## Lecture 14: Markov and Chebyshev's Inequalities

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We continue our discussion of Variance, and talk about using it to prove our first concentration bound.

In the last lecture, we discussed how to calculate the variance of a random variable $\operatorname{Var}[X]=\mathbb{E}\left[X^{2}\right]$. We showed that the variance of the sum of two independent random variables is just the sum of the variance. Another nice way to calculate the variance is:

Fact 1. $\operatorname{Var}[X]=\mathbb{E}\left[X^{2}\right]-\mathbb{E}[X]^{2}$.
Proof. Let $\mu=\mathbb{E}[X]$. Then

$$
\begin{array}{rlr}
\operatorname{Var}[X] & =\mathbb{E}\left[(X-\mu)^{2}\right] \\
& =\mathbb{E}\left[X^{2}+\mu^{2}-2 X \mu\right] \\
& =\mathbb{E}\left[X^{2}\right]+\mu^{2}-2 \mu \mathbb{E}[X] & \text { by linearity of expectation } \\
& =\mathbb{E}\left[X^{2}\right]+\mu^{2}-2 \mu^{2}=\mathbb{E}\left[X^{2}\right]-\mu^{2}
\end{array}
$$

Definition 2. The standard deviation of the random variable $X$ is defined to be

$$
\sigma(X)=\sqrt{\operatorname{Var}[X]} .
$$

## Example: $n$ coin tosses

Suppose we toss a fair coin $n$ times. What is the variance of the number of heads?

Let $X$ denote the number of heads. As usual, it is best to represent $X$ using the indicator random variables for each coin toss being a heads:

$$
X_{i}= \begin{cases}1 & \text { if } i^{\prime} \text { th toss is heads } \\ 0 & \text { otherwise }\end{cases}
$$

Then we have $X=X_{1}+X_{2}+\ldots+X_{n}$. Since $X_{1}, \ldots, X_{n}$ are independent, we have

$$
\operatorname{Var}[X]=\operatorname{Var}\left[X_{1}\right]+\operatorname{Var}\left[X_{2}\right]+\ldots+\operatorname{Var}\left[X_{n}\right]
$$

Now we can calculate:

$$
\operatorname{Var}\left[X_{i}\right]=\mathbb{E}\left[X_{i}^{2}\right]-\mathbb{E}\left[X_{i}\right]^{2}
$$

First let us calculate $\mathbb{E}\left[X_{i}^{2}\right]$,

$$
\begin{gathered}
\mathbb{E}\left[X_{i}^{2}\right]=(1 / 2) \cdot 1^{2}+(1 / 2) \cdot 0=1 / 2 \\
\mathbb{E}\left[X_{i}\right]^{2}=((1 / 2) \cdot 1+(1 / 2) \cdot 0)^{2}=(1 / 2)^{2}=1 / 4
\end{gathered}
$$

so we get

$$
\operatorname{Var}\left[X_{i}\right]=\mathbb{E}\left[X_{i}^{2}\right]-\mathbb{E}\left[X_{i}\right]^{2}=1 / 2-1 / 4=1 / 4
$$

So, we conclude that

$$
\operatorname{Var}[X]=\operatorname{Var}\left[X_{1}\right]+\operatorname{Var}\left[X_{2}\right]+\ldots+\operatorname{Var}\left[X_{n}\right]=n / 4
$$

The standard deviation is

$$
\sigma(X)=\sqrt{n} / 2
$$

The variance gives a powerful way to measure the probability that a random variable deviates from its expectation by a lot. As we have seen, $\mathbb{E}[X]$ does not tell us anything about how far $X$ can be from its expectation. However, we do have the following simple inequality:

Fact 3 (Markov's inequality). If $X$ is a non-negative random variable, then

$$
p(X \geq \alpha) \leq \frac{\mathbb{E}[X]}{\alpha}
$$

Proof. We can use conditional expectation to express:

$$
\begin{aligned}
\mathbb{E}[X] & =p(X \geq \alpha) \cdot \mathbb{E}[X \mid X \geq \alpha]+p(X<\alpha) \cdot \mathbb{E}[X \mid X<\alpha] \\
& \geq p(X \geq \alpha) \cdot \mathbb{E}[X \mid X \geq \alpha] \\
& \geq \alpha \cdot p(X \geq \alpha)
\end{aligned}
$$

proving that

$$
p(X \geq \alpha) \leq \frac{\mathbb{E}[X]}{\alpha}
$$

Applying Markov's inequality to the variance gives us Chebyshev's inequality:

Fact 4 (Chebyshev's inequality).

$$
p(|X-\mathbb{E}[X]| \geq \alpha) \leq \frac{\operatorname{Var}[X]}{\alpha^{2}}
$$

Proof.

$$
\begin{aligned}
p(|X-\mathbb{E}[X]| \geq \alpha) & =p\left((X-\mathbb{E}[X])^{2} \geq \alpha^{2}\right) \\
& \leq \frac{\mathbb{E}\left[(X-\mathbb{E}[X])^{2}\right]}{\alpha^{2}} \\
& =\frac{\operatorname{Var}[X]}{\alpha^{2}}
\end{aligned}
$$

by Markov's inequality applied to the non-negative random variable $(X-\mathbb{E}[X])^{2}$.

Let us apply Markov and Chebyshev's inequality to some common distributions.

## Example: Bernoulli Distribution

The Bernoulli distribution is the distribution of a coin toss that has a probability $p$ of giving heads. Let $X$ denote the number of heads. Then we have

$$
\begin{gathered}
\mathbb{E}[X]=p \\
\operatorname{Var}[X]=p-p^{2}
\end{gathered}
$$

Markov's inequality gives

$$
p(X=1)=p(X \geq 1) \leq \frac{\mathbb{E}[X]}{1}=p
$$

Chebyshev's inequality gives

$$
p(X=1)=p(|X-p| \geq 1-p) \leq \frac{p-p^{2}}{(1-p)^{2}}=\frac{p}{1-p}
$$

Neither of these inequalities are very interesting, because it is much easier to see directly that $p(X=1)=p$.

## Example: Binomial Distribution

The binomial distribution is the outcome of tossing $n$ independent coins, each of which gives heads with probability $p$. Let $X$ be the number of heads.

Suppose that $p<1 / 2$. We have:

$$
\begin{gathered}
\mathbb{E}[X]=p n \\
\operatorname{Var}[X]=\left(p-p^{2}\right) n
\end{gathered}
$$

Markov's inequality gives

$$
p(X \geq 2 p n) \leq \frac{\mathbb{E}[X]}{2 p n}=\frac{p n}{2 p n}=\frac{1}{2}
$$

Chebyshev's inequality gives

$$
\begin{aligned}
p(X \geq 2 p n) & =p(|X-n p| \geq p n)) \\
& \leq \frac{\operatorname{Var}[X]}{p^{2} n^{2}}=\frac{p(1-p) n}{p^{2} n^{2}}=\frac{1-p}{p n}
\end{aligned}
$$

We see that $n$ gets large, this probability goes to 0 .

These are much more interesting inequalities, because it is hard to calculate $p(X \geq 2 p n)$ directly.

