

## Randomized Algorithms

Primality testing: Given a number  $n \geq 2$ , is it prime?

If  $n > 2$  and prime,  $n-1 = 2^s \cdot d$   
where  $d$  is odd.

Then:

$$(1) \quad a^d \not\equiv 1 \pmod{n} \quad \forall a \not\equiv 1 \pmod{n}$$

$$(2) \quad a^{2^r \cdot d} \not\equiv -1 \pmod{n} \quad \forall a$$

$r = 1, 2, \dots, \log_2 n$

Choose  $a \in \{2, 3, \dots, n-1\}$   
unif. at random and test that  
(1) and (2) hold.

If  $n$  is not prime:  

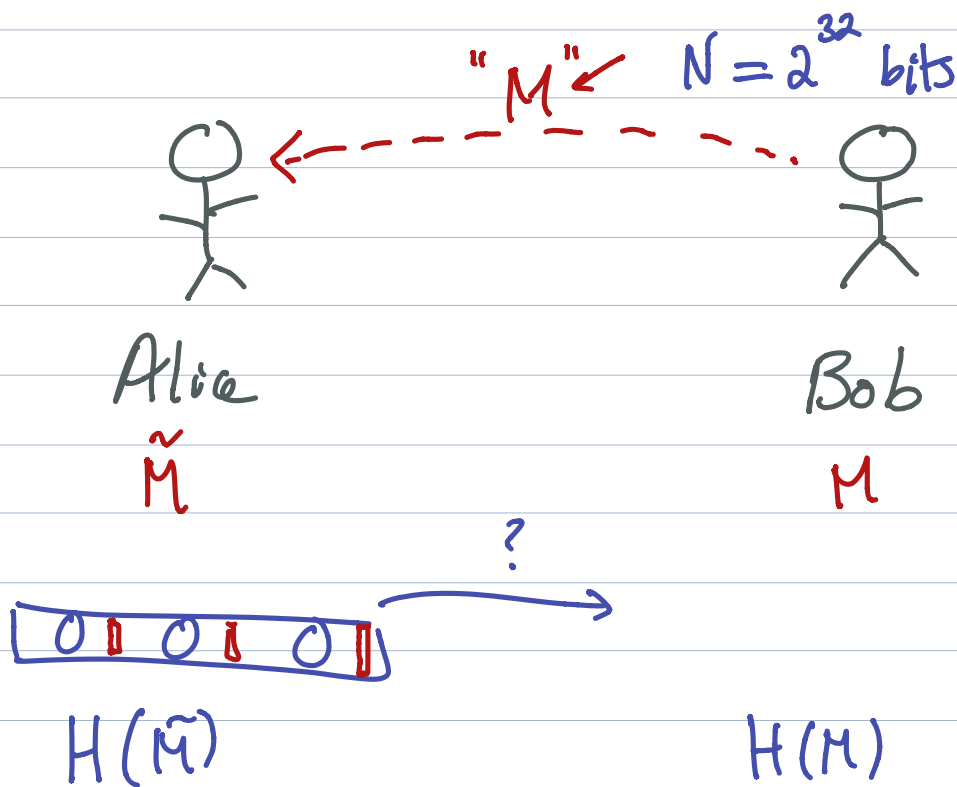
$$\mathbb{P}_a \left[ \begin{array}{l} (1) \text{ or } (2) \\ \text{is violated} \end{array} \right] \geq \frac{1}{4}$$

$n$  not prime  
 $\Rightarrow \mathbb{P} \left[ \begin{array}{l} \text{some } a_i \text{ violates} \\ (1) \text{ or } (2) \end{array} \right]$   
 $a_1, a_2, \dots, a_r \in \{0, 1, \dots, n-1\}$

indep.

$$P[\text{don't find a violation}] \leq \left(\frac{3}{4}\right)^k \quad k=20$$

"correctness amplification"



Have Alice generate a random prime  $p$  on  $\approx \log_2 N$  bits

$$H(x) = x \bmod p$$

$$H(x) = H(x') \iff p \mid x - x'$$

So if  $x \neq x'$ , then

$$P[H(x) \neq H(x')]$$

$$\leq \frac{\# \text{ primes that divide } x-x'}{\# \text{ } c \log_2 N \text{-bit primes}}$$

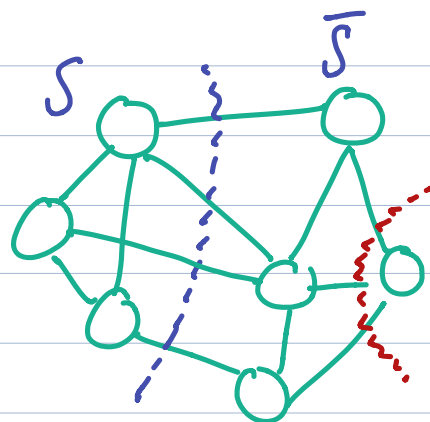
$$\leq \frac{\log_2 N}{c \log_2 N / \log \log N}$$

Choosing  $c \approx \log \log N$  ( $< 10$  in practice)  
gives small error probability.

## Min-cut problem:


Input:  $G$  undirected graph

Output: Minimum cut in  $G$



Algorithm:

- Choose an edge  $e$  uniformly at random in  $G$ .
- $G \leftarrow G/e$

Do this until the graph has two vertices left.  
Return the unique cut remaining. 

Claim: Let  $S^*$  be a minimum cut in  $G$ .

$$\mathbb{P}[\text{output } S^*] \geq \frac{2}{n(n-1)}$$

where  $n = \#$  vertices in  $G$ .

Suppose  $S^*$  has  $k$  edges



$$\mathbb{P} \left[ \begin{array}{l} \text{contract an edge} \\ \text{of } S^* \\ \text{in one step} \end{array} \right] = \frac{k}{\text{total \#edges}}$$

$$\# \text{edges} \geq \frac{kn}{2}$$

$$\leq \frac{k}{\frac{kn}{2}} = \frac{2}{n}$$

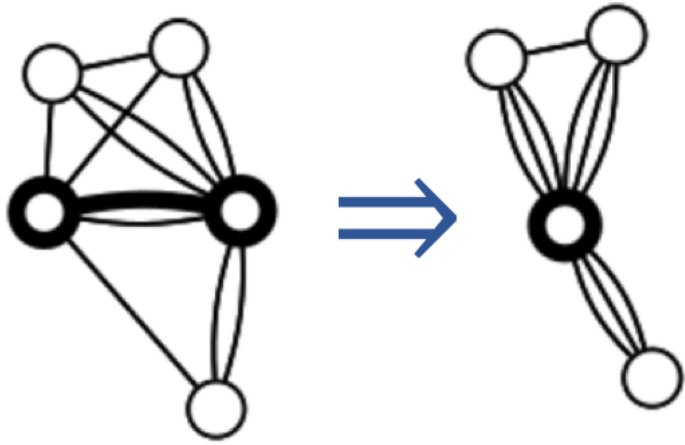
$$\mathbb{P} \left[ \begin{array}{l} \text{don't contract} \\ \text{an edge of } S^* \\ \text{at any step} \end{array} \right]$$

$$\geq \left(1 - \frac{2}{n}\right) \left(1 - \frac{2}{n-1}\right) \left(1 - \frac{2}{n-2}\right)$$

$$\dots \left(1 - \frac{2}{3}\right)$$

$$= \frac{\cancel{n-2}}{n} \frac{\cancel{n-3}}{n-1} \frac{\cancel{n-4}}{\cancel{n-2}} \dots \frac{2}{4} \frac{1}{3}$$

$$= \frac{2}{n(n-1)}$$



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