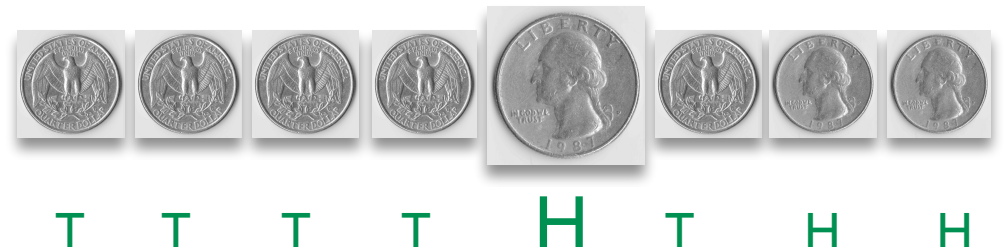


## 6. random variables



let  $X_i$  = index of

# Random Variables—Intro

A *random variable* is a numeric function of the outcome of an experiment, not the outcome itself. (Technically, neither random nor a variable, but...)

Ex.

Let  $H$  be the *number* of Heads when 20 coins are tossed

Let  $T$  be the *total* of 2 dice rolls

Let  $X$  be the *number* of coin tosses needed to see 1<sup>st</sup> head

**Note:** even if the underlying experiment has “equally likely outcomes,” an associated random variable *may not*

Outcome	$X = \#H$	$P(X)$
TT	0	$P(X=0) = 1/4$
TH	1	} $P(X=1) = 1/2$
HT	1	
HH	2	$P(X=2) = 1/4$

20 balls numbered 1, 2, ..., 20

Draw 3 without replacement

Let  $X$  = the maximum of the numbers on those 3 balls

What is  $P(X \geq 17)$

$$P(X = 20) = \binom{19}{2} / \binom{20}{3} = \frac{3}{20} = 0.150$$

$$P(X = 19) = \binom{18}{2} / \binom{20}{3} = \frac{18 \cdot 17 / 2!}{20 \cdot 19 \cdot 18 / 3!} \approx 0.134$$

$$\vdots$$

$$\sum_{i=17}^{20} P(X = i) \approx 0.508$$

Alternatively:

$$P(X \geq 17) = 1 - P(X < 17) = 1 - \binom{16}{3} / \binom{20}{3} \approx 0.508$$



Flip a (biased) coin repeatedly until 1<sup>st</sup> head observed

How many flips? Let  $X$  be that number.

$$P(X=1) = P(H) = p$$

$$P(X=2) = P(TH) = (1-p)p$$

$$P(X=3) = P(TTH) = (1-p)^2p$$

...

$$\sum_{i \geq 0} x^i = \frac{1}{1-x},$$

when  $|x| < 1$

memorize me!

Check that it is a valid probability distribution:

$$1) \quad \forall i \geq 1, P(\{X = i\}) \geq 0$$

$$2) \quad P\left(\bigcup_{i \geq 1} \{X = i\}\right) = \sum_{i \geq 1} (1-p)^{i-1}p = p \sum_{i \geq 0} (1-p)^i = p \frac{1}{1-(1-p)} = 1$$

A *discrete* random variable is one taking on a *countable* number of possible values.

Ex:

$X$  = sum of 3 dice,  $3 \leq X \leq 18$ ,  $X \in \mathbb{N}$

$Y$  = number of 1<sup>st</sup> head in seq of coin flips,  $1 \leq Y$ ,  $Y \in \mathbb{N}$

$Z$  = largest prime factor of  $(1+Y)$ ,  $Z \in \{2, 3, 5, 7, 11, \dots\}$

**Definition:** If  $X$  is a discrete random variable taking on values from a countable set  $T \subseteq \mathcal{R}$ , then

$$p(a) = \begin{cases} P(X = a) & \text{for } a \in T \\ 0 & \text{otherwise} \end{cases}$$

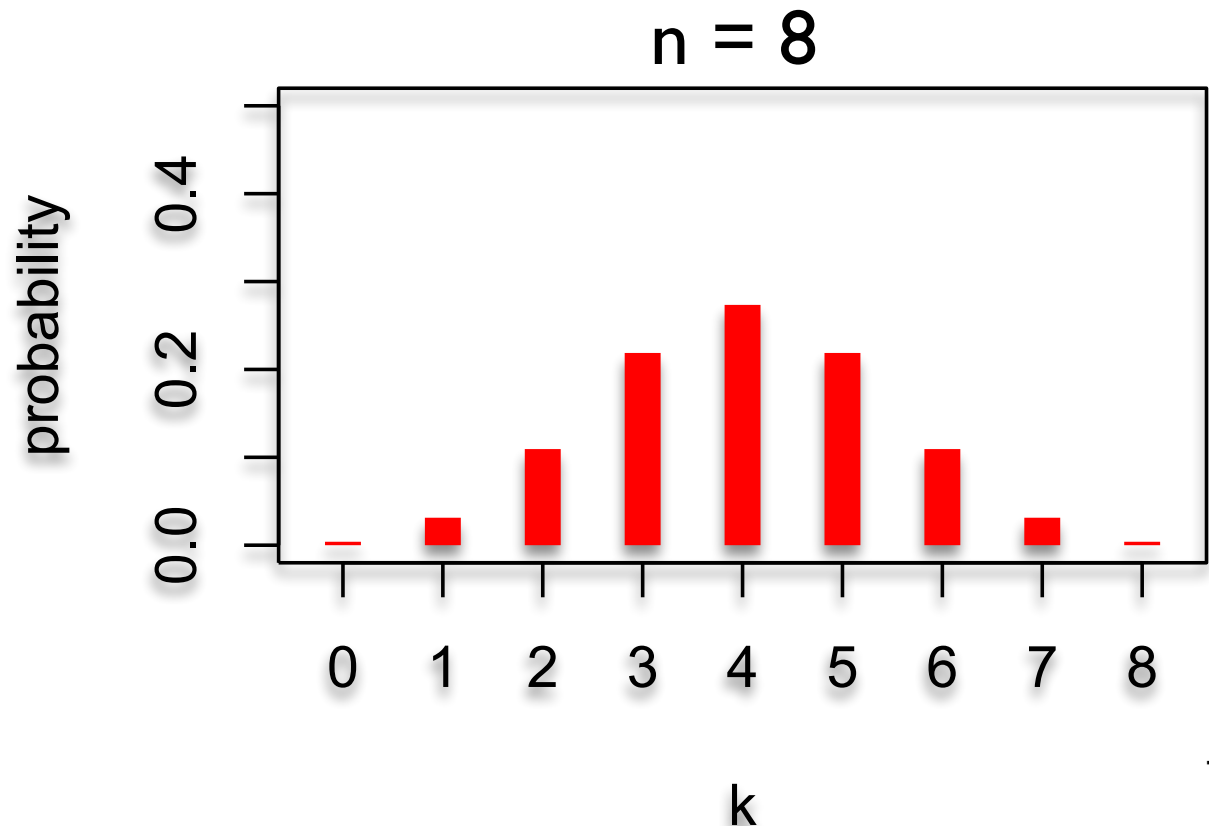
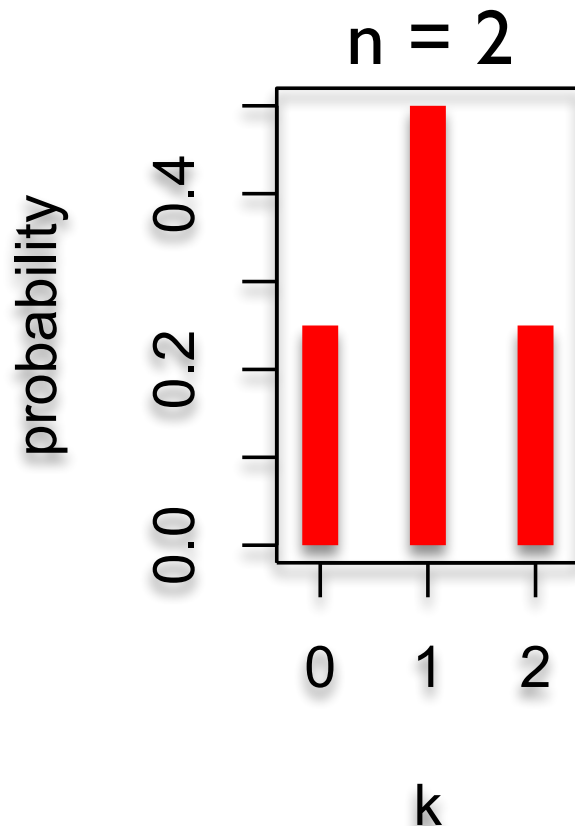
is called the *probability mass function*. Note:  $\sum_{a \in T} p(a) = 1$

---

Let  $X$  be the number of heads observed in  $n$  coin flips

$$P(X = k) = \binom{n}{k} p^k (1 - p)^{n-k}, \text{ where } p = P(H)$$

Probability mass function ( $p = 1/2$ ):



## cumulative distribution function

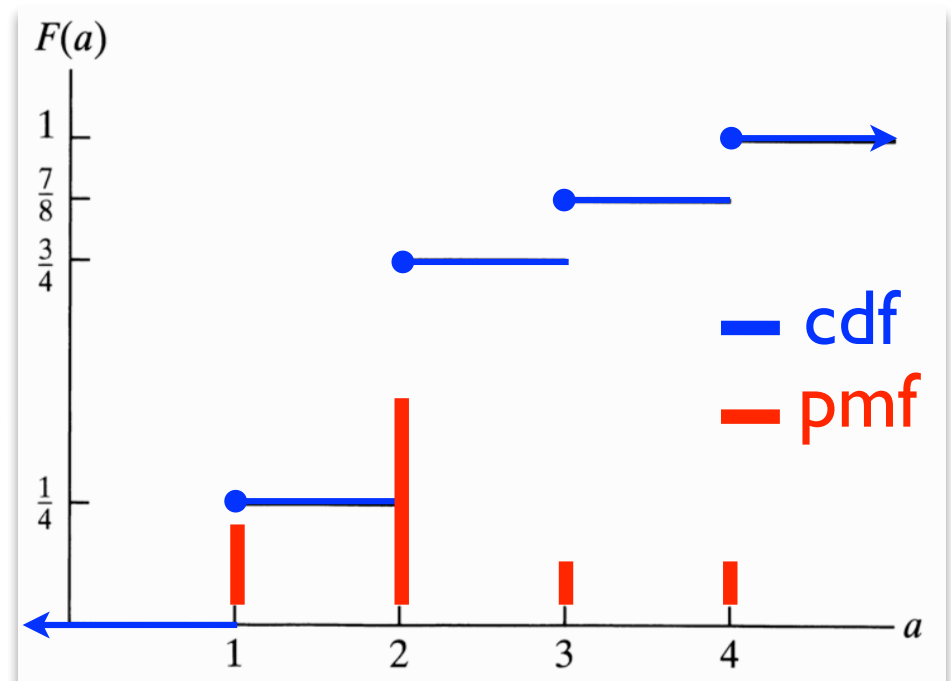
The *cumulative distribution function* for a random variable  $X$  is the function  $F: \mathcal{R} \rightarrow [0,1]$  defined by

$$F(a) = P[X \leq a]$$

Ex: if  $X$  has **probability mass function** given by:

$$p(1) = \frac{1}{4} \quad p(2) = \frac{1}{2} \quad p(3) = \frac{1}{8} \quad p(4) = \frac{1}{8}$$

$$F(a) = \begin{cases} 0 & a < 1 \\ \frac{1}{4} & 1 \leq a < 2 \\ \frac{3}{4} & 2 \leq a < 3 \\ \frac{7}{8} & 3 \leq a < 4 \\ 1 & 4 \leq a \end{cases}$$



NB: for discrete random variables, be careful about “ $\leq$ ” vs “ $<$ ”

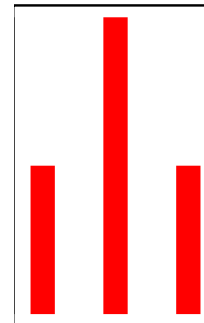
## Why use random variables?

### A. Often we just care about numbers

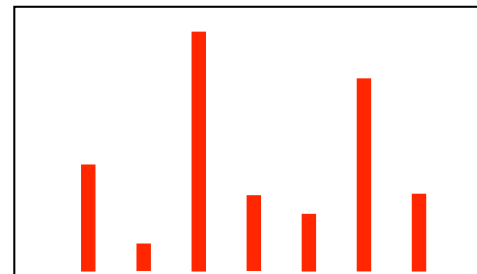
If I win \$1 per head when 20 coins are tossed, what is my average winnings? What is the most likely number? What is the probability that I win < \$5? ...

### B. It cleanly abstracts away unnecessary detail about the experiment/sample space; PMF is all we need.

Outcome	$x=\#H$	$P(X)$
TT	0	$P(X=0) = 1/4$
TH	1	$P(X=1) = 1/2$
HT	1	
HH	2	$P(X=2) = 1/4$



Flip 7 coins, roll 2 dice, and throw a dart; if dart landed in sector = dice roll mod #heads, then  $X = \dots$



expectation

For a discrete r.v.  $X$  with p.m.f.  $p(\bullet)$ , the *expectation of  $X$* , aka *expected value* or *mean*, is

$$E[X] = \sum_x xp(x)$$

average of random values, weighted  
by their respective probabilities

For the equally-likely outcomes case, this is just the average of the possible random values of  $X$

For *unequally-likely* outcomes, it is again the average of the possible random values of  $X$ , **weighted by their respective probabilities**

Ex 1: Let  $X$  = value seen rolling a fair die  $p(1), p(2), \dots, p(6) = 1/6$

$$E[X] = \sum_{i=1}^6 ip(i) = \frac{1}{6}(1 + 2 + \dots + 6) = \frac{21}{6} = 3.5$$

Ex 2: Coin flip;  $X = +1$  if H (win \$1),  $-1$  if T (lose \$1)

$$E[X] = (+1) \cdot p(+1) + (-1) \cdot p(-1) = 1 \cdot (1/2) + (-1) \cdot (1/2) = 0$$

For a discrete r.v.  $X$  with p.m.f.  $p(\bullet)$ , the *expectation of  $X$* , aka *expected value* or *mean*, is

$$E[X] = \sum_x xp(x)$$

average of random values, weighted  
by their respective probabilities

**Another view:** A 2-person gambling game. If  $X$  is how much you win playing the game once, how much would you expect to win, on average, per game, when repeatedly playing?

Ex 1: Let  $X$  = value seen rolling a fair die  $p(1), p(2), \dots, p(6) = 1/6$   
If you win  $X$  dollars for that roll, how much do you expect to win?

$$E[X] = \sum_{i=1}^6 ip(i) = \frac{1}{6}(1 + 2 + \dots + 6) = \frac{21}{6} = 3.5$$

Ex 2: Coin flip;  $X = +1$  if H (win \$1),  $-1$  if T (lose \$1)

$$E[X] = (+1) \cdot p(+1) + (-1) \cdot p(-1) = 1 \cdot (1/2) + (-1) \cdot (1/2) = 0$$

“a fair game”: in repeated play you expect to win as much as you lose. Long term net gain/loss = 0.



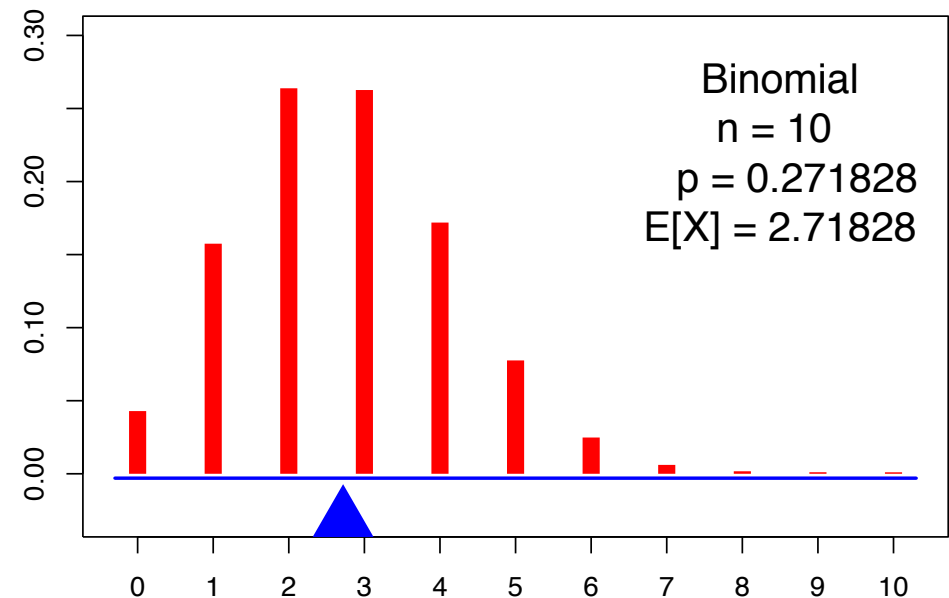
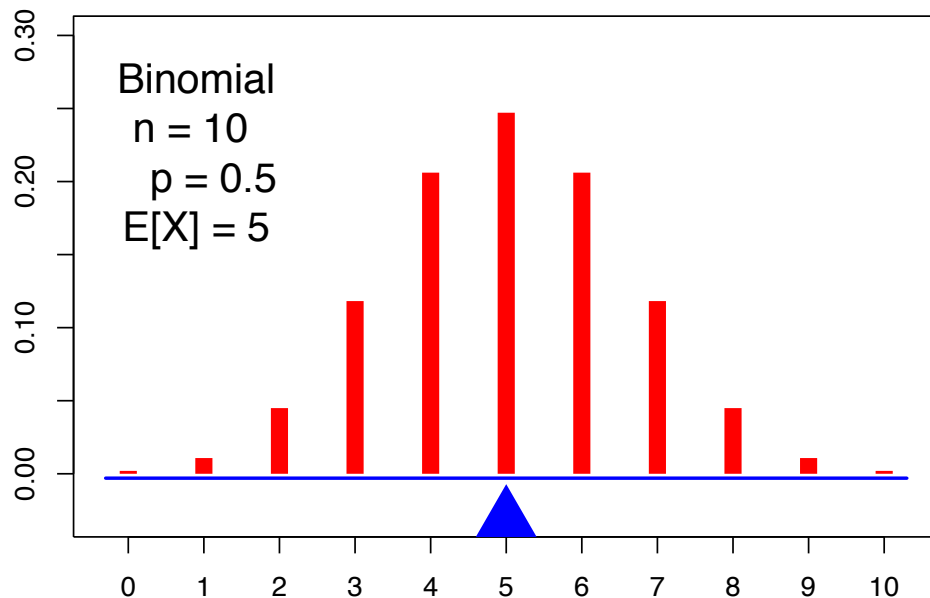
For a discrete r.v.  $X$  with p.m.f.  $p(\bullet)$ , the *expectation of  $X$* , aka *expected value* or *mean*, is

$$E[X] = \sum_x xp(x)$$

average of random values, weighted  
by their respective probabilities

**A third view:**  $E[X]$  is the “balance point” or “center of mass” of the probability mass function

Ex: Let  $X$  = number of heads seen when flipping 10 coins



Let  $X$  be the number of flips up to & including 1<sup>st</sup> head observed in repeated flips of a biased coin. If I pay you \$1 per flip, how much money would you expect to make?

$$P(H) = p; \quad P(T) = 1 - p = q$$

$$p(i) = pq^{i-1} \quad \leftarrow \text{PMF}$$

$$E[X] = \sum_{i \geq 1} ip(i) = \sum_{i \geq 1} ipq^{i-1} = p \sum_{i \geq 1} iq^{i-1} \quad (*)$$

A calculus trick:

$$\sum_{i \geq 1} iy^{i-1} = \sum_{i \geq 1} \frac{d}{dy} y^i = \sum_{i \geq 0} \frac{d}{dy} y^i = \frac{d}{dy} \sum_{i \geq 0} y^i = \frac{d}{dy} \frac{1}{1-y} = \frac{1}{(1-y)^2}$$

So (\*) becomes:

$$E[X] = p \sum_{i \geq 1} iq^{i-1} = \frac{p}{(1-q)^2} = \frac{p}{p^2} = \frac{1}{p}$$

E.g.:

$p=1/2$ ; on average head every 2<sup>nd</sup> flip

$p=1/10$ ; on average, head every 10<sup>th</sup> flip.

How much  
would you  
pay to play?

Let  $X$  be the number of heads observed in  $n$  repeated flips of a biased coin. If I pay you \$1 per head, how much money would you expect to make?

E.g.:

$p=1/2$ ; on average,  
n/2 heads

$p=1/10$ ; on average,  
n/10 heads

How much would  
you pay to play?

$$\begin{aligned} E[X] &= \sum_{i=0}^n i \binom{n}{i} p^i (1-p)^{n-i} \\ &= \sum_{i=1}^n i \binom{n}{i} p^i (1-p)^{n-i} \\ &= \sum_{i=1}^n n \binom{n-1}{i-1} p^i (1-p)^{n-i} \\ &= np \sum_{i=1}^n \binom{n-1}{i-1} p^{i-1} (1-p)^{n-i} \\ &= np \sum_{j=0}^{n-1} \binom{n-1}{j} p^j (1-p)^{n-1-j} \\ &= np(p + (1-p))^{n-1} = np \end{aligned}$$

# expectation of a *function* of a random variable

Calculating  $E[g(X)]$ :

$Y=g(X)$  is a new r.v. Calculate  $P[Y=j]$ , then apply defn:

$X$  = sum of 2 dice rolls

$i$	$p(i) = P[X=i]$	$i \cdot p(i)$
2	1/36	2/36
3	2/36	6/36
4	3/36	12/36
5	4/36	20/36
6	5/36	30/36
7	6/36	42/36
8	5/36	40/36
9	4/36	36/36
10	3/36	30/36
11	2/36	22/36
12	1/36	12/36

$$E[X] = \sum_i i p(i) = \frac{252}{36} = 7$$

$Y = g(X) = X \bmod 5$

$j$	$q(j) = P[Y = j]$	$j \cdot q(j)$
0	$4/36 + 3/36 = 7/36$	0/36
1	$5/36 + 2/36 = 7/36$	7/36
2	$1/36 + 6/36 + 1/36 = 8/36$	16/36
3	$2/36 + 5/36 = 7/36$	21/36
4	$3/36 + 4/36 = 7/36$	28/36

$$E[Y] = \sum_j j q(j) = \frac{72}{36} = 2$$

Way 2

## expectation of a *function* of a random variable

Calculating  $E[g(X)]$ : Another way – *add in a different order*,  
using  $P[X=...]$  instead of calculating  $P[Y=...]$

$X = \text{sum of 2 dice rolls}$

$i$	$p(i) = P[X=i]$	$g(i) \cdot p(i)$
2	1/36	2/36
3	2/36	6/36
4	3/36	12/36
5	4/36	0/36
6	5/36	5/36
7	6/36	12/36
8	5/36	15/36
9	4/36	16/36
10	3/36	0/36
11	2/36	2/36
12	1/36	2/36

$Y = g(X) = X \bmod 5$

$j$	$q(j) = P[Y = j]$	$j \cdot q(j)$
0	$4/36 + 3/36 = 7/36$	0/36
1	$5/36 + 2/36 = 7/36$	7/36
2	$1/36 + 6/36 + 1/36 = 8/36$	16/36
3	$2/36 + 5/36 = 7/36$	21/36
4	$3/36 + 4/36 = 7/36$	28/36

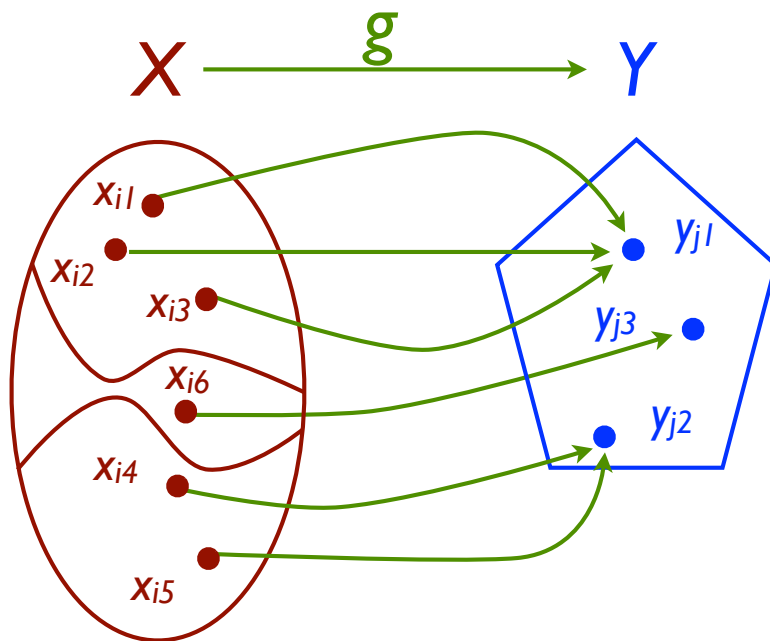
$$E[Y] = \sum_j j q(j) = \frac{72}{36} = 2$$

$$E[g(X)] = \sum_i g(i) p(i) = \frac{72}{36} = 2$$

Above example is not a fluke!

**Theorem:** if  $Y = g(X)$ , then  $E[Y] = \sum_i g(x_i)p(x_i)$ , where  $x_i, i = 1, 2, \dots$  are all possible values of  $X$ .

**Proof:** Let  $y_j, j = 1, 2, \dots$  be all possible values of  $Y$ .



Note that  $S_j = \{x_i \mid g(x_i) = y_j\}$  is a *partition* of the domain of  $g$ .

$$\begin{aligned} \sum_i g(x_i)p(x_i) &= \sum_j \sum_{i:g(x_i)=y_j} g(x_i)p(x_i) \\ &= \sum_j \sum_{i:g(x_i)=y_j} y_j p(x_i) \\ &= \sum_j y_j \sum_{i:g(x_i)=y_j} p(x_i) \\ &= \sum_j y_j P\{g(X) = y_j\} \\ &= E[g(X)] \end{aligned}$$

Above  $E[X \bmod 5] = (E[X]) \bmod 5$

Is that a Law or a Coincidence?

Try  $X \bmod 2$ ,  $X \bmod 3$ ,  $X \bmod 4$ , ...

## properties of expectation

A & B each bet \$1, then flip 2 coins:

HH	A wins \$2
HT	Each takes back \$1
TH	
TT	B wins \$2

Let  $X$  be A's net gain: +1, 0, -1, resp.:

$$\begin{aligned}P(X = +1) &= 1/4 \\P(X = 0) &= 1/2 \\P(X = -1) &= 1/4\end{aligned}$$

What is  $E[X]$ ?

$$E[X] = 1 \cdot 1/4 + 0 \cdot 1/2 + (-1) \cdot 1/4 = 0$$

What is  $E[X^2]$ ?

$$E[X^2] = 1^2 \cdot 1/4 + 0^2 \cdot 1/2 + (-1)^2 \cdot 1/4 = 1/2$$

Big Deal Note:  
 $E[X^2] \neq E[X]^2$



### Linearity of expectation, I

For any constants  $a, b$ :  $E[aX + b] = aE[X] + b$

Proof:

$$\begin{aligned} E[aX + b] &= \sum_x (ax + b) \cdot p(x) \\ &= a \sum_x xp(x) + b \sum_x p(x) \\ &= aE[X] + b \end{aligned}$$

## properties of expectation—example

A & B each bet \$1, then flip 2 coins:

HH	A wins \$2
HT	Each takes back \$1
TH	
TT	B wins \$2

Let  $X$  = A's net gain: +1, 0, -1, resp.:

$$P(X = +1) = 1/4$$

$$P(X = 0) = 1/2$$

$$P(X = -1) = 1/4$$

What is  $E[X]$ ?

$$E[X] = 1 \cdot 1/4 + 0 \cdot 1/2 + (-1) \cdot 1/4 = 0$$

What is  $E[X^2]$ ?

$$E[X^2] = 1^2 \cdot 1/4 + 0^2 \cdot 1/2 + (-1)^2 \cdot 1/4 = 1/2$$

What is  $E[2X+1]$ ?

$$E[2X + 1] = 2E[X] + 1 = 2 \cdot 0 + 1 = 1$$

From slide 20

Example:

Caezzo's Palace Casino offers the following game: They flip a biased coin ( $P(\text{Heads}) = 0.10$ ) until the first Head comes up. "You're on a hot streak now! The more Tails the more you win!" Let  $X$  be the number of flips up to & including 1<sup>st</sup> head. They will pay you \$2 per flip, i.e.,  $2X$  dollars. They charge you \$25 to play.

Q: Is it a fair game? On average, how much would you expect to win/lose per game, if you play it repeatedly?

A: Not fair. Your net winnings per game is  $2X - 25$ , and  
 $E[2X - 25] = 2 E[X] - 25 = 2(1/0.10) - 25 = -5$ ,  
i.e., you lose \$5 per game on average

## Linearity, II

Let  $X$  and  $Y$  be two random variables derived from outcomes of a *single* experiment. Then

$$E[X+Y] = E[X] + E[Y]$$

True even if  $X, Y$  dependent

**Proof:** Assume the sample space  $S$  is countable. (The result is also true for uncountable  $S$ .) Let  $X(s)$ ,  $Y(s)$  be the values of these r.v.'s for outcome  $s \in S$ .

Claim:  $E[X] = \sum_{s \in S} X(s) \cdot p(s)$

Proof: similar to that for “expectation of a function of an r.v.,” i.e., the events “ $X=x$ ” partition  $S$ , so sum above can be rearranged to match the definition of  $E[X] = \sum_x x \cdot P(X = x)$

Then:

$$\begin{aligned} E[X+Y] &= \sum_{s \in S} (X[s] + Y[s]) p(s) \\ &= \sum_{s \in S} X[s] p(s) + \sum_{s \in S} Y[s] p(s) = E[X] + E[Y] \end{aligned}$$

# properties of expectation-example

A & B each bet \$1, then flip 2 coins:

HH	A wins \$2
HT	Each takes back \$1
TH	
TT	B wins \$2

Let  $X$  = A's net gain: +1, 0, -1, resp.:

$$P(X = +1) = 1/4$$

$$P(X = 0) = 1/2$$

$$P(X = -1) = 1/4$$

What is  $E[X]$ ?

$$E[X] = 1 \cdot 1/4 + 0 \cdot 1/2 + (-1) \cdot 1/4 = 0$$

What is  $E[X^2]$ ?

$$E[X^2] = 1^2 \cdot 1/4 + 0^2 \cdot 1/2 + (-1)^2 \cdot 1/4 = 1/2$$

What is  $E[X^2 + 2X + 1]$ ?

$$E[X^2 + 2X + 1] = E[X^2] + 2E[X] + 1 = 1/2 + 2 \cdot 0 + 1 = 1.5$$

Intuitively, not independent

From slide 20

### Example

$X = \#$  of heads in *one* coin flip, where  $P(X=1) = p$ .

What is  $E(X)$ ?

$$E[X] = 1 \cdot p + 0 \cdot (1-p) = p$$

👉 defn of  $E[ ]$

Let  $X_i, 1 \leq i \leq n$ , be  $\#$  of H in flip of coin with  $P(X_i=1) = p_i$

What is the expected number of heads when all are flipped?

$$E[\sum_i X_i] = \sum_i E[X_i] = \sum_i p_i$$

Special case:  $p_1 = p_2 = \dots = p$  :

$$E[\# \text{ of heads in } n \text{ flips}] = pn$$

👉 Compare to slide 15

## Note:

Linearity is special!

It is *not* true in general that

$$E[X \cdot Y] = E[X] \cdot E[Y]$$

$$E[X^2] = E[X]^2$$

$$E[X/Y] = E[X] / E[Y]$$

$$E[\sinh(X)] = \sinh(E[X])$$

•

•

•

← counterexample above

variance



Alice & Bob are gambling (again).  $X$  = Alice's gain per flip:

$$X = \begin{cases} +1 & \text{if Heads} \\ -1 & \text{if Tails} \end{cases}$$

$$E[X] = 0$$

... Time passes ...

Alice (yawning) says “let's raise the stakes”

$$Y = \begin{cases} +1000 & \text{if Heads} \\ -1000 & \text{if Tails} \end{cases}$$

$$E[Y] = 0, \text{ as before.}$$

Are you (Bob) equally happy to play the new game?

$E[X]$  measures the “average” or “central tendency” of  $X$ .

What about its *variability*? E.g., is  $X$  usually near average, or far above/below it?

If  $E[X] = \mu$ , then  $E[|X - \mu|]$  seems like a natural quantity to look at: how much do we expect (on average)  $X$  to deviate from its average.

Unfortunately, it's a bit inconvenient mathematically; following is nicer/easier/much more common.

## Definitions

The *variance* of a random variable  $X$  with mean  $E[X] = \mu$  is

$$\text{Var}[X] = E[(X-\mu)^2],$$

often denoted  $\sigma^2$ .

The *standard deviation* of  $X$  is

$$\sigma = \sqrt{\text{Var}[X]}$$

## what does variance tell us?

The *variance* of a random variable  $X$  with mean  $E[X] = \mu$  is  $\text{Var}[X] = E[(X-\mu)^2]$ , often denoted  $\sigma^2$ .

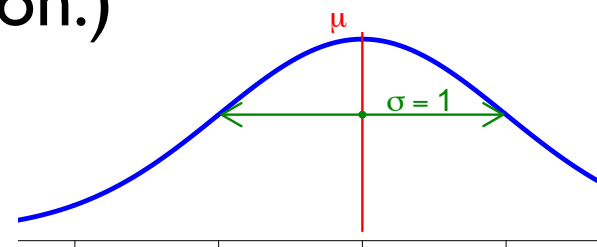
I: Square always  $\geq 0$ , and exaggerated as  $X$  moves away from  $\mu$ , so  $\text{Var}[X]$  emphasizes *deviation* from the mean.

II: Numbers vary a lot depending on exact distribution of  $X$ , but it is common that  $X$  is

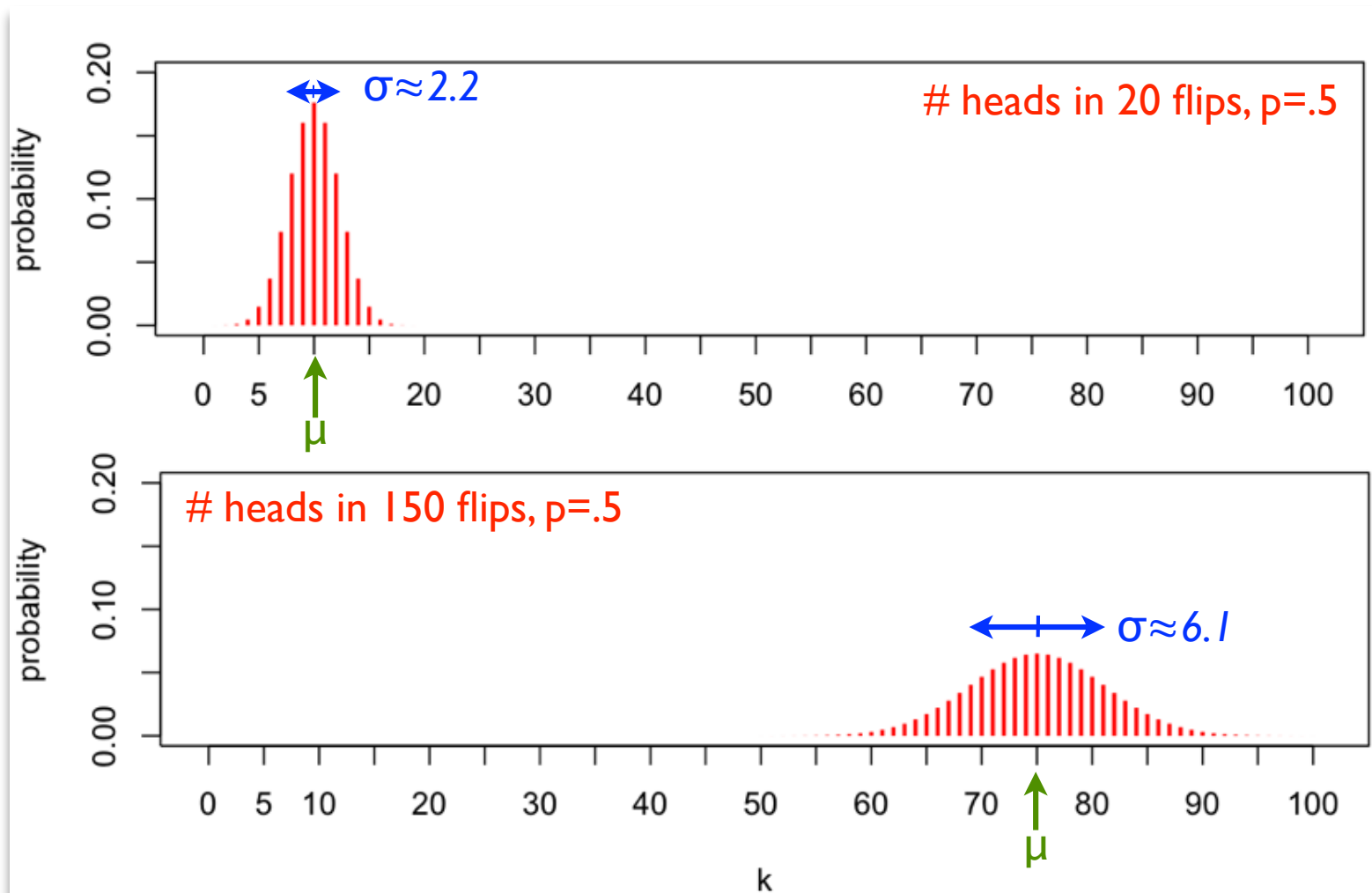
within  $\mu \pm \sigma$  ~66% of the time, and

within  $\mu \pm 2\sigma$  ~95% of the time.

(We'll see the reasons for this soon.)



$\mu = E[X]$  is about *location*;  $\sigma = \sqrt{\text{Var}(X)}$  is about *spread*



Blue arrows denote the interval  $\mu \pm \sigma$

(and note  $\sigma$  bigger in absolute terms in second ex., but smaller as a proportion of  $\mu$  or max.) 33

Alice & Bob are gambling (again).  $X$  = Alice's gain per flip:

$$X = \begin{cases} +1 & \text{if Heads} \\ -1 & \text{if Tails} \end{cases}$$

$$E[X] = 0$$

$$\underline{\text{Var}[X] = 1}$$

... Time passes ...

Alice (yawning) says “let's raise the stakes”

$$Y = \begin{cases} +1000 & \text{if Heads} \\ -1000 & \text{if Tails} \end{cases}$$

$$E[Y] = 0, \text{ as before.}$$

$$\underline{\text{Var}[Y] = 1,000,000}$$

Are you (Bob) equally happy to play the new game?

Two games:

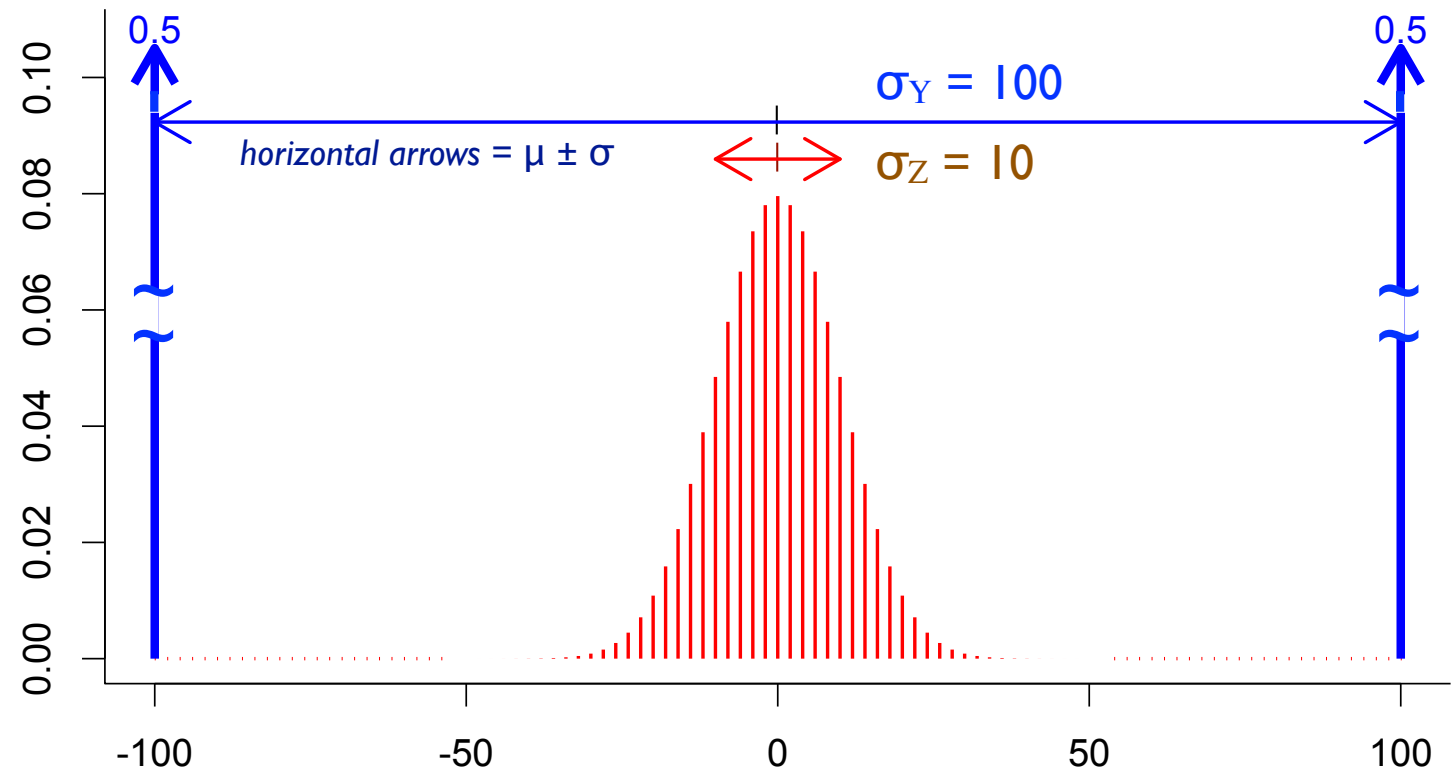
a) flip 1 coin, win  $Y = \$100$  if heads,  $\$-100$  if tails

b) flip 100 coins, win  $Z = (\text{\#(heads)} - \text{\#(tails)})$  dollars

Same expectation in both:  $E[Y] = E[Z] = 0$

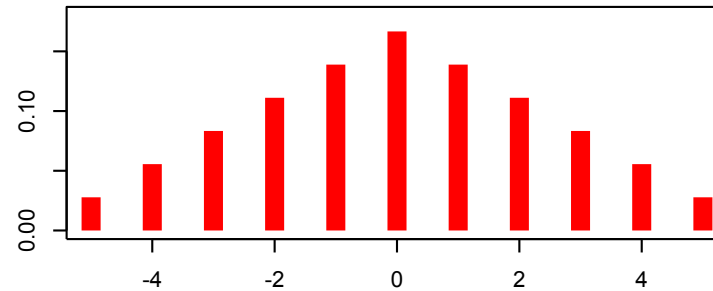
Same extremes in both: max gain =  $\$100$ ; max loss =  $\$100$

But  
variability  
is very  
different:



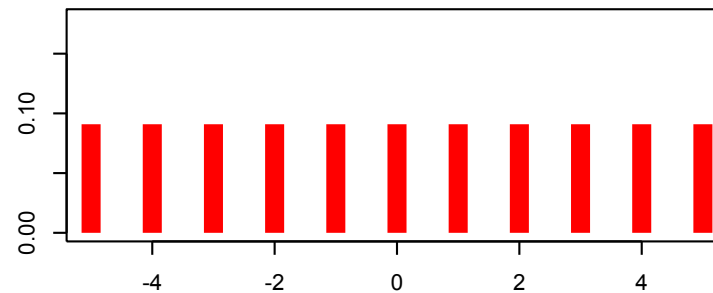
## more variance examples

$X_1$  = sum of 2 fair dice, minus 7



$$\sigma^2 = 5.83$$

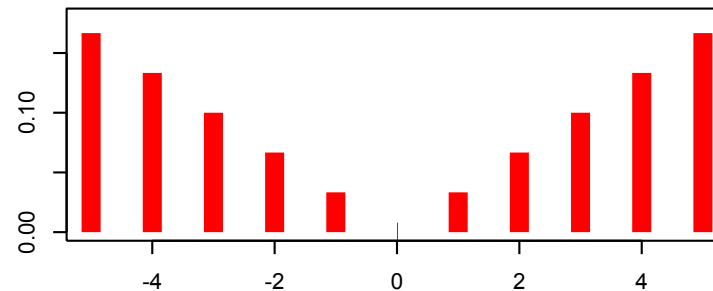
$X_2$  = fair 11-sided die labeled -5, ..., 5



$$\sigma^2 = 10$$

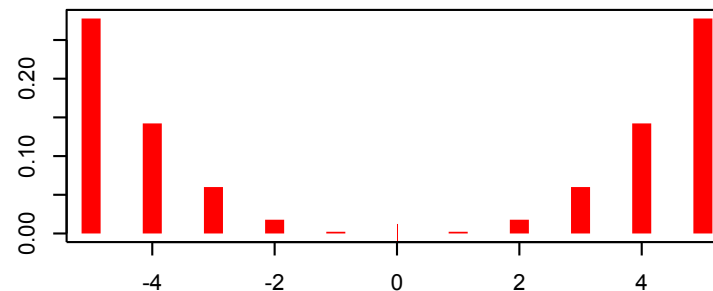
-1, 0, +1

$X_3$  =  $Y - 6 \cdot \text{signum}(Y)$ , where  $Y$  is the difference of 2 fair dice, given no doubles



$$\sigma^2 = 15$$

$X_4$  =  $X_3$  when 3 pairs of dice all give same  $X_3$



$$\sigma^2 = 19.7$$

NB: Wow,  
kinda complex;  
see [slide 9](#)



$$\text{Var}(X) = E[X^2] - (E[X])^2$$

$$\begin{aligned}\text{Var}(X) &= E[(X - \mu)^2] \\ &= E[X^2 - 2\mu X + \mu^2] \\ &= E[X^2] - 2\mu E[X] + \mu^2 \\ &= E[X^2] - 2\mu^2 + \mu^2 \\ &= E[X^2] - \mu^2 \\ &= E[X^2] - (E[X])^2\end{aligned}$$

### Example:

What is  $\text{Var}[X]$  when  $X$  is outcome of one fair die?

$$\begin{aligned} E[X^2] &= 1^2 \left(\frac{1}{6}\right) + 2^2 \left(\frac{1}{6}\right) + 3^2 \left(\frac{1}{6}\right) + 4^2 \left(\frac{1}{6}\right) + 5^2 \left(\frac{1}{6}\right) + 6^2 \left(\frac{1}{6}\right) \\ &= \left(\frac{1}{6}\right) (91) \end{aligned}$$

$E[X] = 7/2$ , so

$$\text{Var}(X) = \frac{91}{6} - \left(\frac{7}{2}\right)^2 = \frac{35}{12}$$

## properties of variance

$$\text{Var}[aX+b] = a^2 \text{Var}[X]$$

NOT linear;  
insensitive to location (b),  
quadratic in scale (a)

$$\begin{aligned}\text{Var}(aX + b) &= E[(aX + b - a\mu - b)^2] \\ &= E[a^2(X - \mu)^2] \\ &= a^2 E[(X - \mu)^2] \\ &= a^2 \text{Var}(X)\end{aligned}$$

Ex:

$$X = \begin{cases} +1 & \text{if Heads} \\ -1 & \text{if Tails} \end{cases} \quad \begin{aligned} E[X] &= 0 \\ \text{Var}[X] &= 1 \end{aligned}$$

$$Y = \begin{cases} +1000 & \text{if Heads} \\ -1000 & \text{if Tails} \end{cases} \quad \begin{aligned} Y &= 1000 X \\ E[Y] &= E[1000 X] = 1000 E[X] = 0 \\ \text{Var}[Y] &= \text{Var}[10^3 X] = 10^6 \text{Var}[X] = 10^6 \end{aligned}$$

In general:

$$\text{Var}[X+Y] \neq \text{Var}[X] + \text{Var}[Y]$$

^^^^^^

NOT linear

Ex 1:

Let  $X = \pm 1$  based on 1 coin flip

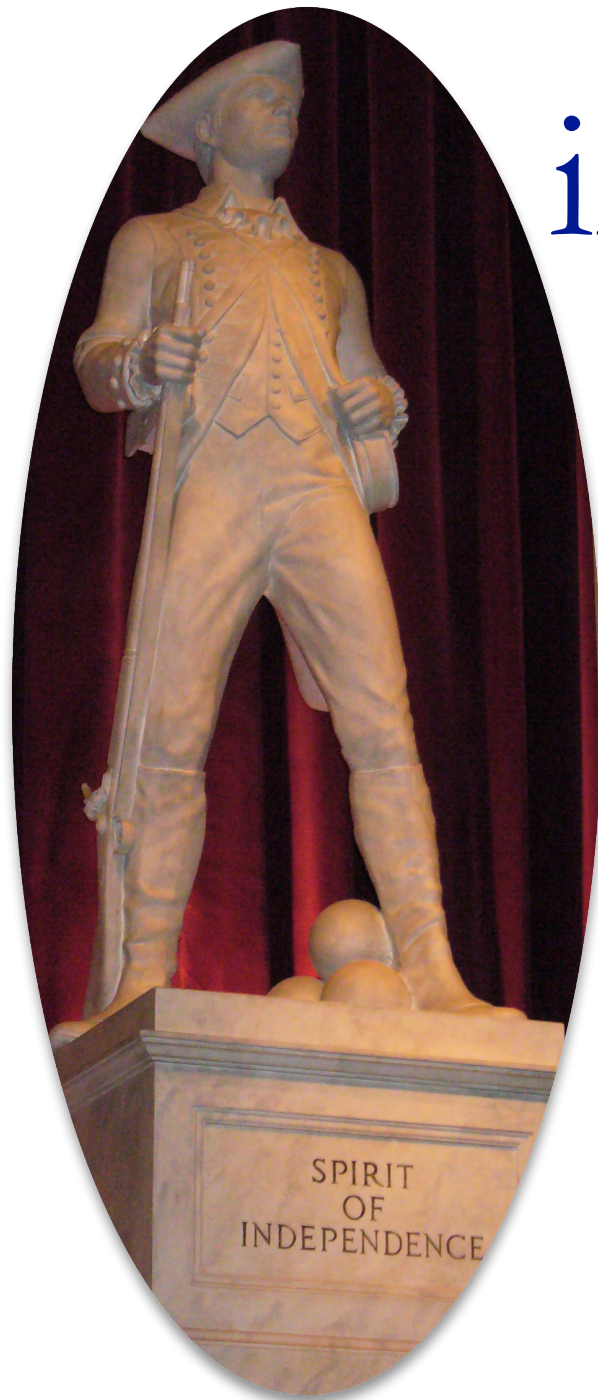
As shown above,  $E[X] = 0, \text{Var}[X] = 1$

Let  $Y = -X$ ; then  $\text{Var}[Y] = (-1)^2 \text{Var}[X] = 1$

But  $X+Y = 0$ , always, so  $\text{Var}[X+Y] = 0$

Ex 2:

As another example, is  $\text{Var}[X+X] = 2\text{Var}[X]$ ?



independence

and

joint

distributions



**Defn:** Random variable  $X$  and event  $E$  are independent if event  $E$  is independent of event  $\{X=x\}$  (for all fixed  $x$ ), i.e.

$$\forall x \ P(\{X = x\} \ \& \ E) = P(\{X=x\}) \cdot P(E)$$

**Defn:** Two random variables  $X$  and  $Y$  are independent if the events  $\{X=x\}$  and  $\{Y=y\}$  are independent (for all fixed  $x, y$ ), i.e.

$$\forall x, y \ P(\{X = x\} \ \& \ \{Y=y\}) = P(\{X=x\}) \cdot P(\{Y=y\})$$

Intuition as before: knowing  $X$  doesn't help you guess  $Y$  or  $E$  and vice versa.

Random variable  $X$  and event  $E$  are independent if

$$\forall x \ P(\{X = x\} \ \& \ E) = P(\{X=x\}) \cdot P(E)$$

Ex 1: Roll a fair die to obtain a random number  $1 \leq X \leq 6$ , then flip a fair coin  $X$  times. Let  $E$  be the event that the number of heads is even.

$$P(\{X=x\}) = 1/6 \text{ for any } 1 \leq x \leq 6,$$

$$P(E) = 1/2$$

$$P(\{X=x\} \ \& \ E) = 1/12, \text{ so they are independent}$$

Ex 2: as above, and let  $F$  be the event that the total number of heads = 6.

$P(F) = 2^{-6}/6 > 0$ , and considering, say,  $X=4$ , we have  $P(X=4) = 1/6 > 0$  (as above), but  $P(\{X=4\} \ \& \ F) = 0$ , since you can't see 6 heads in 4 flips. So  $X$  &  $F$  are *dependent*. (Knowing that  $X$  is  $<6$  renders  $F$  impossible; knowing that  $F$  happened means  $X$  must be 6.)

Two random variables  $X$  and  $Y$  are independent if the events  $\{X=x\}$  and  $\{Y=y\}$  are independent (for any  $x, y$ ), i.e.

$$\forall x, y \ P(\{X = x\} \ \& \ \{Y=y\}) = P(\{X=x\}) \cdot P(\{Y=y\})$$

Ex: Let  $X$  be number of heads in first  $n$  of  $2n$  coin flips,  $Y$  be number in the last  $n$  flips, and let  $Z$  be the total.  $X$  and  $Y$  are independent:

$$P(X = j) = \binom{n}{j} 2^{-n}$$

$$P(Y = k) = \binom{n}{k} 2^{-n}$$

$$P(X = j \wedge Y = k) = \binom{n}{j} \binom{n}{k} 2^{-2n} = P(X = j)P(Y = k)$$

But  $X$  and  $Z$  are *not* independent, since, e.g., knowing that  $X = 0$  precludes  $Z > n$ . E.g.,  $P(X = 0)$  and  $P(Z = n+1)$  are both positive, but  $P(X = 0 \ \& \ Z = n+1) = 0$ .



Independence simplifies some  $E[ ]$  and  $\text{Var}[ ]$  calculations.

(Jump to slide 60)

Often, several random variables are *simultaneously* observed

$X$  = height and  $Y$  = weight

$X$  = cholesterol and  $Y$  = blood pressure

$X_1, X_2, X_3$  = work loads on servers A, B, C

*Joint* probability mass function:

$$f_{XY}(x, y) = P(\{X = x\} \& \{Y = y\})$$

*Joint* cumulative distribution function:

$$F_{XY}(x, y) = P(\{X \leq x\} \& \{Y \leq y\})$$

## Two joint PMFs

W \ Z	1	2	3
1	2/24	2/24	2/24
2	2/24	2/24	2/24
3	2/24	2/24	2/24
4	2/24	2/24	2/24

X \ Y	1	2	3
1	4/24	1/24	1/24
2	0	3/24	3/24
3	0	4/24	2/24
4	4/24	0	2/24

$$P(W = Z) = 3 * 2/24 = 6/24$$

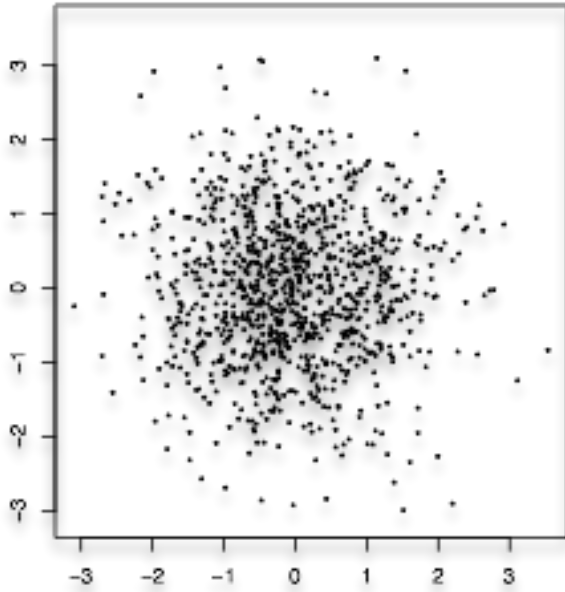
$$P(X = Y) = (4 + 3 + 2)/24 = 9/24$$

Can look at arbitrary relationships among variables this way

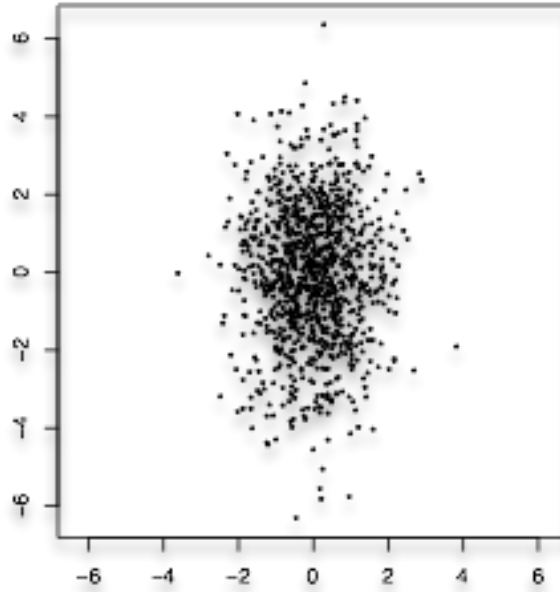
# sampling from a joint distribution

Top row: independent variables

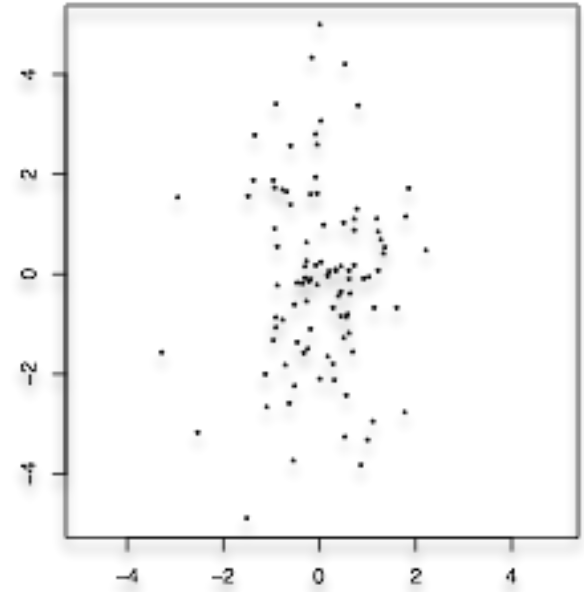
$\text{var}(x)=1, \text{var}(y)=1, \text{cov}=0, n=1000$



$\text{var}(x)=1, \text{var}(y)=3, \text{cov}=0, n=1000$

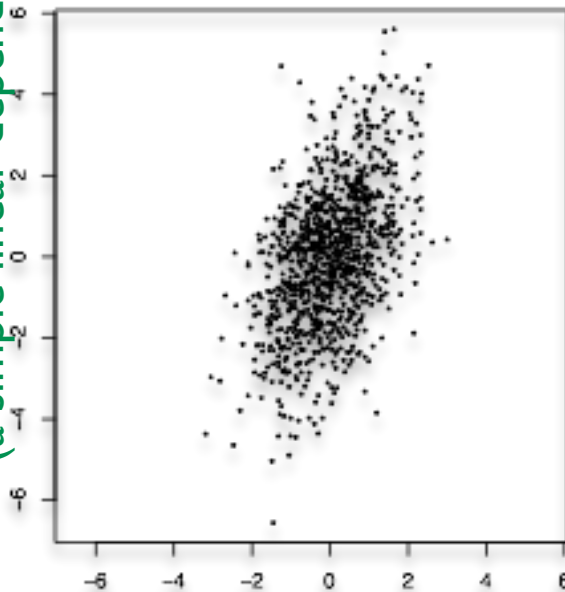


$\text{var}(x)=1, \text{var}(y)=3, \text{cov}=0, n=100$

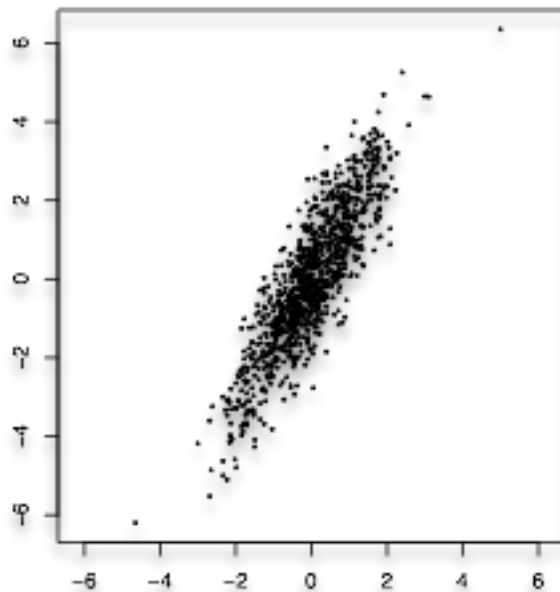


bottom row: dependent variables  
(a simple linear dependence)

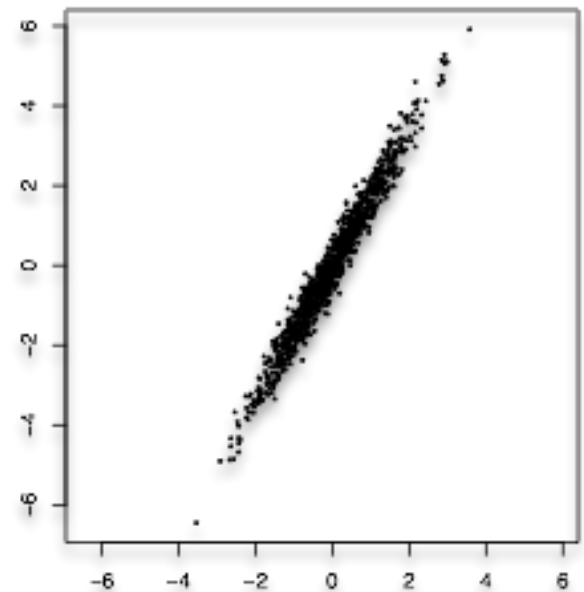
$\text{var}(x)=1, \text{var}(y)=3, \text{cov}=0.8, n=1000$



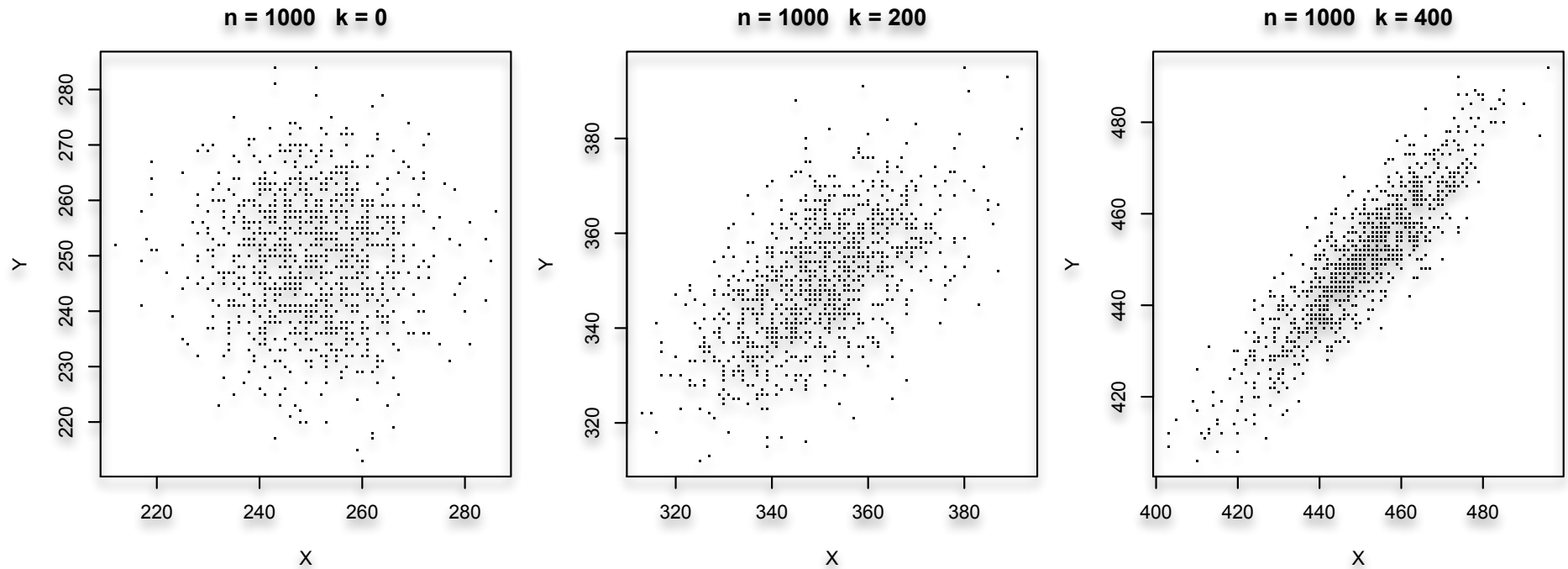
$\text{var}(x)=1, \text{var}(y)=3, \text{cov}=1.5, n=1000$



$\text{var}(x)=1, \text{var}(y)=3, \text{cov}=1.7, n=1000$



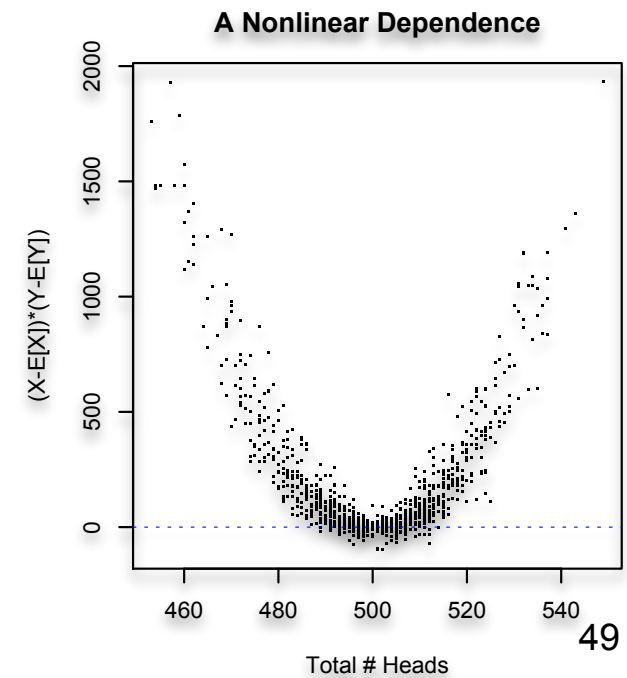
## another example



Flip  $n$  fair coins

$X = \text{\#Heads seen in first } n/2+k$

$Y = \text{\#Heads seen in last } n/2+k$



## Two joint PMFs

$W \backslash Z$	1	2	3	$f_W(w)$
1	2/24	2/24	2/24	6/24
2	2/24	2/24	2/24	6/24
3	2/24	2/24	2/24	6/24
4	2/24	2/24	2/24	6/24
$f_Z(z)$	8/24	8/24	8/24	

$X \backslash Y$	1	2	3	$f_X(x)$
1	4/24	1/24	1/24	6/24
2	0	3/24	3/24	6/24
3	0	4/24	2/24	6/24
4	4/24	0	2/24	6/24
$f_Y(y)$	8/24	8/24	8/24	

**Marginal PMF of one r.v.: sum over the other** (Law of total probability)

$$f_Y(y) = \sum_x f_{XY}(x, y)$$

$$f_X(x) = \sum_y f_{XY}(x, y)$$

**Question:** Are  $W$  &  $Z$  independent? Are  $X$  &  $Y$  independent?

**Repeating the Definition:** Two random variables  $X$  and  $Y$  are independent if the events  $\{X=x\}$  and  $\{Y=y\}$  are independent (for any fixed  $x, y$ ), i.e.

$$\forall x, y \ P(\{X = x\} \ \& \ \{Y=y\}) = P(\{X=x\}) \cdot P(\{Y=y\})$$

**Equivalent Definition:** Two random variables  $X$  and  $Y$  are independent if their *joint* probability mass function is the product of their *marginal* distributions, i.e.

$$\forall x, y \ f_{XY}(x,y) = f_X(x) \cdot f_Y(y)$$

Exercise: Show that this is also true of their *cumulative* distribution functions

## expectation of a function of 2 r.v.'s

A function  $g(X,Y)$  defines a new random variable.

Its expectation is:

$$E[g(X, Y)] = \sum_x \sum_y g(x, y) f_{XY}(x, y)$$

👉 like [slide 18](#)

Expectation is linear. E.g., if  $g$  is linear:

$$E[g(X, Y)] = E[a X + b Y + c] = a E[X] + b E[Y] + c$$

Example:

$$g(X, Y) = 2X - Y$$

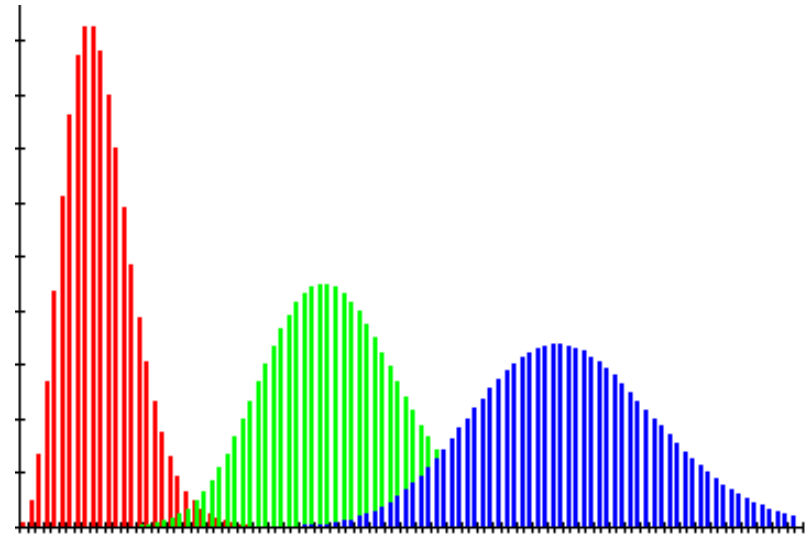
$$E[g(X, Y)] = 72/24 = 3$$

$$\begin{aligned} E[g(X, Y)] &= 2 \cdot E[X] - E[Y] \\ &= 2 \cdot \underline{2.5} - \underline{2} = 3 \end{aligned}$$

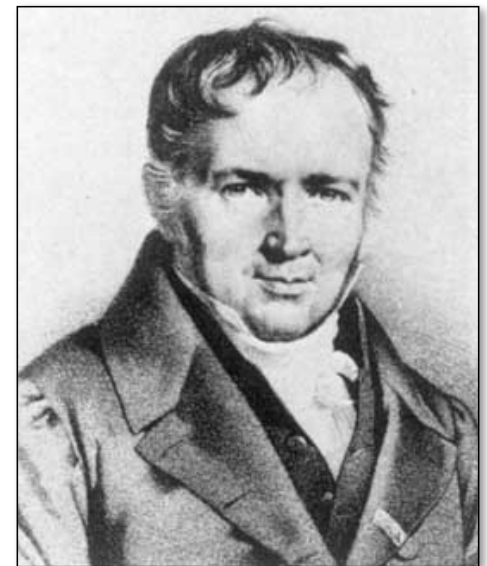
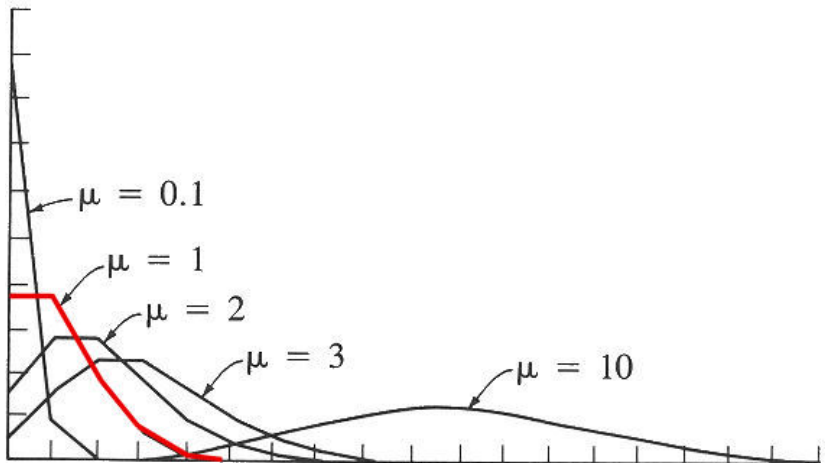
X \ Y	1	2	3
1	1 • 4/24	0 • 1/24	-1 • 1/24
2	3 • 0/24	2 • 3/24	1 • 3/24
3	5 • 0/24	4 • 4/24	3 • 2/24
4	7 • 4/24	6 • 0/24	5 • 2/24

recall both marginals are uniform





# a zoo of (discrete) random variables



## discrete uniform random variables

A discrete random variable  $X$  **equally likely** to take any (integer) value between integers  $a$  and  $b$ , inclusive, is **uniform**.

Notation:  $X \sim \text{Unif}(a,b)$

Probability:  $P(X = i) = \frac{1}{b - a + 1}$

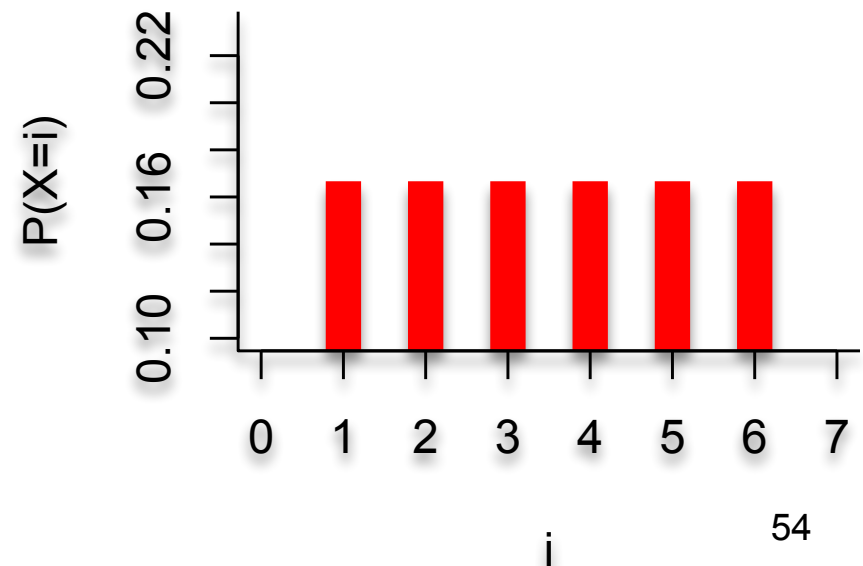
Mean, Variance:  $E[X] = \frac{a + b}{2}$ ,  $\text{Var}[X] = \frac{(b - a)(b - a + 1)}{12}$

*Example:* value shown on one roll of a fair die is  $\text{Unif}(1,6)$ :

$$P(X=i) = 1/6$$

$$E[X] = 7/2$$

$$\text{Var}[X] = 35/12$$



# Bernoulli random variables

An experiment results in “Success” or “Failure”

$X$  is a random *indicator variable* (1 = success, 0 = failure)

$$P(X=1) = p \quad \text{and} \quad P(X=0) = 1-p$$

$X$  is called a *Bernoulli* random variable:  $X \sim \text{Ber}(p)$

$$E[X] = E[X^2] = p$$

$$\text{Var}(X) = E[X^2] - (E[X])^2 = p - p^2 = p(1-p)$$

Examples:

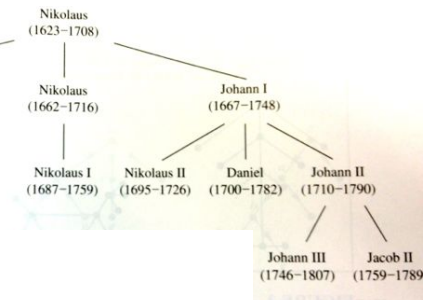
coin flip

random binary digit

whether a disk drive crashed



Jacob (aka James, Jacques)  
Bernoulli, 1654 – 1705



Consider  $n$  independent random variables  $Y_i \sim \text{Ber}(p)$

$X = \sum_i Y_i$  is the number of successes in  $n$  trials

$X$  is a *Binomial* random variable:  $X \sim \text{Bin}(n, p)$

$$P(X = i) = \binom{n}{i} p^i (1 - p)^{n-i} \quad i = 0, 1, \dots, n$$

N.B., by Binomial theorem,  $\sum_{i=0}^n P(X = i) = 1$

Examples

# of heads in  $n$  coin flips

# of 1's in a randomly generated length  $n$  bit string

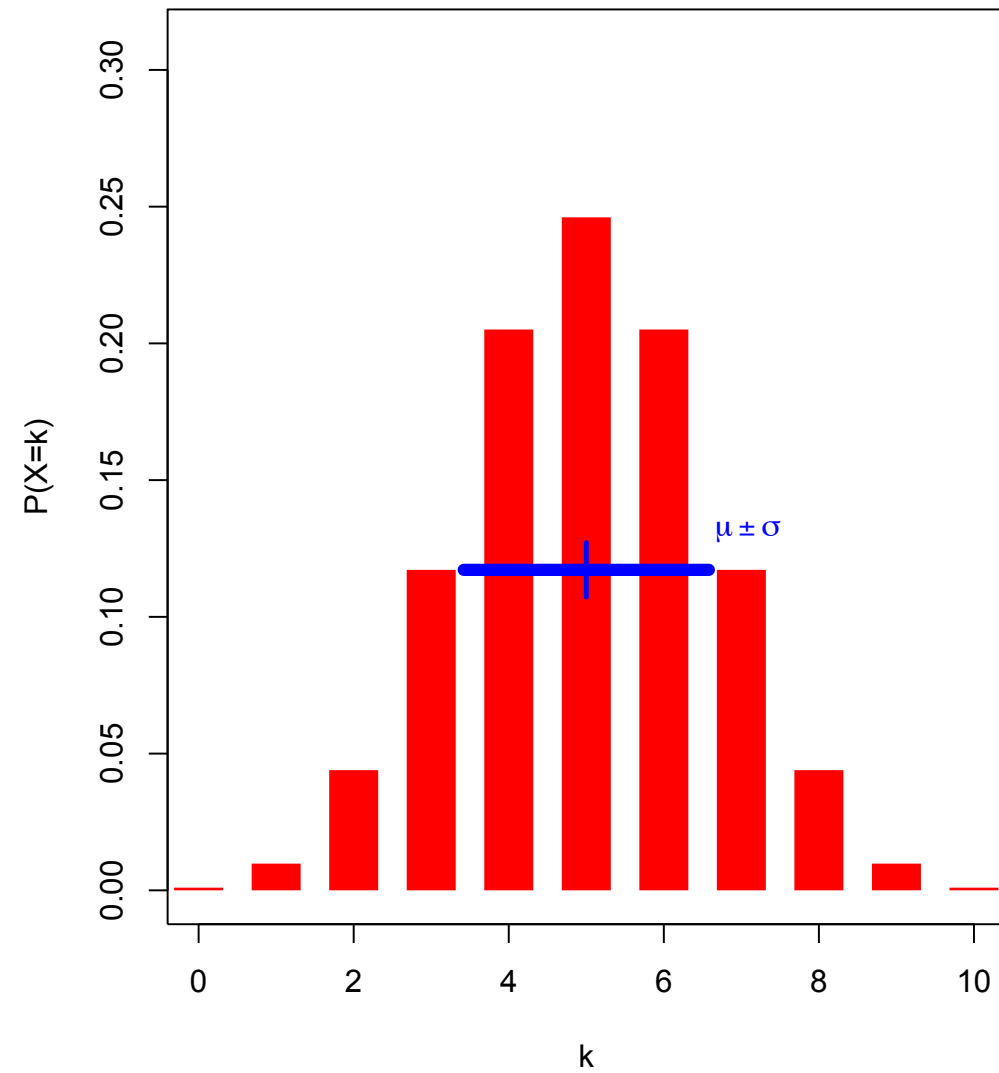
# of disk drive crashes in a 1000 computer cluster

$$E[X] = np$$

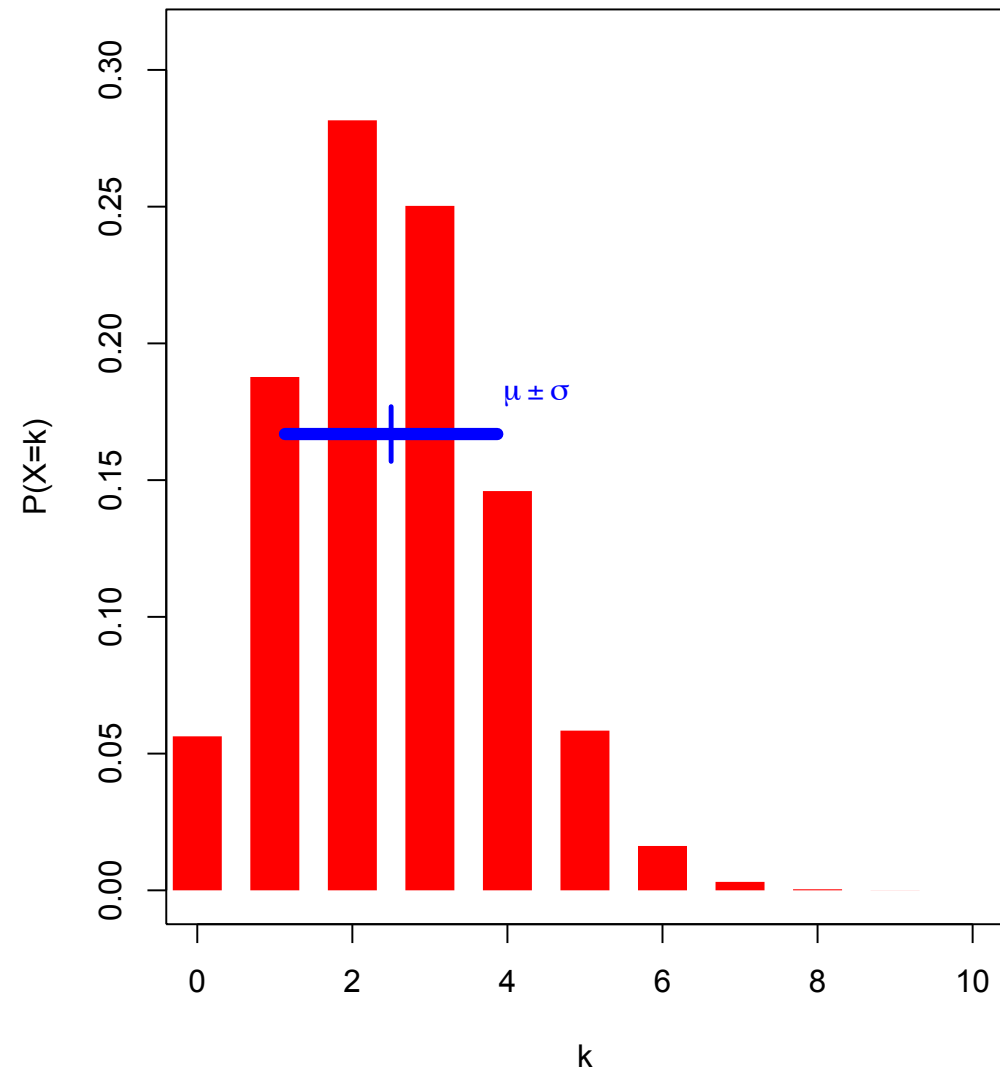
$$\text{Var}(X) = np(1-p)$$

←(proof below, twice)

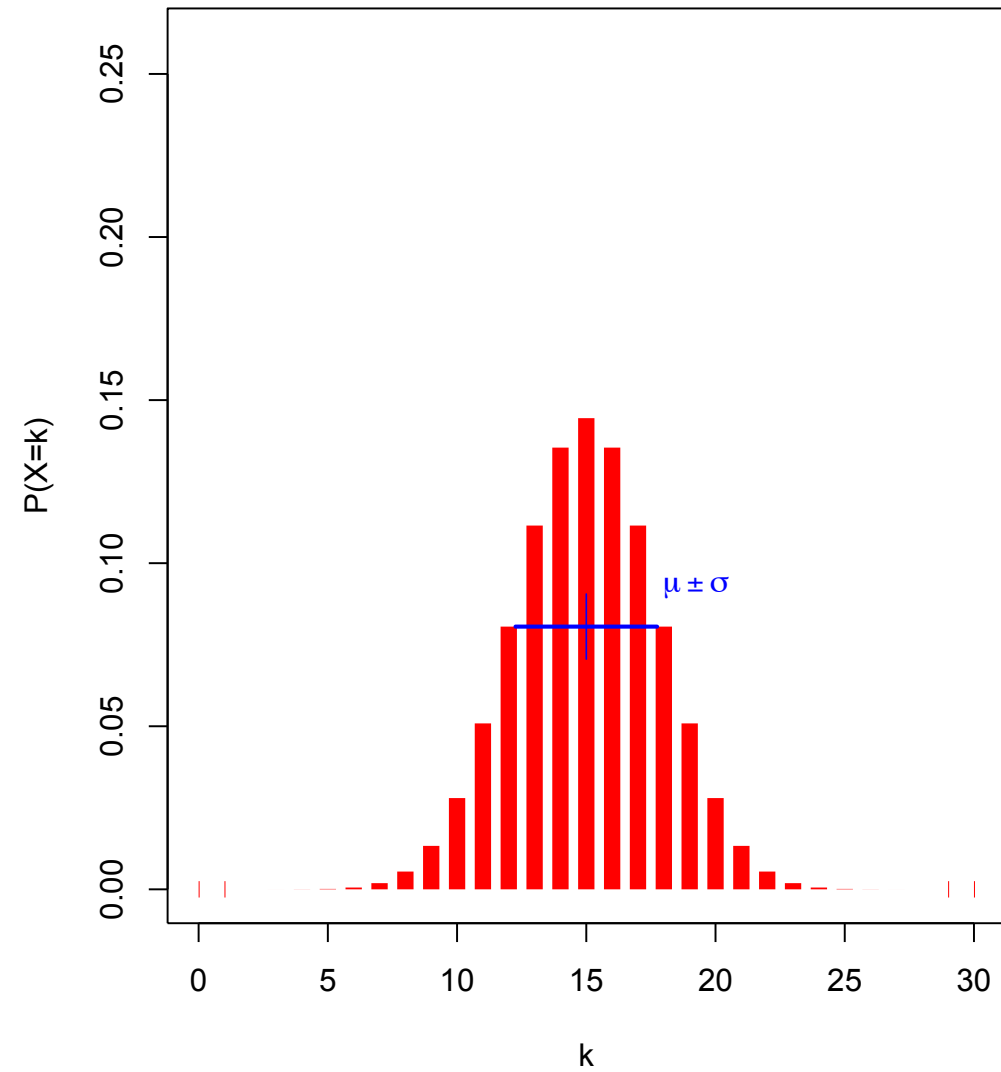
PMF for  $X \sim \text{Bin}(10, 0.5)$



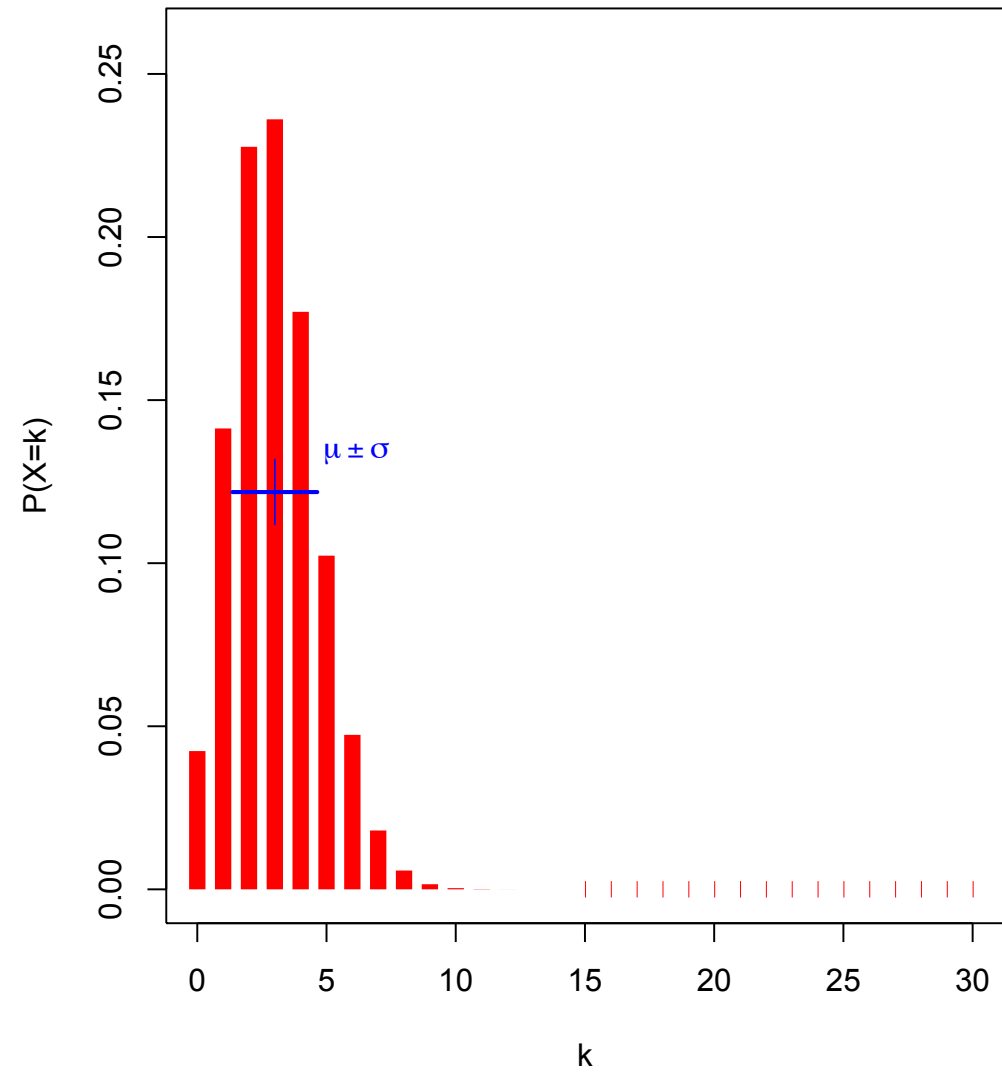
PMF for  $X \sim \text{Bin}(10, 0.25)$



PMF for  $X \sim \text{Bin}(30, 0.5)$



PMF for  $X \sim \text{Bin}(30, 0.1)$



# mean and variance of the binomial (I)

$$E[X^k] = \sum_{i=0}^n i^k \binom{n}{i} p^i (1-p)^{n-i}$$

👉 generalizes slide 15

$$= \sum_{i=1}^n i^k \binom{n}{i} p^i (1-p)^{n-i}$$

$$= np \sum_{i=1}^n i^{k-1} \binom{n-1}{i-1} p^{i-1} (1-p)^{n-i}$$

using  $i \binom{n}{i} = n \binom{n-1}{i-1}$

$$= np \sum_{j=0}^{n-1} (j+1)^{k-1} \binom{n-1}{j} p^j (1-p)^{n-1-j}$$

using  $j = i - 1$

$$= np E[(Y+1)^{k-1}]$$

using defn of  $E[\cdot]$ ,  
where  $Y \sim \text{Bin}(n-1, p)$

$$k = 1 \text{ gives: } \boxed{E[X] = np} ; \quad k = 2 \text{ gives: } \boxed{E[X^2] = np((n-1)p + 1)}$$

$$np (E[Y] + 1)$$

$$\begin{aligned} \text{Var}[X] &= E[X^2] - (E[X])^2 \\ &= np((n-1)p + 1) - (np)^2 \\ &= np(1-p) \end{aligned}$$

Theorem: If  $X$  &  $Y$  are *independent*, then  $E[X \cdot Y] = E[X] \cdot E[Y]$

Proof:

any dist, not just binomial

Let  $x_i, y_i, i = 1, 2, \dots$  be the possible values of  $X, Y$ .

$$\begin{aligned} E[X \cdot Y] &= \sum_i \sum_j x_i \cdot y_j \cdot P(X = x_i \wedge Y = y_j) \\ &= \sum_i \sum_j x_i \cdot y_j \cdot P(X = x_i) \cdot P(Y = y_j) \\ &= \sum_i x_i \cdot P(X = x_i) \cdot \left( \sum_j y_j \cdot P(Y = y_j) \right) \\ &= E[X] \cdot E[Y] \end{aligned}$$

independence

Note: *NOT* true in general; see earlier example  $E[X^2] \neq E[X]^2$



## variance of independent r.v.s is additive

(Bienaymé, 1853)

Theorem: If  $X$  &  $Y$  are *independent*, (any dist, not just binomial) then variance is *additive*:  $\text{Var}[X+Y] = \text{Var}[X] + \text{Var}[Y]$

Proof: Let

$$\begin{aligned}\hat{X} &= X - E[X] & \hat{Y} &= Y - E[Y] \\ E[\hat{X}] &= 0 & E[\hat{Y}] &= 0 \\ \text{Var}[\hat{X}] &= \text{Var}[X] & \text{Var}[\hat{Y}] &= \text{Var}[Y]\end{aligned}$$

$\text{Var}(aX+b) = a^2\text{Var}(X)$

$$\begin{aligned}\text{Var}[X + Y] &= \text{Var}[\hat{X} + \hat{Y}] \\ &= E[(\hat{X} + \hat{Y})^2] - (E[\hat{X} + \hat{Y}])^2 \\ &= E[\hat{X}^2 + 2\hat{X}\hat{Y} + \hat{Y}^2] - 0 \\ &= E[\hat{X}^2] + 2E[\hat{X}\hat{Y}] + E[\hat{Y}^2] \\ &= \text{Var}[\hat{X}] + 0 + \text{Var}[\hat{Y}] \\ &= \text{Var}[X] + \text{Var}[Y]\end{aligned}$$

previous slide

Q: Why are  $\hat{X}, \hat{Y}$  independent?  
A: See HW

## variance of independent r.v.s is additive

(Bienaymé, 1853)

Theorem: If  $X$  &  $Y$  are *independent*, (any dist, not just binomial) then

$$\text{Var}[X+Y] = \text{Var}[X] + \text{Var}[Y]$$

Alternate Proof:

$$\text{Var}[X + Y]$$

$$= E[(X + Y)^2] - (E[X + Y])^2$$

$$= E[X^2 + 2XY + Y^2] - (E[X] + E[Y])^2$$

$$= E[X^2] + 2E[XY] + E[Y^2] - ((E[X])^2 + 2E[X]E[Y] + (E[Y])^2)$$

$$= E[X^2] - (E[X])^2 + E[Y^2] - (E[Y])^2 + 2(E[XY] - E[X]E[Y])$$

$$= \text{Var}[X] + \text{Var}[Y] + 2(E[XY] - E[X]E[Y])$$

$$= \text{Var}[X] + \text{Var}[Y]$$

FYI, the quantity  $E[XY] - E[X]E[Y]$  is called the *covariance* of  $X, Y$ . As shown, it is 0 if  $X, Y$  are independent; if not zero it is a useful measure of their degree of dependence.

slide 60

## mean, variance of the binomial (II)

If  $Y_1, Y_2, \dots, Y_n \sim \text{Ber}(p)$  and independent,

then  $X = \sum_{i=1}^n Y_i \sim \text{Bin}(n, p)$ .

The  $Y_i$ 's are i.i.d.:  
Independent and  
Identically Distributed

$$E[X] = np$$

$$E[X] = E \left[ \sum_{i=1}^n Y_i \right] = \sum_{i=1}^n E[Y_i] = nE[Y_1] = np$$

$$\text{Var}[X] = np(1 - p)$$

$$\text{Var}[X] = \text{Var} \left[ \sum_{i=1}^n Y_i \right] = \sum_{i=1}^n \text{Var}[Y_i] = n\text{Var}[Y_1] = np(1 - p)$$

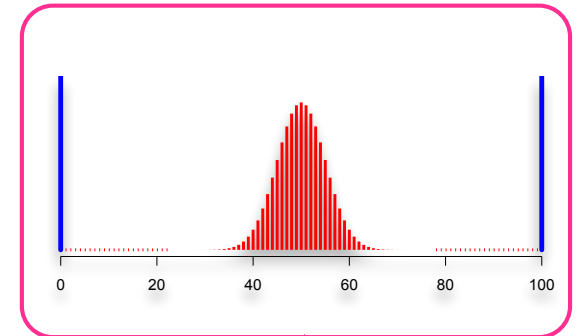
## mean, variance of the binomial (II)

If  $Y_1, Y_2, \dots, Y_n \sim \text{Ber}(p)$  and independent,

then  $X = \sum_{i=1}^n Y_i \sim \text{Bin}(n, p)$ .

$$E[X] = E\left[\sum_{i=1}^n Y_i\right] = nE[Y_1] = np$$

$$\text{Var}[X] = \text{Var}\left[\sum_{i=1}^n Y_i\right] = n\text{Var}[Y_1] = np(1-p)$$



Note :

$$E\left[\sum_{i=1}^n Y_i\right] = \sum_{i=1}^n E[Y_i] = nE[Y_7] = E[nY_7]$$

but

$$\text{Var}\left[\sum_{i=1}^n Y_i\right] = \sum_{i=1}^n \text{Var}[Y_i] = n\text{Var}[Y_7] \ll \text{Var}[nY_7] = n^2\text{Var}[Y_7]$$

Q. Why the big difference? A.

Indp random fluctuations tend to cancel when added; dependent ones may reinforce; “ $nY_7$ ”: no such cancelation; much variation

A RAID-like disk array consists of  $n$  drives, each of which will fail independently with probability  $p$ . Suppose it can operate effectively if at least one-half of its components function, e.g., by “majority vote.”



For what values of  $p$  is a 5-component system more likely to operate effectively than a 3-component system?

$X_5 = \# \text{ failed in 5-component system} \sim \text{Bin}(5, p)$

$X_3 = \# \text{ failed in 3-component system} \sim \text{Bin}(3, p)$

$X_5 = \# \text{ failed in 5-component system} \sim \text{Bin}(5, p)$

$X_3 = \# \text{ failed in 3-component system} \sim \text{Bin}(3, p)$

$P(\text{5 component system effective}) = P(X_5 < 5/2)$

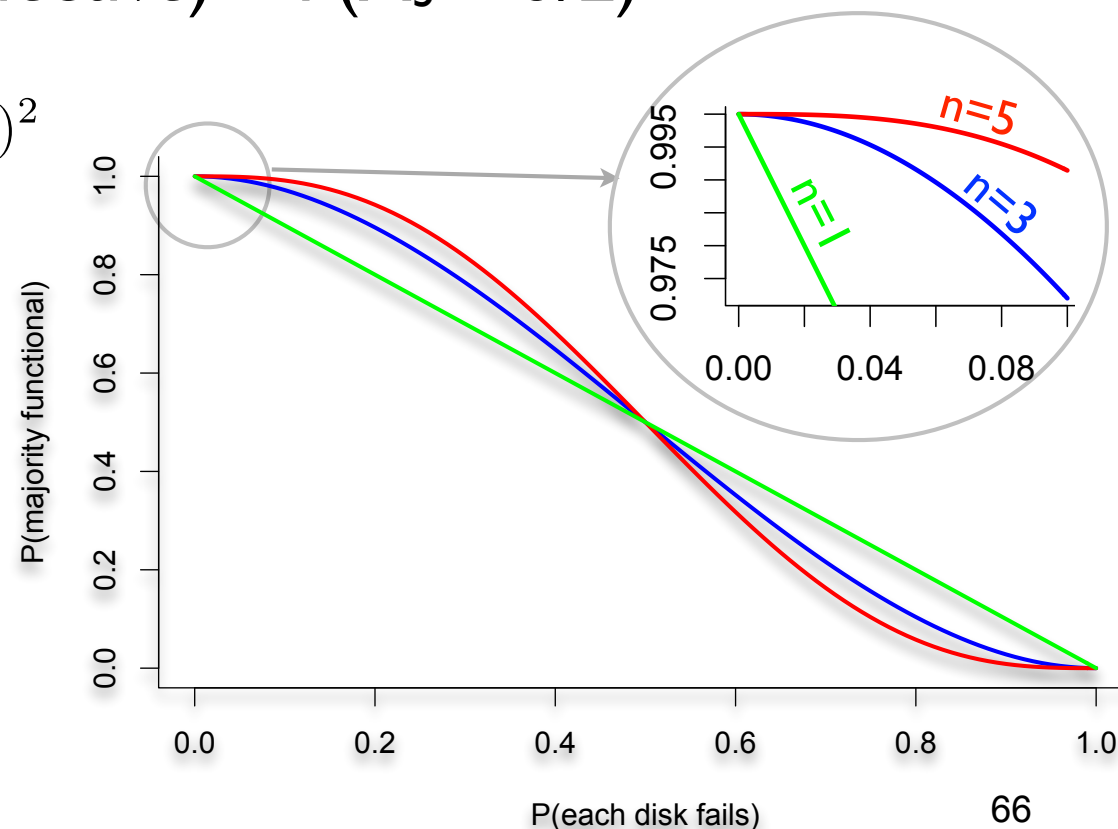
$$\binom{5}{0}p^0(1-p)^5 + \binom{5}{1}p^1(1-p)^4 + \binom{5}{2}p^2(1-p)^3$$

$P(\text{3 component system effective}) = P(X_3 < 3/2)$

$$\binom{3}{0}p^0(1-p)^3 + \binom{3}{1}p^1(1-p)^2$$

*Calculation:*

5-component system  
is better iff  $p < 1/2$



Goal: send a 4-bit message over a noisy communication channel.

Say, 1 bit in 10 is flipped in transit, independently.

What is the probability that the message arrives correctly?

Let  $X$  = # of errors;  $X \sim \text{Bin}(4, 0.1)$

$P(\text{correct message received}) = P(X=0)$

$$P(X = 0) = \binom{4}{0} (0.1)^0 (0.9)^4 = 0.6561$$

Can we do better? Yes: error correction via redundancy.

E.g., send every bit in triplicate; use majority vote.

Let  $Y$  = # of errors in one trio;  $Y \sim \text{Bin}(3, 0.1)$ ;  $P(\text{a trio is OK}) =$

$$P(Y \leq 1) = \binom{3}{0} (0.1)^0 (0.9)^3 + \binom{3}{1} (0.1)^1 (0.9)^2 = 0.972$$

If  $X'$  = # errors in triplicate msg,  $X' \sim \text{Bin}(4, 0.028)$ , and

$$P(X' = 0) = \binom{4}{0} (0.028)^0 (0.972)^4 = 0.8926168$$

The Hamming(7,4) code:

Have a 4-bit string to send over the network (or to disk)

Add 3 “parity” bits, and send 7 bits total

If bits are  $b_1b_2b_3b_4$  then the three parity bits are

$$\text{parity}(b_1b_2b_3), \text{parity}(b_1b_3b_4), \text{parity}(b_2b_3b_4)$$

Each bit is independently corrupted (flipped) in transit with probability 0.1

$Z = \text{number of bits corrupted} \sim \text{Bin}(7, 0.1)$

The Hamming code allow us to *correct* all 1 bit errors.

(E.g., if  $b_1$  flipped, 1st 2 parity bits, but not 3rd, will look wrong; the only single bit error causing this symptom is  $b_1$ . Similarly for any other single bit being flipped. Some, but not all, multi-bit errors can be detected, but not corrected.)

$$P(\text{correctable message received}) = P(Z \leq 1)$$



“Parity(x,y,z)” is perhaps best defined as  $(x+y+z+1) \bmod 2$

I.e., make sure that there are an odd number of one-bits among x,y,z,parity. Why?  
“Stuck at zero” faults are a common error mode in digital systems, so it’s best if the parity check on 000 is 1. I.e., 0001 is OK but 0000 would be recognized as faulty.

Suppose the message you want to send is ‘1011’

Instead, you send ‘1011 1 0 1’ (via rules on prev slide)

If your partner receives a 1-bit corruption of this, e.g.,

**0**011 1 0 1

then both underlined parity bits are incorrect: the quadruples defined above (incl the parity bit) have even parity, but should have odd parity. Studying the rules on the prev slide, this is the *ONLY* single bit corruption displaying this pattern, so you know to “correct” the initial 0 bit to 1, recovering the 1011 message.

Exercise: try all 6 other single bit errors; you should see that each has a distinct pattern of “parity errors,” hence is correctable. (But 2 or more errors leave you in deep doo doo...)

Using Hamming error-correcting codes:  $Z \sim \text{Bin}(7, 0.1)$

$$P(Z \leq 1) = \binom{7}{0} (0.1)^0 (0.9)^7 + \binom{7}{1} (0.1)^1 (0.9)^6 \approx 0.8503$$

Recall, uncorrected success rate is

$$P(X = 0) = \binom{4}{0} (0.1)^0 (0.9)^4 = 0.6561$$

And triplicate code success rate is:

$$P(X' = 0) = \binom{4}{0} (0.028)^0 (0.972)^4 = 0.8926168$$

Hamming code is nearly as reliable as the triplicate code, with  $5/12 \approx 42\%$  fewer bits. (& better with longer codes; overhead is  $O(\log n)$  bits for  $n$  bit messages.)

Sending a bit string over the network

$n = 4$  bits sent, each corrupted with probability  $0.1$

$X = \#$  of corrupted bits,  $X \sim \text{Bin}(4, 0.1)$

In real networks, large bit strings (length  $n \approx 10^4$ )

Corruption probability is very small:  $p \approx 10^{-6}$

$X \sim \text{Bin}(10^4, 10^{-6})$  is unwieldy to compute

Extreme  $n$  and  $p$  values arise in many cases

- # bit errors in file written to disk

- # of typos in a book

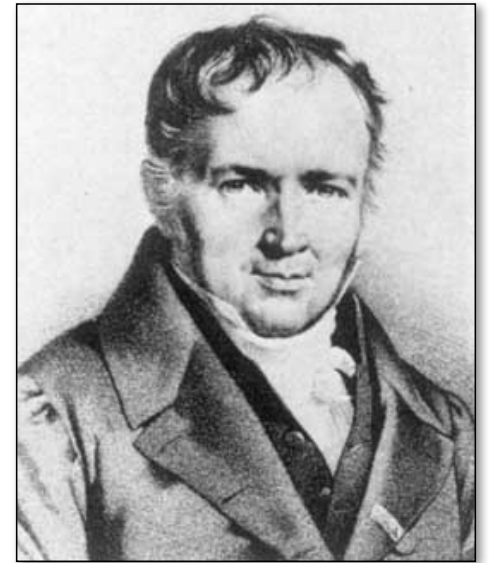
- # of elements in particular bucket of large hash table

- # of server crashes per day in giant data center

- # facebook login requests sent to a particular server

Suppose “events” happen, independently, at an *average* rate of  $\lambda$  per unit time. Let  $X$  be the *actual* number of events happening in a given time unit. Then  $X$  is a *Poisson* r.v. with *parameter*  $\lambda$  (denoted  $X \sim \text{Poi}(\lambda)$ ) and has distribution (PMF):

$$P(X = i) = e^{-\lambda} \frac{\lambda^i}{i!}$$



Siméon Poisson, 1781-1840

Examples:

# of alpha particles emitted by a lump of radium in 1 sec.

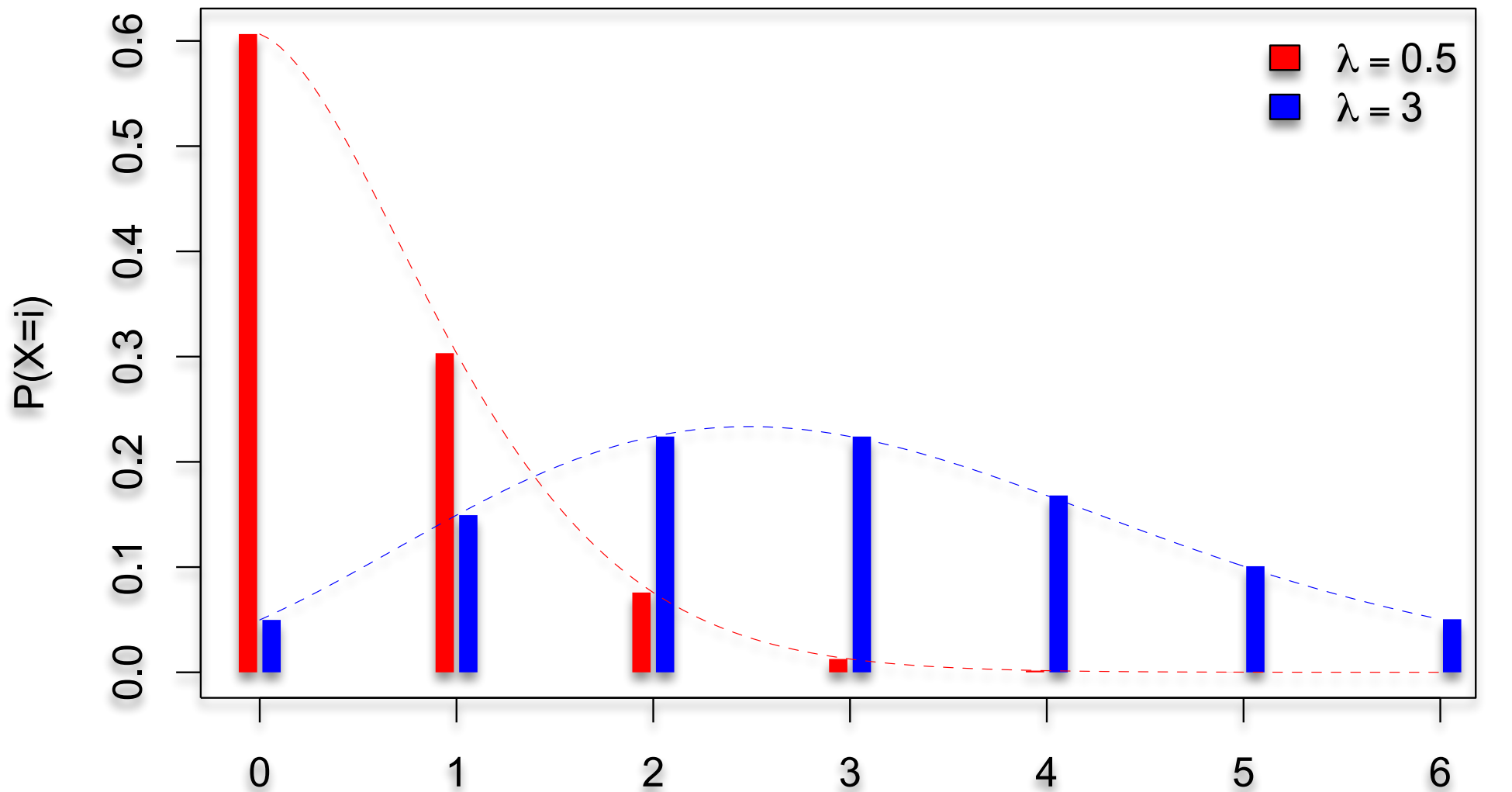
# of traffic accidents in Seattle in one year

# of babies born in a day at UW Med center

# of visitors to my web page today

See B&T Section 6.2 for more on theoretical basis for Poisson.

$$P(X = i) = e^{-\lambda} \frac{\lambda^i}{i!}$$



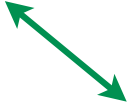
$X$  is a Poisson r.v. with parameter  $\lambda$  if it has PMF:

$$P(X = i) = e^{-\lambda} \frac{\lambda^i}{i!}$$

Is it a valid distribution? Recall Taylor series:

$$e^{\lambda} = \frac{\lambda^0}{0!} + \frac{\lambda^1}{1!} + \dots = \sum_{0 \leq i} \frac{\lambda^i}{i!}$$

So

$$\sum_{0 \leq i} P(X = i) = \sum_{0 \leq i} e^{-\lambda} \frac{\lambda^i}{i!} = e^{-\lambda} \sum_{0 \leq i} \frac{\lambda^i}{i!} = e^{-\lambda} e^{\lambda} = 1$$


## expected value of poisson r.v.s

$$\begin{aligned} E[X] &= \sum_{0 \leq i} i \cdot e^{-\lambda} \frac{\lambda^i}{i!} \\ &= \sum_{1 \leq i} i \cdot e^{-\lambda} \frac{\lambda^i}{i!} && \text{ } i = 0 \text{ term is zero} \\ &= \lambda e^{-\lambda} \sum_{1 \leq i} \frac{\lambda^{i-1}}{(i-1)!} && j = i-1 \\ &= \lambda e^{-\lambda} \sum_{0 \leq j} \frac{\lambda^j}{j!} \\ &= \lambda e^{-\lambda} e^{\lambda} \\ &= \lambda && \text{As expected, given definition in terms of "average rate } \lambda \text{"} \end{aligned}$$

(Var[X] =  $\lambda$ , too; proof similar, see B&T example 6.20)

## binomial random variable is poisson in the limit

---

Poisson approximates binomial when  $n$  is large,  $p$  is small, and  $\lambda = np$  is “moderate”

Different interpretations of “moderate,” e.g.

$n > 20$  and  $p < 0.05$

$n > 100$  and  $p < 0.1$

Formally, Binomial is Poisson in the limit as

$n \rightarrow \infty$  (equivalently,  $p \rightarrow 0$ ) while holding  $np = \lambda$



## binomial $\rightarrow$ poisson in the limit

---

$X \sim \text{Binomial}(n, p)$

$$\begin{aligned} P(X = i) &= \binom{n}{i} p^i (1 - p)^{n-i} \\ &= \frac{n!}{i!(n-i)!} \left(\frac{\lambda}{n}\right)^i \left(1 - \frac{\lambda}{n}\right)^{n-i}, \text{ where } \lambda = pn \\ &= \frac{n(n-1) \cdots (n-i+1)}{n^i} \frac{\lambda^i}{i!} \frac{(1 - \lambda/n)^n}{(1 - \lambda/n)^i} \\ &= \underbrace{\frac{n(n-1) \cdots (n-i+1)}{(n-\lambda)^i}}_{\approx 1} \cdot \frac{\lambda^i}{i!} \cdot \underbrace{(1 - \lambda/n)^n}_{\approx e^{-\lambda}} \\ &\approx 1 \cdot \frac{\lambda^i}{i!} \cdot e^{-\lambda} \end{aligned}$$

I.e., Binomial  $\approx$  Poisson for large  $n$ , small  $p$ , moderate  $i$ ,  $\lambda$ .

Handy: Poisson has only 1 parameter—the expected # of successes

## sending data on a network, again

---

Recall example of sending bit string over a network

Send bit string of length  $n = 10^4$

Probability of corruption is  $p = 10^{-6}$  per bit (independent)

What is probability that message arrives uncorrupted?

Binomial Model:

Number of corrupt bits  $Y \sim \text{Bin}(10^4, 10^{-6})$ :

$$P(Y=0) \approx 0.990049829$$

Poisson Approximation (where “unit time” =  $10^4$  bits):

Number of corrupt bits  $X \sim \text{Poi}(\lambda = 10^4 \cdot 10^{-6} = 0.01)$

$$P(X = 0) = e^{-\lambda} \frac{\lambda^0}{0!} = e^{-0.01} \frac{0.01^0}{0!} \approx 0.990049834$$

I.e., Poisson approx (here) is accurate to  $\sim 5$  parts per billion

## Hamming codes and reality

Remember Hamming? Generalized Hamming code adds 14 code bits to correct 1 error in 10000 bit message. Message arrives correctly, or is correctable, if at most 1 bit is in error.

*Binomial model:*

Number of corrupt bits  $Z \sim \text{Bin}(10014, \epsilon)$ ,  $\epsilon = 10^{-6}$

$$P(Z \leq 1) = \binom{10014}{0} \epsilon^0 (1 - \epsilon)^{10014} + \binom{10014}{1} \epsilon^1 (1 - \epsilon)^{10013} = 0.999950198$$

*Poisson approximation:*

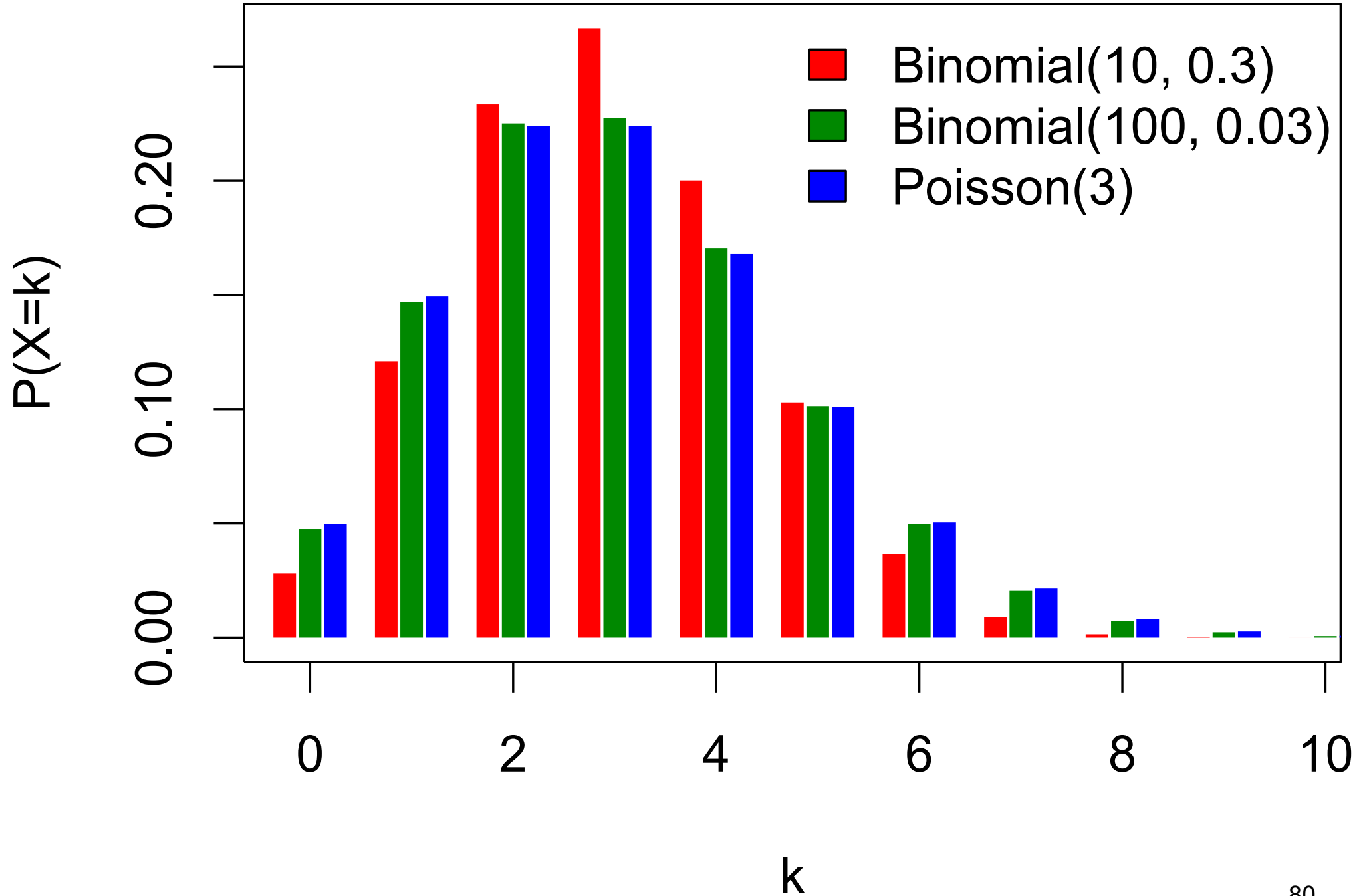
Number of corrupt bits  $Z' \sim \text{Poi}(0.010014)$

$$P(Z' \leq 1) = e^{-0.010014} \left( \frac{(0.010014)^0}{0!} + \frac{(0.010014)^1}{1!} \right) = 0.999950193$$

Two takeaways:

1. Again, Poisson approximation is good to about 5 ppb
2. And 0.14% overhead yields  $\approx 200\times$  reduction in fraction of erroneous messages (.5 in  $10^4$  vs 1 in  $10^2$ )

## binomial vs poisson



## expectation and variance of a poisson

---

Recall: if  $Y \sim \text{Bin}(n,p)$ , then:

$$E[Y] = np$$

$$\text{Var}[Y] = np(1-p)$$

And if  $X \sim \text{Poi}(\lambda)$  where  $\lambda = np$  ( $n \rightarrow \infty, p \rightarrow 0$ ) then

$$E[X] = \lambda = np = E[Y]$$

$$\text{Var}[X] = \lambda \approx \lambda(1-\lambda/n) = np(1-p) = \text{Var}[Y]$$

Expectation and variance of Poisson are the same ( $\lambda$ )

Expectation is the same as corresponding binomial

Variance almost the same as corresponding binomial

Note: when two different distributions share the same mean & variance, it suggests (but doesn't prove) that one may be a good approximation for the other.

Suppose a server can process 2 requests per second

Requests arrive at random at an average rate of 1/sec

Unprocessed requests are held in a *buffer*

*Q. How big a buffer do we need to avoid ever dropping a request?*

**A. Infinite**

*Q. How big a buffer do we need to avoid dropping a request more often than once a day?*

**A. (approximate)** If  $X$  is the number of arrivals in a second, then  $X$  is Poisson ( $\lambda=1$ ). We want  $b$  s.t.

$$P(X > b) < 1/(24*60*60) \approx 1.2 \times 10^{-5}$$

$$P(X = b) = e^{-1}/b! \quad \sum_{i \geq 8} P(X=i) \approx P(X=8) \approx 10^{-5}, \text{ so } b \approx 8$$

Above necessary but not sufficient; also check prob of 10 arrivals in 2 seconds, 12 in 3, etc.  
See BT p366 for a possible approach to fully solving it.

In a series  $X_1, X_2, \dots$  of Bernoulli trials with success probability  $p$ , let  $Y$  be the index of the first success, i.e.,

$$X_1 = X_2 = \dots = X_{Y-1} = 0 \text{ \& } X_Y = 1$$

Then  $Y$  is a *geometric* random variable with parameter  $p$ .

Examples:

Number of coin flips until first head


Number of blind guesses on LSAT until I get one right

Number of darts thrown until you hit a bullseye

Number of random probes into hash table until empty slot

Number of wild guesses at a password until you hit it

$$P(Y=k) = (1-p)^{k-1}p; \text{ Mean } 1/p; \text{ Variance } (1-p)/p^2$$

 see [slide 14](#); see also [slide 86](#),  
BT p105 for slick alt. proof

## interlude: more on conditioning

---

Recall: conditional probability

$$P(X | A) = P(X \& A)/P(A)$$

A note about notation: When  $X$  is an r.v., take this as either shorthand for “ $\forall x P(X=x \dots)$ ” or as defining the conditional PMF  $p(x|A)$  from the joint PMF

Conditional probability is a probability, i.e.

1. it's nonnegative
2. it's normalized
3. it's happy with the axioms, etc.

Define: The *conditional expectation* of  $X$

$$E[X | A] = \sum_x x \cdot p(X = x | A)$$

I.e., the value of r.v.  $X$  averaged over outcomes where I know event  $A$  happened



Recall: the law of total probability

$$p(X) = p(X | A) \cdot P(A) + p(X | A^c) \cdot P(A^c)$$

Again,  
“ $\forall x P(X=x \dots)$ ” or  
“unconditional PMF  
is weighted avg of  
conditional PMFs”  
←

I.e., unconditional probability is the weighted average of conditional probabilities, weighted by the probabilities of the conditioning events

## The Law of Total Expectation

$$E[X] = E[X | A] \cdot P(A) + E[X | A^c] \cdot P(A^c)$$

I.e., unconditional expectation is the weighted average of conditional expectations, weighted by the probabilities of the conditioning events

## Proof of the Law of Total Expectation:

$$\begin{aligned} E[X] &= \sum_x xp(x) \\ &= \sum_x x(p(x \mid A)P(A) + p(x \mid \bar{A})P(\bar{A})) \\ &= \sum_x xp(x \mid A)P(A) + \sum_x xp(x \mid \bar{A})P(\bar{A}) \\ &= \left( \sum_x xp(x \mid A) \right) P(A) + \left( \sum_x xp(x \mid \bar{A}) \right) P(\bar{A}) \\ &= E[X \mid A]P(A) + E[X \mid \bar{A}]P(\bar{A}) \end{aligned}$$

$$X \sim \text{geo}(p)$$

$$E[X] = E[X \mid X=1] \cdot P(X=1) + E[X \mid X>1] \cdot P(X>1)$$

$$= 1 \cdot p + (1 + E[X]) \cdot (1-p)$$

$\vdots$   simple algebra

$$E[X] = 1/p$$

memorylessness: after flipping one tail, *remaining* waiting time until 1<sup>st</sup> head is exactly the same as starting from scratch

E.g., if  $p=1/2$ , expect to wait 2 flips for 1<sup>st</sup> head;  
 $p=1/10$ , expect to wait 10 flips.

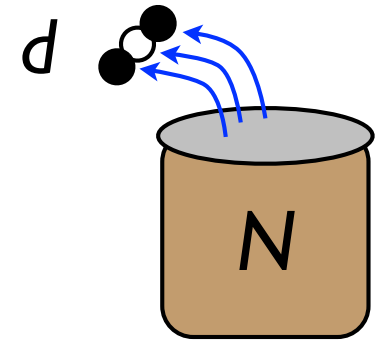
(Similar derivation for variance:  $(1-p)/p^2$ )

# balls in urns – the hypergeometric distribution

B&T, exercise 1.61

Draw  $d$  balls (without replacement) from an urn containing  $N$ , of which  $w$  are white, the rest black.

Let  $X$  = number of white balls drawn



$$P(X = i) = \frac{\binom{w}{i} \binom{N-w}{d-i}}{\binom{N}{d}}, \quad i = 0, 1, \dots, d$$

[note:  $\binom{n}{k} = 0$  if  $k < 0$  or  $k > n$ ]

$E[X] = dp$ , where  $p = w/N$  (the fraction of white balls)

proof: Let  $X_j$  be 0/1 indicator for  $j$ -th ball is white,  $X = \sum X_j$

The  $X_j$  are *dependent*, but  $E[X] = E[\sum X_j] = \sum E[X_j] = dp$

$\text{Var}[X] = dp(1-p)(1-(d-1)/(N-1))$

like  
binomial  
(almost)

$N \approx 22500$  human genes, many of unknown function

Suppose in some experiment,  $d = 1588$  of them were observed (say, they were all switched on in response to some drug)

A big question: What are they doing?

One idea: The Gene Ontology Consortium ([www.geneontology.org](http://www.geneontology.org)) has grouped genes with known functions into categories such as “muscle development” or “immune system.” Suppose 26 of your  $d$  genes fall in the “muscle development” category.

Just chance?

Or call Coach (& see if he wants to dope some athletes)?

Hypergeometric: GO has 116 genes in the muscle development category. If those are the white balls among 22500 in an urn, what is the probability that you would see 26 of them in 1588 draws?

**Table 2. Gene Ontology Analysis on Differentially Bound Peaks in Myoblasts versus Myotubes**

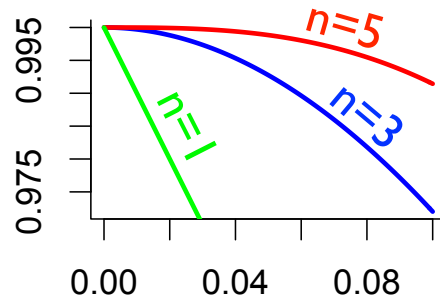
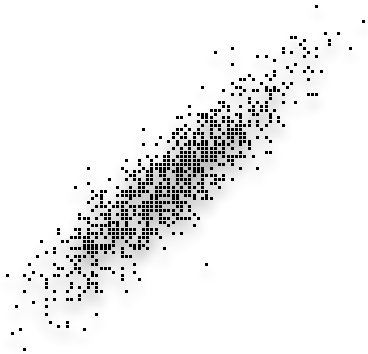
GO Categories Enriched in Genes Associated with Myotube-Increased Peaks

GOID	Term	P Value	OR <sup>a</sup>	Count <sup>b</sup>	Size <sup>c</sup>	Ont <sup>d</sup>
GO:0005856	cytoskeleton	2.05E-11	2.40	94	490	CC
GO:0043292	contractile fiber	6.98E-09	5.85	22	58	CC
GO:0030016	myofibril	1.96E-08	5.74	21	56	CC
GO:0044449	contractile fiber part	2.58E-08	5.97	20	52	CC
GO:0030017	sarcomere	4.95E-08	6.04	19	49	CC
GO:0008092	probability of seeing this many genes from a set of this size by chance according to the hypergeometric distribution.					MF
GO:0007519	skeletal muscle development	2.50E-16	4.13	20	65	BP
GO:0015629	actin cytoskeleton	4.73E-06	3.08	27	111	CC
GO:0003779	actin binding	1.13E-06	3.08	27	159	MF
GO:0006936	muscle contraction	1.35E-05	3.35	15	111	BP
GO:0044430	cytoskeleton part	1.35E-05	3.35	15	294	CC
GO:0031674	I band	2.27E-05	5.67	12	32	CC
GO:0003012	muscle system process	2.54E-05	4.11	16	52	BP
GO:0030029	actin filament-based process	2.89E-05	2.73	27	119	BP
GO:0007517	muscle development	5.06E-05	2.69	26	116	BP

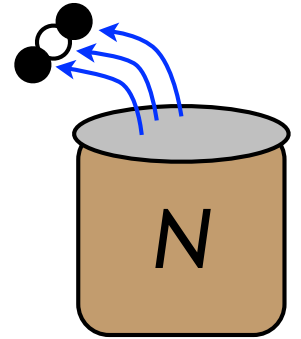
probability of seeing this many genes from  
a set of this size by chance according to  
the hypergeometric distribution.

E.g., if you draw 1588 balls from an urn containing 490 white balls  
and  $\approx 22000$  black balls,  $P(94 \text{ white}) \approx 2.05 \times 10^{-11}$

So, are genes flagged by this experiment specifically related to muscle development? This doesn't prove that they are, but it does say that there is an exceedingly small probability that so many would cluster in the "muscle development" group purely by chance.

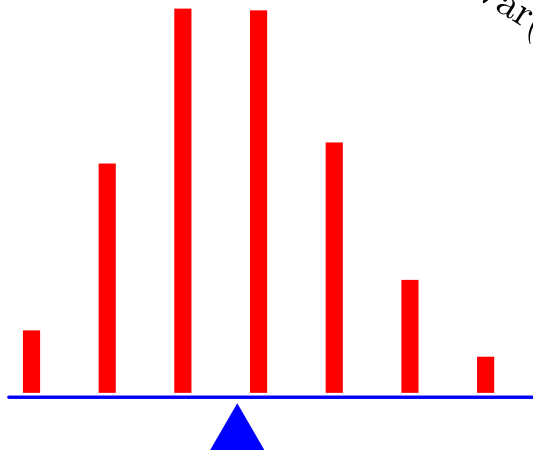


$$E[X+Y] = E[X] + E[Y]$$



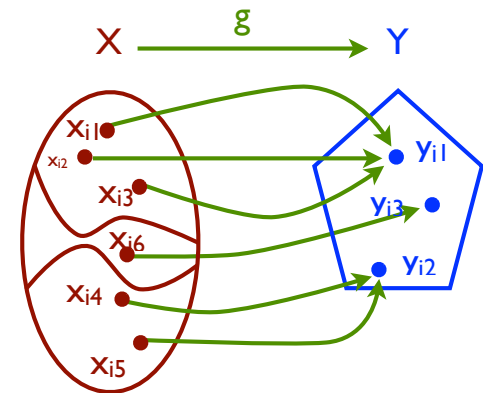
$$\sum_{i=-\infty}^{\infty} =$$

$$P(X=i) = e^{-\lambda} \frac{\lambda^i}{i!}$$



$$\text{Var}(X) = E[X^2] - (E[X])^2$$

$$\text{Var}[aX+b] = a^2 \text{Var}[X]$$



## random variables – summary

---

**RV:** a numeric function of the outcome of an experiment

**Probability Mass Function**  $p(x)$ : prob that  $RV = x$ ;  $\sum p(x) = 1$

**Cumulative Distribution Function**  $F(x)$ : probability that  $RV \leq x$

Generalize to *joint* distributions; *independence* & *marginals*

Expectation:

mean, average, “center of mass,” fair price for a game of chance

of a random variable:  $E[X] = \sum_x xp(x)$

of a function: if  $Y = g(X)$ , then  $E[Y] = \sum_x g(x)p(x)$

← (probability)-weighted average

linearity:

$$E[aX + b] = aE[X] + b$$

$$E[X+Y] = E[X] + E[Y]; \text{ even if dependent}$$

this interchange of “order of operations” is quite special to linear combinations. E.g.,  $E[XY] \neq E[X] \cdot E[Y]$ , in general (but see below)



Conditional Expectation:

$$E[X \mid A] = \sum_x x \cdot P(X=x \mid A)$$

Law of Total Expectation

$$E[X] = E[X \mid A] \cdot P(A) + E[X \mid \neg A] \cdot P(\neg A)$$

Variance:

$$\text{Var}[X] = E[(X - E[X])^2] = E[X^2] - (E[X])^2$$

$$\text{Standard deviation: } \sigma = \sqrt{\text{Var}[X]}$$

$$\text{Var}[aX+b] = a^2 \text{Var}[X] \quad \text{“Variance is insensitive to location, quadratic in scale”}$$

If  $X$  &  $Y$  are *independent*, then

$$E[X \cdot Y] = E[X] \cdot E[Y]$$

$$\text{Var}[X+Y] = \text{Var}[X] + \text{Var}[Y]$$

} (These two equalities hold for *indp* rv's; but not in general.)

### Important Examples:

Uniform(a,b):  $P(X = i) = \frac{1}{b - a + 1}$        $\mu = \frac{a + b}{2}, \sigma^2 = \frac{(b - a)(b - a + 1)}{12}$

Bernoulli:  $P(X = 1) = p, P(X = 0) = 1 - p$      $\mu = p, \sigma^2 = p(1 - p)$

Binomial:  $P(X = i) = \binom{n}{i} p^i (1 - p)^{n - i}$      $\mu = np, \sigma^2 = np(1 - p)$

Poisson:  $P(X = i) = e^{-\lambda} \frac{\lambda^i}{i!}$        $\mu = \lambda, \sigma^2 = \lambda$

$Bin(n, p) \approx Poi(\lambda)$  where  $\lambda = np$  fixed,  $n \rightarrow \infty$  (and so  $p = \lambda/n \rightarrow 0$ )

Geometric  $P(X = k) = (1 - p)^{k-1} p$        $\mu = 1/p, \sigma^2 = (1 - p)/p^2$

Many others, e.g., hypergeometric, negative binomial, ...

# Poisson distributions have no value over negative numbers

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