

independence



Defn: Two events E and F are *independent* if

$$P(EF) = P(E) P(F)$$

If $P(F) > 0$, this is equivalent to: $P(E|F) = P(E)$

Otherwise, they are called *dependent*

Roll two dice, yielding values D_1 and D_2

$$1) E = \{ D_1 = 1 \}$$

$$F = \{ D_2 = 1 \}$$

$$P(E) = 1/6, P(F) = 1/6, P(EF) = 1/36$$

$$P(EF) = P(E) \cdot P(F) \Rightarrow E \text{ and } F \text{ independent}$$

Intuitive; the two dice are not physically coupled

$$2) G = \{ D_1 + D_2 = 5 \} = \{(1,4),(2,3),(3,2),(4,1)\}$$

$$P(E) = 1/6, P(G) = 4/36 = 1/9, P(EG) = 1/36$$

not independent!

E, G are dependent events

The dice are still not physically coupled, but “ $D_1 + D_2 = 5$ ” couples them mathematically: info about D_1 constrains D_2 . (But dependence/independence not always intuitively obvious; “use the definition, Luke”.)



Two events E and F are *independent* if

$$P(EF) = P(E) P(F)$$

If $P(F) > 0$, this is equivalent to: $P(E|F) = P(E)$

Otherwise, they are called *dependent*

Three events E, F, G are independent if

$$P(EF) = P(E) P(F)$$

$$P(EG) = P(E) P(G) \quad \text{and} \quad P(EFG) = P(E) P(F) P(G)$$

$$P(FG) = P(F) P(G)$$

Example: Let X, Y be each $\{-1, 1\}$ with equal prob

$$E = \{X = 1\}, F = \{Y = 1\}, G = \{XY = 1\}$$

$$P(EF) = P(E)P(F), P(EG) = P(E)P(G), P(FG) = P(F)P(G)$$

but $P(EFG) = 1/4$!!! (because $P(G|EF) = 1$)

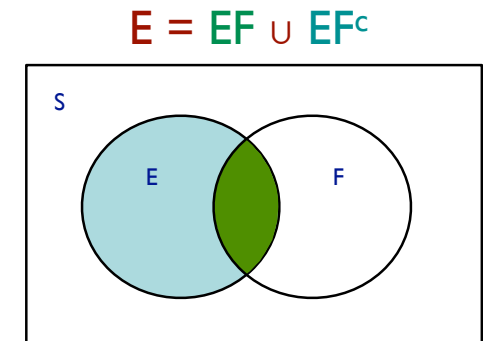
In general, events E_1, E_2, \dots, E_n are independent if for *every subset* S of $\{1, 2, \dots, n\}$, we have

$$P\left(\bigcap_{i \in S} E_i\right) = \prod_{i \in S} P(E_i)$$

(Sometimes this property holds only for small subsets S . E.g., E, F, G on the previous slide are *pairwise* independent, but not fully independent.)

Theorem: E, F independent $\Rightarrow E, F^c$ independent

Proof:
$$\begin{aligned} P(EF^c) &= P(E) - P(EF) \\ &= P(E) - P(E)P(F) \\ &= P(E)(1 - P(F)) \\ &= P(E)P(F^c) \end{aligned}$$



Theorem: if $P(E) > 0, P(F) > 0$, then

$$E, F \text{ independent} \Leftrightarrow P(E|F) = P(E) \Leftrightarrow P(F|E) = P(F)$$

Proof: Note $P(EF) = P(E|F)P(F)$, regardless of in/dep.
Assume independent. Then

$$P(E)P(F) = P(EF) = P(E|F)P(F) \Rightarrow P(E|F) = P(E) \quad (\div \text{ by } P(F))$$

$$\text{Conversely, } P(E|F) = P(E) \Rightarrow P(E)P(F) = P(EF) \quad (\times \text{ by } P(F))$$

Suppose a biased coin comes up heads with probability p ,
independent of other flips

$$P(n \text{ heads in } n \text{ flips}) = p^n$$

$$P(n \text{ tails in } n \text{ flips}) = (1-p)^n$$

$$P(\text{exactly } k \text{ heads in } n \text{ flips}) = \binom{n}{k} p^k (1-p)^{n-k}$$

Aside: note that the probability of some number of heads = $\sum_k \binom{n}{k} p^k (1-p)^{n-k} = (p + (1-p))^n = 1$
as it should, by the binomial theorem.



Suppose a biased coin comes up heads with probability p , *independent* of other flips

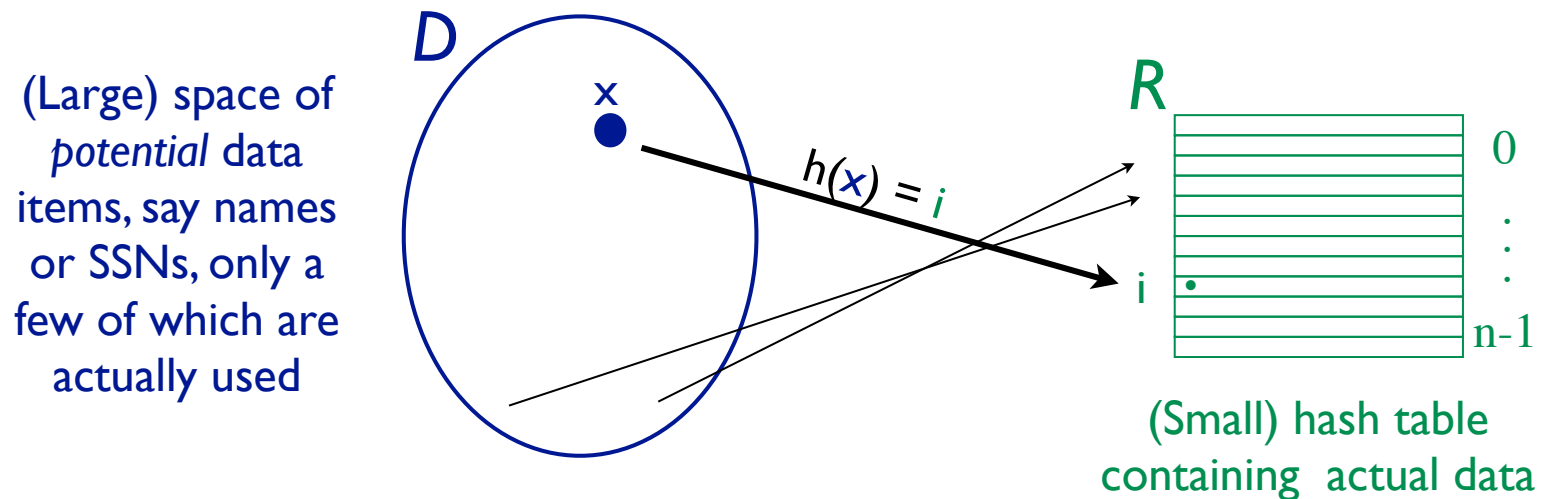


$$P(\text{exactly } k \text{ heads in } n \text{ flips}) = \binom{n}{k} p^k (1 - p)^{n-k}$$

Note when $p=1/2$, this is the same result we would have gotten by considering n flips in the “equally likely outcomes” scenario. But $p \neq 1/2$ makes that inapplicable. Instead, the *independence* assumption allows us to conveniently assign a probability to each of the 2^n outcomes, e.g.:

$$\Pr(\text{HHTHTTT}) = p^2(1-p)p(1-p)^3 = p^{\#H}(1-p)^{\#T}$$

A data structure problem: *fast* access to *small* subset of data drawn from a *large* space.



A solution: *hash function* $h: D \rightarrow \{0, \dots, n-1\}$ crunches/scrambles names from large space into small one. E.g., if x is integer:

$$h(x) = x \bmod n$$

Good hash functions *approximately* randomize placement.

m strings hashed (uniformly) into a table with n buckets

Each string hashed is an *independent* trial

E = at least one string hashed to first bucket

What is $P(E)$?

Solution:

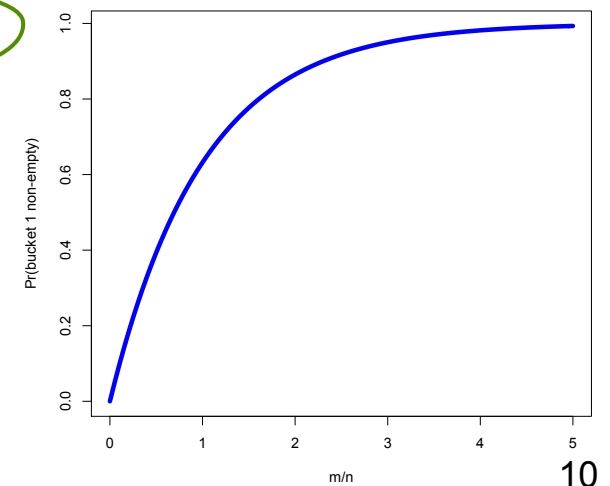
F_i = string i *not* hashed into first bucket ($i=1,2,\dots,m$)

$P(F_i) = 1 - 1/n = (n-1)/n$ for all $i=1,2,\dots,m$

Event $(F_1 F_2 \dots F_m)$ = no strings hashed to first bucket

$$\begin{aligned}
 P(E) &= 1 - P(F_1 F_2 \dots F_m) \\
 &= 1 - P(F_1) P(F_2) \dots P(F_m) \\
 &= 1 - ((n-1)/n)^m \\
 &\approx 1 - \exp(-m/n)
 \end{aligned}$$

indp



m strings hashed (non-uniformly) to table w/ n buckets

Each string hashed is an *independent* trial, with probability p_i of getting hashed to bucket i

E = At least 1 of buckets 1 to k gets ≥ 1 string

What is $P(E)$?

Solution:

F_i = at least one string hashed into i-th bucket

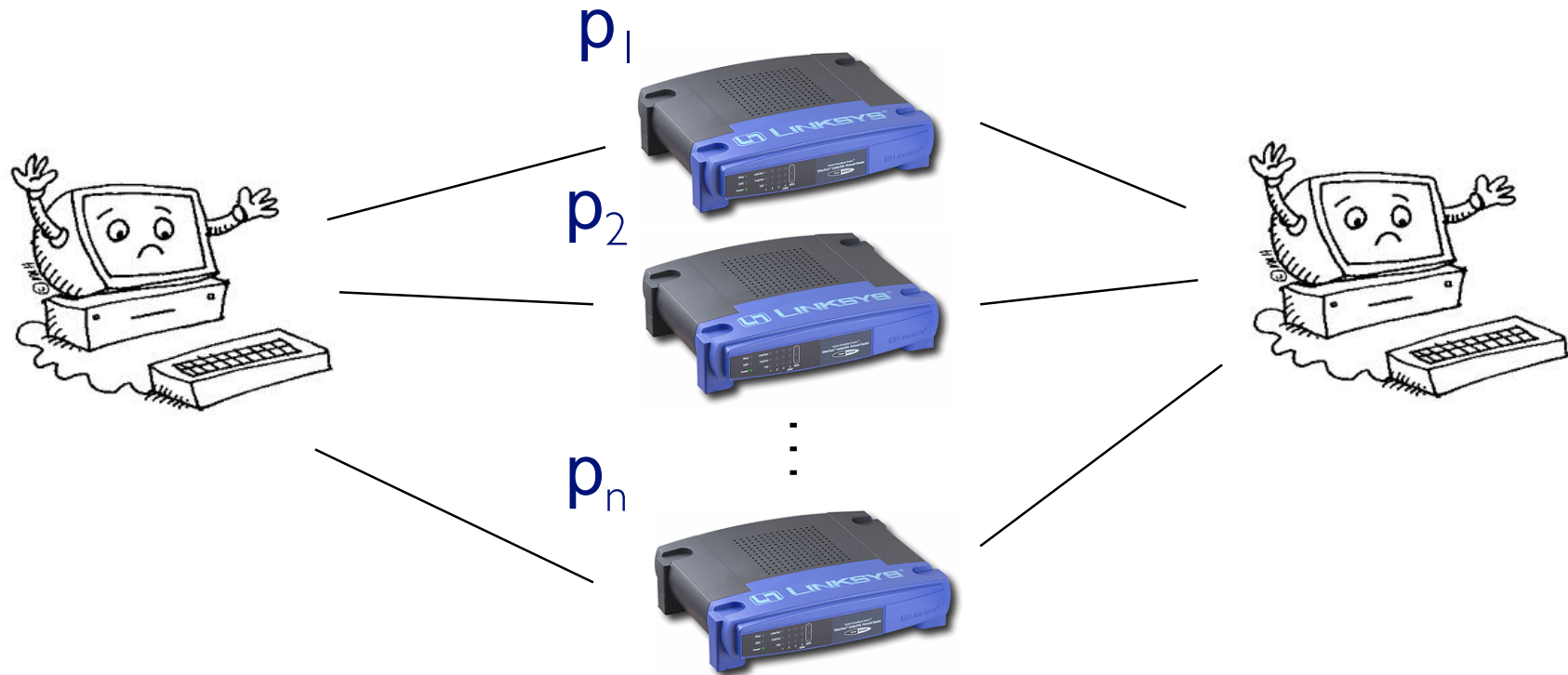
$$P(E) = P(F_1 \cup \dots \cup F_k) = 1 - P((F_1 \cup \dots \cup F_k)^c)$$

$$= 1 - P(F_1^c \cap F_2^c \cap \dots \cap F_k^c)$$

$$= 1 - P(\text{no strings hashed to buckets 1 to k})$$

$$= 1 - (1 - p_1 - p_2 - \dots - p_k)^m$$

Consider the following parallel network

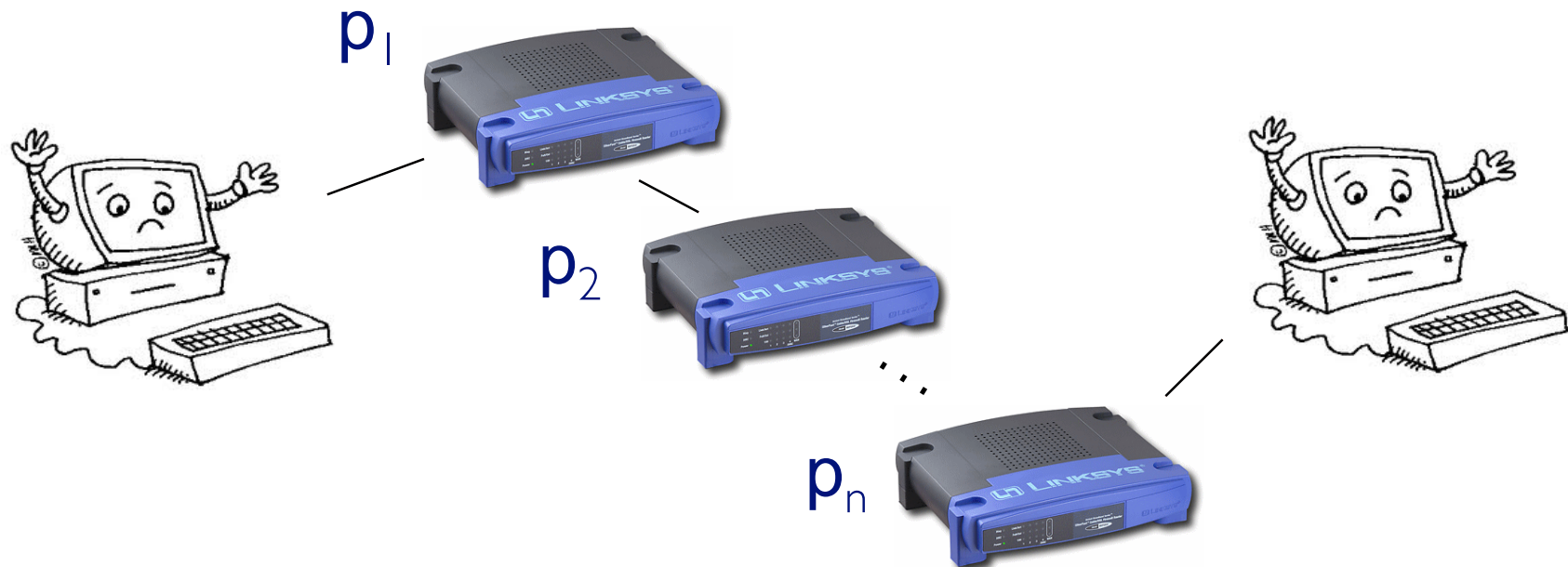


n routers, i^{th} has probability p_i of failing, independently

$P(\text{there is functional path}) = 1 - P(\text{all routers fail})$

$$= 1 - p_1 p_2 \cdots p_n$$

Contrast: a series network



n routers, i^{th} has probability p_i of failing, independently

$P(\text{there is functional path}) =$

$$P(\text{no routers fail}) = (1 - p_1)(1 - p_2) \cdots (1 - p_n)$$

Recall: Two events E and F are independent if

$$P(EF) = P(E) P(F)$$

If E & F are independent, does that tell us anything about

$$P(EF|G), P(E|G), P(F|G),$$

when G is an arbitrary event? In particular, is

$$P(EF|G) = P(E|G) P(F|G) ?$$

In general, *no*.

Roll two 6-sided dice, yielding values D_1 and D_2

$$E = \{ D_1 = 1 \}$$

$$F = \{ D_2 = 6 \}$$

$$G = \{ D_1 + D_2 = 7 \}$$

E and F are independent

$$P(E|G) = 1/6$$

$$P(F|G) = 1/6, \text{ but}$$

$$P(EF|G) = 1/6, \text{ not } 1/36$$

so $E|G$ and $F|G$ are not independent!

Two events E and F are called *conditionally independent given G* , if

$$P(EF|G) = P(E|G) P(F|G)$$

Or, equivalently (assuming $P(F)>0$, $P(G)>0$),

$$P(E|FG) = P(E|G)$$

do CSE majors get fewer A's?

Say you are in a dorm with 100 students

10 are CS majors: $P(C) = 0.1$

30 get straight A's: $P(A) = 0.3$

3 are CS majors who get straight A's

$P(CA) = 0.03$

$P(CA) = P(C) P(A)$, so C and A *independent*

At faculty night, only CS majors and A students show up

So 37 students arrive

Of 37 students, 10 are CS \Rightarrow

$P(C | C \text{ or } A) = 10/37 = 0.27 < .3 = P(A)$

Seems CS major *lowers* your chance of straight A's ☹️

Weren't they supposed to be independent?

In fact, CS and A are *conditionally dependent* at fac night

conditioning can also break DEPENDENCE

Randomly choose a day of the week

$A = \{ \text{It is not a Monday} \}$

$B = \{ \text{It is a Saturday} \}$

$C = \{ \text{It is the weekend} \}$

A and B are dependent events

$$P(A) = 6/7, P(B) = 1/7, P(AB) = 1/7.$$

Now condition both A and B on C:

$$P(A|C) = 1, P(B|C) = 1/2, P(AB|C) = 1/2$$

$$P(AB|C) = P(A|C) P(B|C) \Rightarrow A|C \text{ and } B|C \text{ independent}$$



Dependent events can become independent
by conditioning on additional information!

Events E & F are *independent* if

$$P(EF) = P(E) P(F), \text{ or, equivalently } P(E|F) = P(E) \text{ (if } p(E) > 0)$$

More than 2 events are indep if, for *all subsets*, joint probability = product of separate event probabilities

Independence can greatly simplify calculations

For fixed G, conditioning on G gives a probability measure, $P(E|G)$

But “conditioning” and “independence” are orthogonal:

Events E & F that are (unconditionally) independent may become dependent when conditioned on G

Events that are (unconditionally) dependent may become independent when conditioned on G