let X be a random variable. For any $k > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq k\sqrt{\text{Var}(X)}) \leq \frac{1}{k^2}$$



$$\leq \frac{1}{1^2}$$

$$\leq \frac{1}{2^2}$$

Proof of Chebyshev

$$\mathbb{P}(X \geq t) \leq \frac{\mathbb{E}[X]}{t}$$

Markov's

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

★ Let $Z = X - \mathbb{E}[X]$

$$\mathbb{P}(|Z| \geq t) = \mathbb{P}(Z^2 \geq t^2) \leq \frac{\mathbb{E}[Z^2]}{t^2} = \frac{\mathbb{E}[Z^2] - (\mathbb{E}[Z])^2}{t^2} = \frac{\text{Var}(Z)}{t^2} = \frac{\text{Var}(X)}{t^2}$$

LHS

Markov's

LHS, $\text{Var}(Z) = \text{Var}(X - \mathbb{E}[X]) = \text{Var}(X)$

$$\mathbb{E}[Z] = \mathbb{E}[X - \mathbb{E}[X]] = \mathbb{E}[X] - \mathbb{E}[\mathbb{E}[X]] = \mathbb{E}[X] - \mathbb{E}[X] = 0$$

Proof of Chebyshev

$$\frac{\mathbb{E}[Z^2]}{t^2} = \frac{\mathbb{E}[Z^2] - 0^2}{t^2}$$

Let $Z = \underline{X - \mathbb{E}[X]}$

$$= \frac{\mathbb{E}[Z^2] - \mathbb{E}[Z]^2}{t^2}$$

$$\mathbb{P}(|Z| \geq t) = \mathbb{P}(Z^2 \geq t^2) \leq \frac{\mathbb{E}[Z^2]}{t^2} = \frac{\mathbb{E}[Z^2] - (\mathbb{E}[Z])^2}{t^2} = \frac{\text{Var}(Z)}{t^2} = \frac{\text{Var}(X)}{t^2}$$

$$\mathbb{E}[Z] = 0$$

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

$\text{Var}(X+5) = \text{Var}(X)$
 $\text{Var}(X+(-5)) = \text{Var}(X)$

$$\text{Var}(Z) = \text{Var}(X - \mathbb{E}[X]) = \text{Var}(X)$$

Proof of Chebyshev

$$Z = X - \mathbb{E}[X]$$

$$\mathbb{E}[Z] = \mathbb{E}[X - \mathbb{E}[X]]$$

$$= \mathbb{E}[X] - \underbrace{\mathbb{E}[\mathbb{E}[X]]}_{\mathbb{E}[X]} = 0$$

Let $Z = X - \mathbb{E}[X]$

Markov's
Inequality

$$\mathbb{E}[Z] = 0$$

$$\mathbb{P}(|Z| \geq t) = \mathbb{P}(Z^2 \geq t^2) \leq \frac{\mathbb{E}[Z^2]}{t^2} = \frac{\mathbb{E}[Z^2] - (\mathbb{E}[Z])^2}{t^2} = \frac{\text{Var}(Z)}{t^2} = \frac{\text{Var}(X)}{t^2}$$

Inequalities are
equivalent (square
each side).

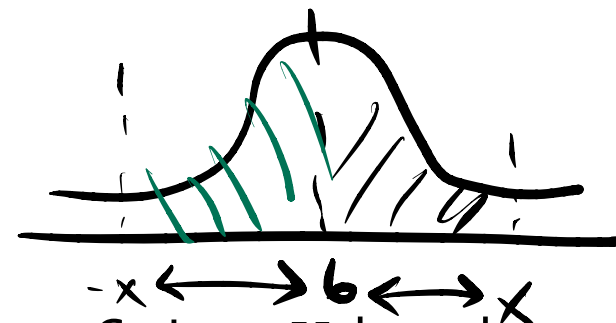
Z is just X shifted.
Variance is
unchanged.

Chebyshev's Inequality

Let X be a random variable. For
any $t > 0$

$$\mathbb{P}(\underbrace{|X - \mathbb{E}[X]|}_{\geq t}) \leq \frac{\text{Var}(X)}{t^2}$$

Example with geometric RV



Suppose you roll a fair (6-sided) die until you see a 6. Let X be the number of rolls.

$$X \sim \text{Geo}\left(\frac{1}{6}\right)$$

Bound the probability that $X \geq 12$,

$$\mathbb{E}[X] = \frac{1}{p} = 6$$

$$\text{Var}(X) = \frac{1-p}{p^2}$$

$$\begin{aligned} & \mathbb{P}(X \geq 12) \\ &= \mathbb{P}(X - 6 \geq 6) \\ &\leq \mathbb{P}(\underbrace{X - 6 \geq 6}_{+(X-6)} + \underbrace{\mathbb{P}(-(X-6) \geq 6)}_{-(X-6)}) \end{aligned}$$

$$\begin{aligned} &= \mathbb{P}(|X - 6| \geq 6) \\ &\leq \frac{\text{Var}(X)}{6^2} \end{aligned}$$

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Example with geometric RV $p = \frac{1}{6}$

Suppose you roll a fair (6-sided) die until you see a 6. Let X be the number of rolls.

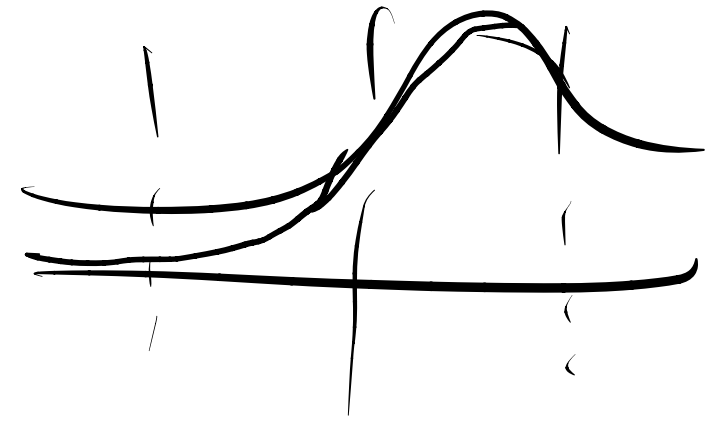
Bound the probability that $X \geq 12$

$$\mathbb{P}(X \geq 12) \leq \mathbb{P}(|X - 6| \geq 6) \leq \frac{\frac{5/6}{1/36}}{6^2} = \frac{5}{6}$$

(Handwritten arrows indicate the substitution of $\frac{5/6}{1/36}$ for $\mathbb{E}[X]$ and 6 for $\text{Var}(X)$ in the Chebyshev inequality formula.)

Not any better than Markov ☹️
(The Markov bound was 0.5)

↗



Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Example with geometric RV

Let X be a geometric rv with parameter p

Bound the probability that $X \geq \frac{2}{p}$

$$\mathbb{P}\left(X - \frac{1}{p} \geq \frac{1}{p}\right)$$

$$\mathbb{P}(X \geq 2/p) \leq \mathbb{P}(|X - 1/p| \geq 1/p) \leq \frac{\frac{1-p}{p^2}}{1/p^2} = 1-p$$

Markov gives:

$$\mathbb{P}\left(X \geq \frac{2}{p}\right) \leq \frac{\mathbb{E}[X]}{2/p} = \frac{1}{p} \cdot \frac{p}{2} = \frac{1}{2} \quad \checkmark$$

For large p , Chebyshev is better.

Geo(p)

$$\mathbb{E}[X] = \frac{1}{p}$$

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Better Example

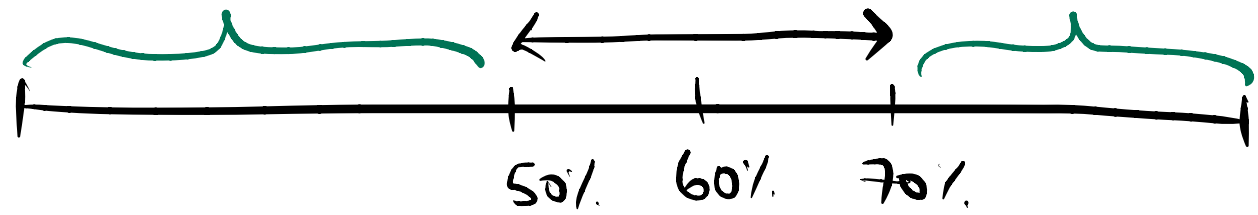
Suppose the average number of ads you see on a website is 25. And the variance of the number of ads is 16. Give an upper bound on the probability of seeing a website with 30 or more ads.

Better Example

Suppose the average number of ads you see on a website is 25. And the variance of the number of ads is 16. Give an upper bound on the probability of seeing a website with 30 or more ads.

$$\mathbb{P}(X \geq 30) \leq \mathbb{P}(|X - 25| \geq 5) \leq \frac{16}{25}$$

Near the mean



Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\bar{X} = \sum X_i / 1000$$

$$\bar{X} = \frac{1}{1000} \sum_{i=1}^{1000} X_i$$

$$\mathbb{E}[\bar{X}] = \frac{1}{1000} \times 1000 \times \mathbb{E}[X_i]$$

$$\text{Var}(\bar{X}) \quad \mathbb{P}(X_i = 1)$$

// //

0.6

$$\frac{1}{1000^2} \cdot 1000 \cdot \text{Var}(X_i)$$

0.6 · 0.4

Chebyshev's Inequality

Let X be a random variable. For

any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Near the mean

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\text{Var}(aX) = a^2 \text{Var}(X)$$

$$\bar{X} = \sum X_i / 1000$$

$$\mathbb{E}[\bar{X}] = 1000 \cdot \frac{.6}{1000} = \frac{3}{5}$$

$$\text{Var}(\bar{X}) = 1000 \frac{.6 \cdot .4}{1000^2} = \frac{3}{12500}$$

$$\text{Var}\left(\frac{1}{1000} X\right)$$

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Near the mean

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\bar{X} = \sum X_i / 1000$$

$$\mathbb{E}[\bar{X}] = 1000 \cdot \frac{.6}{1000} = \frac{3}{5}$$

$$\text{Var}(\bar{X}) = 1000 \frac{.6 \cdot .4}{1000^2} = \frac{3}{12500}$$

$$\mathbb{P}(\underbrace{|\bar{X} - \mathbb{E}[\bar{X}]|}_{\leq} \geq .1) \leq \frac{3/12500}{.1^2} = .024$$

2.4%

Chebyshev's Inequality

Let X be a random variable. For any $t > 0$

$$\mathbb{P}(|X - \mathbb{E}[X]| \geq t) \leq \frac{\text{Var}(X)}{t^2}$$

Chebyshev's – Repeated Experiments

How many coin flips (each head with probability p) are needed until you get n heads.

Let X be the number necessary. What is probability $X \geq 2n/p$?

Markov

$$\frac{\text{Var}(X)}{t^2}$$

$$\downarrow \text{Var}(x)$$

Chebyshev

Chebyshev's – Repeated Experiments

How many coin flips (each head with probability p) are needed until you get n heads.

Let X be the number necessary. What is probability $X \geq 2n/p$?

Markov
$$\mathbb{P}\left(X \geq \frac{2n}{p}\right) \leq \frac{n/p}{2n/p} = \frac{1}{2}$$

Chebyshev
$$\mathbb{P}\left(X \geq \frac{2n}{p}\right) \leq \mathbb{P}\left(\left|X - \frac{n}{p}\right| \geq \frac{n}{p}\right) \leq \frac{\text{Var}(X)}{n^2/p^2} = \frac{n(1-p)/p^2}{n^2/p^2} = \frac{1-p}{n}$$

Takeaway

Chebyshev gets more powerful as the variance shrinks.

Repeated experiments are a great way to cause that to happen.

Markov

$E[X]$

non-neg X

Chebyshev

Symmetric

$\text{Var}(X)$

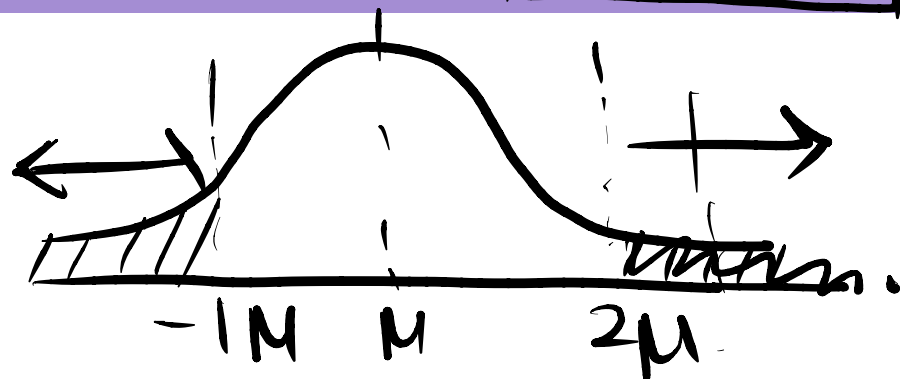
More Assumptions $\xrightarrow{e^{-\frac{\delta^2 \mu}{3}}}$ Better Guarantee $e^{-\frac{\delta^2 \mu}{2}}$

(Multiplicative) Chernoff Bound

Let $X_1, X_2, \dots, X_n$ be independent Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = \mathbb{E}[X]$. For any $0 \leq \delta \leq 1$

$$\mathbb{P}(X \geq (1 + \delta)\mu) \leq \exp\left(-\frac{\delta^2 \mu}{3}\right) \text{ and } \mathbb{P}(X \leq (1 - \delta)\mu) \leq \exp\left(-\frac{\delta^2 \mu}{2}\right)$$



$$(1 + \delta)\mu$$

$$\delta = 1$$

$$e^{-\frac{\delta^2 \mu}{2}}$$

Left

Same Problem, New Solution

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

(Multiplicative) Chernoff Bound

Let $X_1, X_2, \dots, X_n$ be *independent* Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = \mathbb{E}[X]$. For any $0 \leq \delta \leq 1$

$$\mathbb{P}(X \geq (1 + \delta)\mu) \leq \exp\left(-\frac{\delta^2\mu}{3}\right) \text{ and } \mathbb{P}(X \leq (1 - \delta)\mu) \leq \exp\left(-\frac{\delta^2\mu}{2}\right)$$

Right Tail

$$P(\bar{X} \geq 0.6 + 0.1) \quad E[X] = E\left[\sum_{i=1}^{1000} X_i\right]$$

$$P\left(\frac{1}{1000} \sum_{i=1}^{1000} X_i \geq 0.7\right) \quad 1000 \cdot 0.6 = 600$$

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

Want $P\left(\frac{\bar{X}}{1000} \geq .7\right) =$

$$P(X \geq 0.7 \times 1000)$$

$$P(X \geq \boxed{700})$$

$$X \geq (1 + \delta) 600$$

$$(1 + \delta) 600 = 700 \quad \begin{aligned} 1 + \delta &= \frac{7}{6} \\ \delta &= \frac{1}{6} \end{aligned}$$

$$\star P(X \geq (1 + \frac{1}{6}) 600) \leq e^{-\frac{(\frac{1}{6})^2 600}{3}}$$

Chernoff Bound (right tail)

Let $X_1, X_2, \dots, X_n$ be *independent* Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = E[X]$. For any $0 \leq \delta \leq 1$

$$P(X \geq \boxed{(1 + \delta)\mu}) \leq \exp\left(-\frac{\delta^2 \mu}{3}\right)$$

Right Tail

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\text{Want } \mathbb{P}\left(\frac{X}{1000} \geq .7\right) = \mathbb{P}(X \geq .7 \cdot 1000)$$

$$= \mathbb{P}(X \geq (1 + .1/.6) \cdot (.6 \cdot 1000))$$

$$\text{So } \delta = \frac{1}{6} \text{ and } \mu = .6 \cdot 1000$$

$$\mathbb{P}(X \geq 700) \leq \exp\left(-\frac{\frac{1}{6^2} \cdot .6 \cdot 1000}{3}\right)$$

$$\leq \underline{0.0039}$$

Chernoff Bound (right tail)

Let $X_1, X_2, \dots, X_n$ be *independent* Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = \mathbb{E}[X]$. For any $0 \leq \delta \leq 1$

$$\mathbb{P}(X \geq (1 + \delta)\mu) \leq \exp\left(-\frac{\delta^2 \mu}{3}\right)$$

Left Tail

$$\underbrace{\mathbb{P}(\bar{X} \leq 0.6 - 0.1)}_{\mathbb{P}} + \underbrace{\mathbb{P}(\bar{X} \geq 0.6 + 0.1)}$$

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\text{Want } \mathbb{P}\left(\frac{X}{1000} \leq .5\right) = \mathbb{P}(X \leq .5 \cdot 1000)$$

$$\underline{\underline{\mathbb{P}(X \leq 500)}}$$

$$(1 - \delta) 600 = 500$$

$$1 - \delta = \frac{5}{6}$$

$$\frac{1}{6} = \delta$$

$$\mathbb{P}(X \leq (1 - \frac{1}{6}) 600)$$

$$\leq e^{\left(\frac{-(\frac{1}{6})^2 \cdot 600}{2}\right)}$$

Chernoff Bound (left tail)

Let $X_1, X_2, \dots, X_n$ be *independent* Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = \mathbb{E}[X]$. For any $0 \leq \delta \leq 1$

$$\mathbb{P}(X \leq (1 - \delta)\mu) \leq \exp\left(-\frac{\delta^2 \mu}{2}\right)$$

Left Tail

Suppose you run a poll of 1000 people where in the true population 60% of the population supports you. What is the probability that the poll is not within 10-percentage-points of the true value?

$$\text{Want } \mathbb{P}\left(\frac{X}{1000} \leq .5\right) = \mathbb{P}(X \leq .5 \cdot 1000)$$

$$= \mathbb{P}(X \leq (1 - .1/.6) \cdot (.6 \cdot 1000))$$

$$\text{So } \delta = \frac{1}{6} \text{ and } \mu = .6 \cdot 1000$$

$$\mathbb{P}(X \leq 500) \leq \exp\left(-\frac{\frac{1}{6^2} \cdot .6 \cdot 1000}{2}\right)$$

$$\leq 0.0003$$

Chernoff Bound (left tail)

Let $X_1, X_2, \dots, X_n$ be *independent* Bernoulli random variables.

Let $X = \sum X_i$, and $\mu = \mathbb{E}[X]$. For any $0 \leq \delta \leq 1$

$$\mathbb{P}(X \leq (1 - \delta)\mu) \leq \exp\left(-\frac{\delta^2 \mu}{2}\right)$$

Both Tails

Let E be the event that X is not between 500 and 700 (i.e. we're not within 10 percentage points of the true value)

$$\begin{aligned}\mathbb{P}(E) &= \mathbb{P}(X < 500) + \mathbb{P}(X > 700) \\ &\leq .0039 + .0003 = .0042\end{aligned}$$

mutually excl.

Less than 1%. That's a better bound than Chebyshev gave (2.4%)!

Wait a Minute

I asked Wikipedia about the “Chernoff Bound” and I saw something different?

This is the “easiest to use” version of the bound. If you need something more precise, there are other versions.

Wait a Minute

Why are the tails different??

The strongest/original versions of “Chernoff bounds” are symmetric ($1 + \delta$ and $1 - \delta$ correspond), but those bounds are ugly and hard to use.

When computer scientists made the “easy to use versions”, they needed to use some inequalities. The numerators now have plain old δ ’s, instead of $1 +$ or $1 -$. As part of the simplification to this version, there were different inequalities used so you don’t get exactly the same expression.

$$(1 - \delta) (1 + \delta)$$

Wait a Minute

This is just a binomial!

Well if all the X_i have the same probability. It does work if they're independent but have different distributions. But there's bigger reasons to care...

The concentration inequality will let you control n easily, even as a **variable**. That's not easy with the binomial.

What happens when n gets **big**?

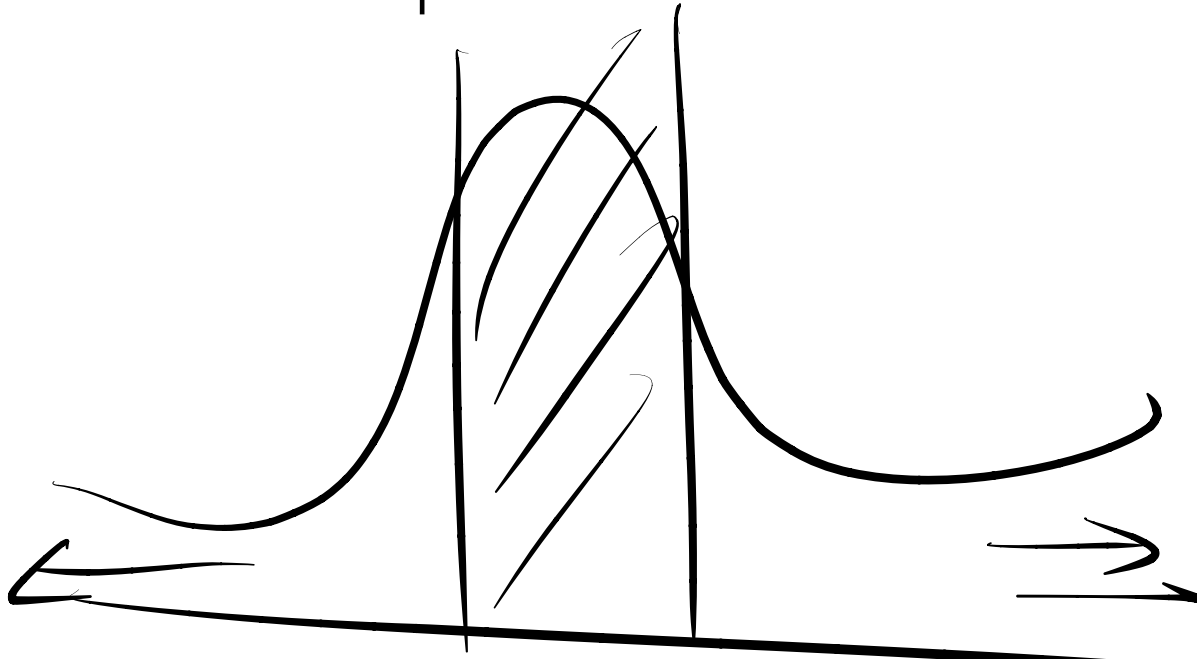
Evaluating $\binom{20000}{10000} \cdot 51^{10000} \cdot 49^{10000}$ is fraught with chances for floating point error and other issues. Chernoff is much better.



But Wait! There's More

For this class, please limit yourself to:
Markov, Chebyshev, and Chernoff, as stated in these slides...

Concentration of measure phenomenon: as n increases, probability that you're near expectation increases



But Wait! There's More

$$(1+\delta)\mu$$

$$\mu + \delta$$

But for your information. There's more.

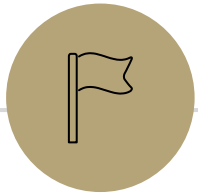
Trying to apply Chebyshev, but only want a "one-sided" bound (and tired of losing that ~~almost~~-factor-of-two) Try [Cantelli's Inequality](#) ★

In a position to use Chernoff, but want additive distance to the mean instead of multiplicative? [They got one of those.](#) ★

Have a sum of independent random variables that aren't indicators, but are bounded, you better believe [Wikipedia's got one](#) ★

Have a sum of random **matrices** instead of a sum of random numbers. Not only is that a thing you can do, but the eigenvalue of the matrix [concentrates](#) ★

There's [a whole book](#) of these!



One More Bound



One More Bound

The Union bound (not a concentration bound)

Union Bound

For any events E, F

$$\mathbb{P}(E \cup F) \leq \mathbb{P}(E) + \mathbb{P}(F)$$

Proof? $\mathbb{P}(E \cup F) = \mathbb{P}(E) + \mathbb{P}(F) - \mathbb{P}(E \cap F)$

And $\mathbb{P}(E \cap F) \geq 0$.

inclusion-exclusion



Concentration Applications

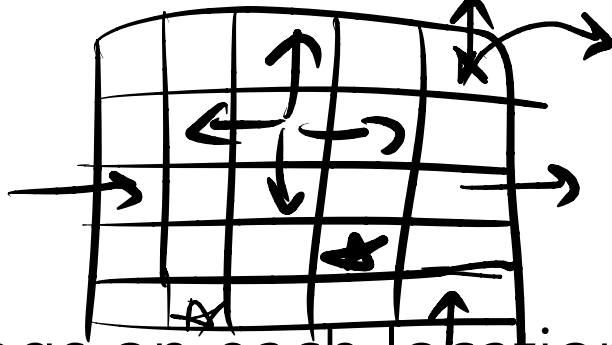
A common pattern:

Figure out “what could possibly go wrong” – often these are dependent.

Use a concentration inequality for each of the things that could go wrong.

Union bound over everything that could go wrong.

Frogs



20 x 25 frogs

There are 20 frogs on each location in a 5x5 grid. Each frog will independently jump to the left, right, up, down, or stay where it is with equal probability. A frog at an edge of the grid magically warps to the corresponding edge (pac-man-style).

Bound the probability that at least one square ends up with at least 36
frogs.

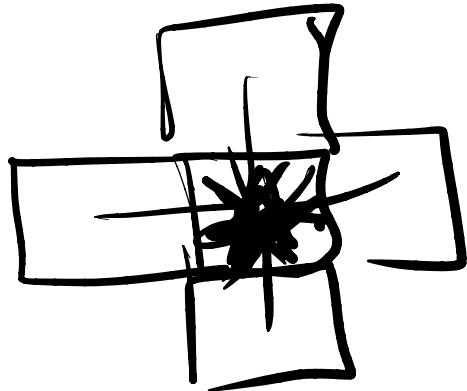
These events are dependent – adjacent squares affect each other!

Frogs

There are 20 frogs on each location in a 5x5 grid. Each frog will independently jump to the left, right, up, down, or stay where it is with equal probability. A frog at an edge of the grid magically warps to the corresponding edge (pac-man-style).

Bound the probability that at least one square ends up with at least 36 frogs.

$$\mathbb{P}(S \geq 36)$$



$$20 \cdot 5 = \underline{\underline{100}}$$

$$\frac{1}{5}$$

Frogs

For an arbitrary location:

There are 100 frogs who could end up there (those above, below, left, right, and at that location). Each with probability .2. Let X be the number that land at the location we're interested in.

Frogs

S_1

$$P(S_1 \cup S_2 \cup S_3 \cup \dots \cup S_{25})$$

$$P(A \cup B) \quad P\left(\bigcup_{i=1}^{25} S_i\right) \leq \sum_{i=1}^{25} P(S_i)$$

For an arbitrary location:

There are 100 frogs who could end up there (those above, below, left, right, and at that location). Each with probability .2. Let X be the number that land at the location we're interested in.

$$P(X \geq 36) = P(X \geq (1 + \delta)20) \leq \exp\left(-\frac{\left(\frac{4}{5}\right)^2 \cdot 20}{3}\right) \leq 0.015 \quad \text{Chernoff} = 25 \cdot 0.015 = 1.5\%$$

$$X = \sum_{i=1}^{100} X_i$$

$$36 = (1 + \delta)20$$

$$\frac{36}{20} - 1 = \frac{16}{20} = \delta = \frac{4}{5}$$

$$X_i \sim \text{Ber}\left(\frac{1}{5}\right)$$

$$E[X] = np = \frac{100}{5} = 20$$

Frogs

For an arbitrary location:

There are 100 frogs who could end up there (those above, below, left, right, and at that location). Each with probability .2. Let X be the number that land at the location we're interested in.

$$\mathbb{P}(X \geq 36) = \mathbb{P}(X \geq (1 + \delta)20) \leq \exp\left(-\frac{\left(\frac{4}{5}\right)^2 \cdot 20}{3}\right) \leq 0.015$$

There are 25 locations. Since all locations are symmetric, by the union bound the probability of at least one location having 36 or more frogs is at most $25 \cdot 0.015 \leq 0.375$.



Probabilistic Algorithms

1. Las Vegas Algorithm:

- > Always correct
- > Varying amounts of time 😊💧

2. Monte Carlo Algorithm:

- > Correct, with some probability 😊💧
- > Bounded runtime

Tail Bounds – Takeaways

Useful when an experiment is complicated and you just need the probability to be small (you don't need the exact value).

Choosing a minimum n for a poll – don't need exact probability of failure, just to make sure it's small.

Designing probabilistic algorithms – just need a guarantee that they'll be extremely accurate

Learning more about the situation (e.g. learning variance instead of just mean, knowing bounds on the support of the starting variables) usually lets you get more accurate bounds.