

CSE 311 Section 08



Regular Expressions,
CFGs, Relations

Administrivia



Announcements & Reminders

- Midterm grades will be released shortly
- Homework 6 was due Wednesday (2/26)
- Homework 7 will be due Wednesday (3/5)
- Check your section participation grade on gradescope
 - If it different than what you expect, let your TA know



Regular Expressions



Regular Expressions

ϵ matches only the **empty string**

a matches only the one-character string a

$A \cup B$ matches all strings that either A matches or B matches (or both)

AB matches all strings that have a first part that A matches followed by a second part that B matches

A^* matches all strings that have any number of strings (even 0) that A matches, one after another ($\epsilon \cup A \cup AA \cup AAA \cup \dots$)

Definition of the *language*
matched by a regular expression

Task 1 – Regular Expressions

- a) Write a regular expression that matches base 10 numbers (e.g., there should be no leading zeroes).
- b) Write a regular expression that matches all base-3 numbers that are divisible by 3.
- c) Write a regular expression that matches all binary strings that contain the substring “111”, but not the substring “000”.

We will do (b) together, then work on c

Task 1 – Regular Expressions

- a) Write a regular expression that matches base 10 numbers (e.g., there should be no leading zeroes).

base-10 numbers:

Our everyday numbers!
Notice we have 10 symbols (0-9) to represent numbers.

$$256: (2 * 10^2) + (5 * 10^1) + (6 * 10^0)$$

base-2 numbers: (binary)

$$10: (1 * 2^1) + (0 * 2^0)$$

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(0 U 1 U 2 U 3 U 4 U 5 U 6 U 7 U 8 U 9)*

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“0101” or “091” **are not** Base-10 numbers

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Hint: you know that Base-10 numbers are divisible by 10 when they end in 0 (10, 20, 30, 40...)

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$0 \cup ((1 \cup 2)(0 \cup 1 \cup 2)^*0)$



all possible Base-3 numbers divisible by 3

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$(0 \cup 1)^* 111 (0 \cup 1)^*$



The Kleene-star has us generating any number of 0's

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Use careful case-work to modify this and produce only 0,1,or two 0's

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⚠ Cannot produce 1's with “0” or “00” like “1011101”

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⚠ Generates “000” like “00 01 111”

$(01 \cup 001 \cup 1)^* (0 \cup 00 \cup \epsilon) 111 (01 \cup 001 \cup 1)^* (0 \cup 00 \cup \epsilon)$ ✓ all binary strings with “111” and without “000”

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all binary strings that contain the substring “111”

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⚠ Cannot produce 1's with “0” or “00” like “1011101”

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Generates “000” like “00 01 111”

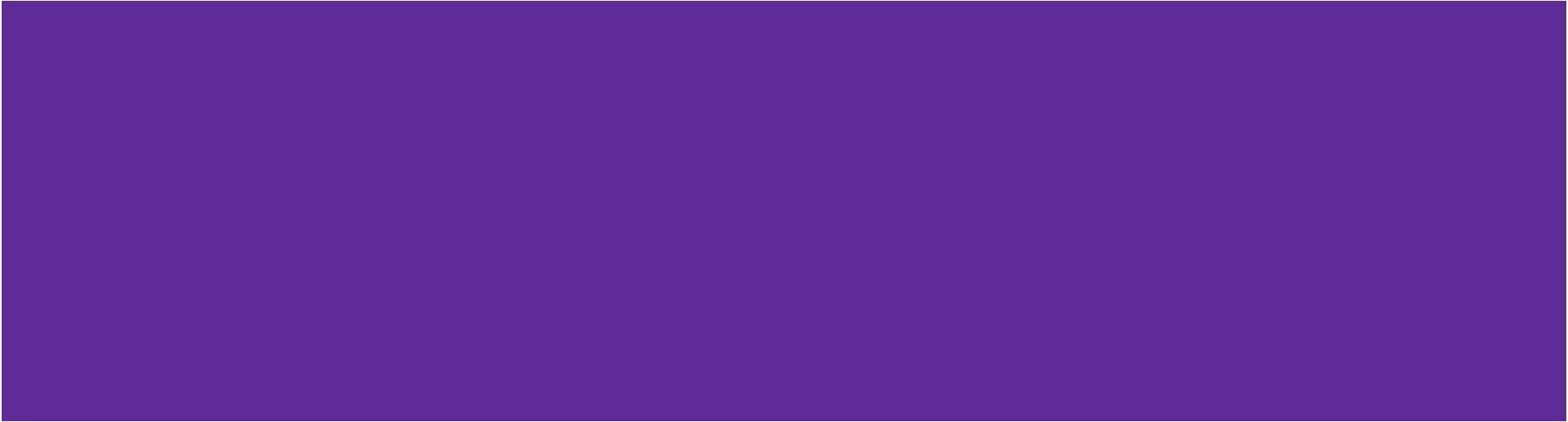
$(01 \cup 001 \cup 1)^* (0 \cup 00 \cup \epsilon) 111 (01 \cup 001 \cup 1)^* (0 \cup 00 \cup \epsilon)$



all binary strings with “111” and without “000”

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Context-Free Grammars



Context-Free Grammars

A context free grammar (CFG) is a finite set of production rules over:

- An alphabet Σ of “terminal symbols”
- A finite set V of “nonterminal symbols”
- A start symbol (one of the elements of V) usually denoted S

Always think back to Regex!

- CFG to match RE **$A \cup B$**

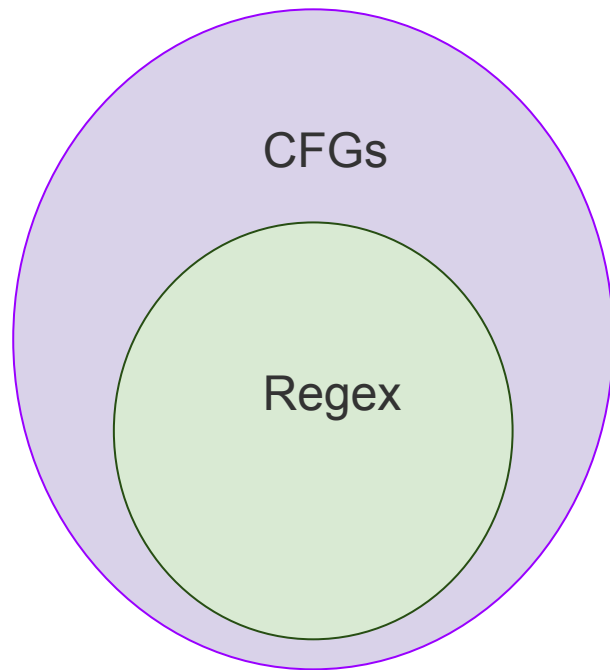
$S \rightarrow S_1 \mid S_2$ + rules from original CFGs

- CFG to match RE **AB**

$S \rightarrow S_1 S_2$ + rules from original CFGs

- CFG to match RE **A^*** $(= \epsilon \cup A \cup AA \cup AAA \cup \dots)$

$S \rightarrow S_1 S \mid \epsilon$ + rules from CFG with S_1



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- CFG to match RE **A** \cup **B**

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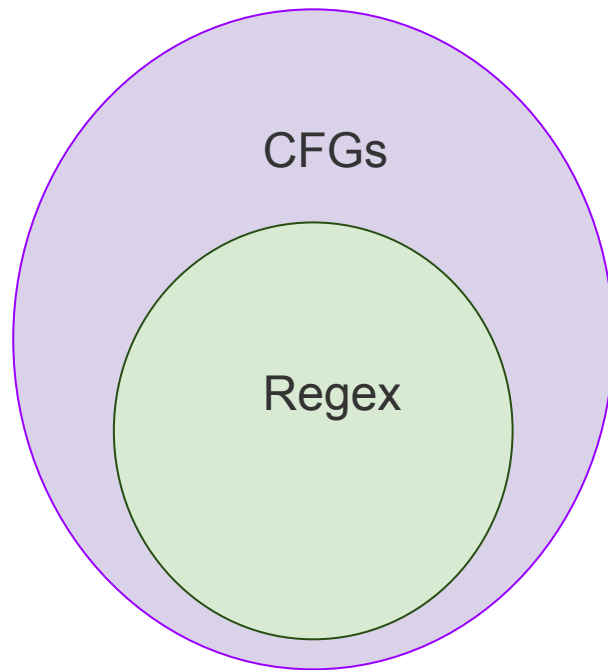
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- CFG to match RE **A**^{*} ($= \epsilon \cup \mathbf{A} \cup \mathbf{AA} \cup \mathbf{AAA} \cup \dots$)

$S \rightarrow S_1 S \mid \epsilon$ + rules from CFG with S_1

CFG or Regex?

“equal number of 0’s and 1’s” (ex. 011010)



Always think back to Regex!

- CFG to match RE **A** \cup **B**

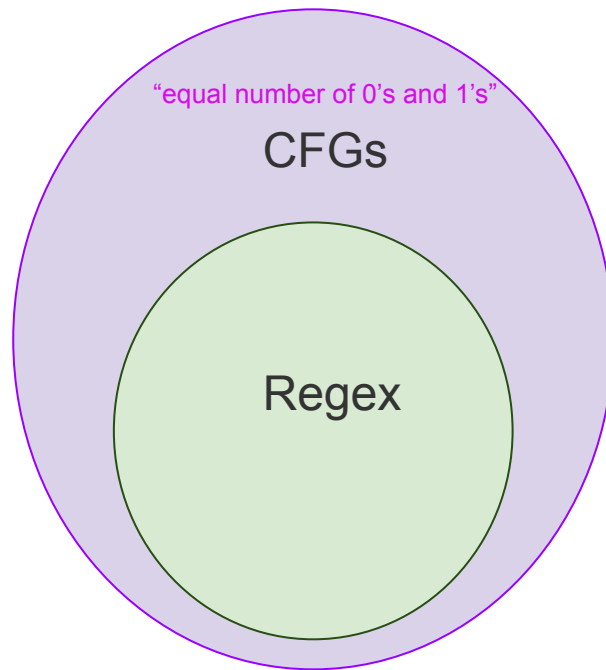
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$S \rightarrow S_1 S_2$ + rules from original CFGs

- CFG to match RE **A**^{*} ($= \epsilon \cup \mathbf{A} \cup \mathbf{AA} \cup \mathbf{AAA} \cup \dots$)

$S \rightarrow S_1 S \mid \epsilon$ + rules from CFG with S_1



Task 2 – CFGs

Write a context-free grammar to match each of these languages.

- a) All binary strings that start with 11.
- b) All binary strings that contain at most one 1.

Work on this problem with the people around you.

Task 2 – CFGs

- a) All binary strings that start with 11.

Task 2 – CFGs

- a) All binary strings that start with 11.

Thinking back to regular expressions...

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11 (0 ∪ 1)*

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Now generate the CFG...

Task 2 – CFGs

- a) All binary strings that start with 11.

Thinking back to regular expressions...

11 (0 \cup 1)*

Now generate the CFG...

S \rightarrow 11**T**

T \rightarrow 1**T** | 0**T** | ϵ

Task 2 – CFGs

- b) All binary strings that contain at most one 1.

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$0^* (1 \cup \epsilon) 0^*$

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- b) All binary strings that contain at most one 1.

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$0^* (1 \cup \epsilon) 0^*$

Now generate the CFG...

$S \rightarrow ABA$

$A \rightarrow 0A \mid \epsilon$

$B \rightarrow 1 \mid \epsilon$

Task 2 – CFGs

- b) All binary strings that contain at most one 1.

Thinking back to Regular expressions...

$0^* (1 \cup \epsilon) 0^*$

Now generate the CFG...

$S \rightarrow ABA$

$A \rightarrow 0A \mid \epsilon$

$B \rightarrow 1 \mid \epsilon$

Alternative solution:

$S \rightarrow 0S \mid S0 \mid 1 \mid 0 \mid \epsilon$

Relations



Relations

R is **reflexive** iff $(a,a) \in R$ for every $a \in A$

R is **symmetric** iff $(a,b) \in R$ implies $(b,a) \in R$

R is **antisymmetric** iff $(a,b) \in R$ and $a \neq b$ implies $(b,a) \notin R$

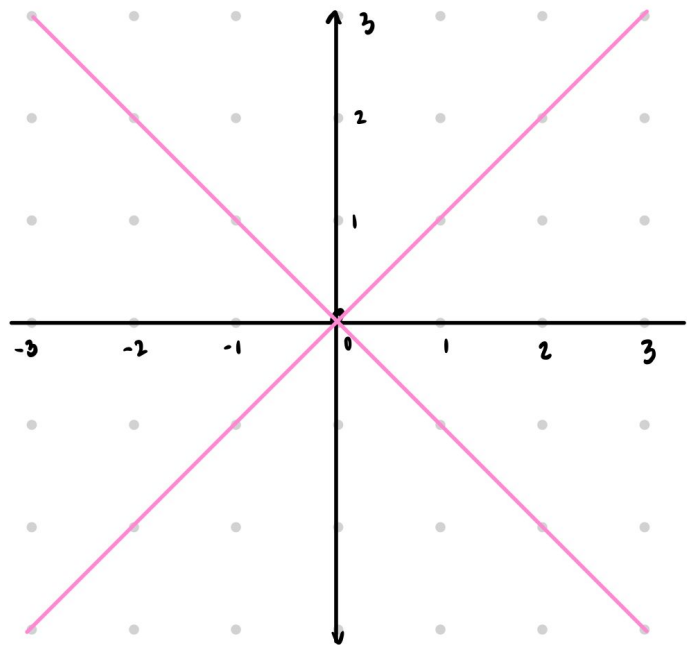
R is **transitive** iff $(a,b) \in R$ and $(b,c) \in R$ implies $(a,c) \in R$

Task 4b

Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .

Task 4b

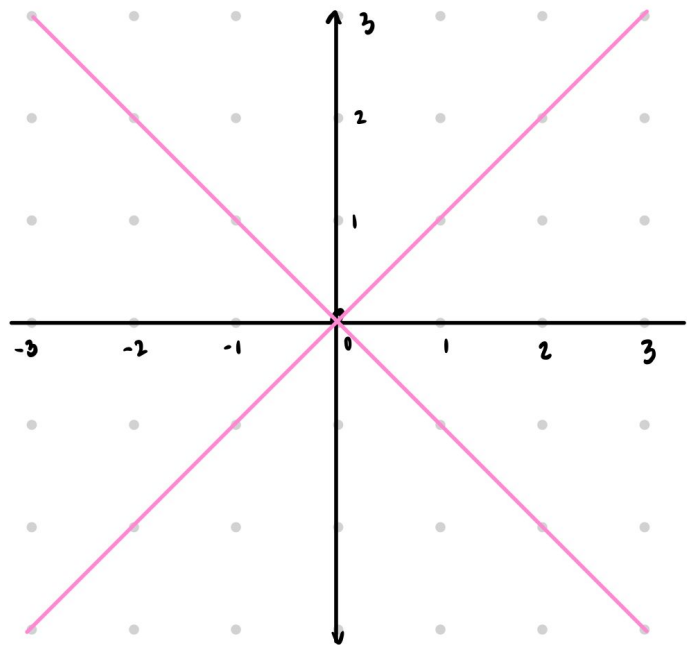
Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .



We can graph the points of R

Task 4b

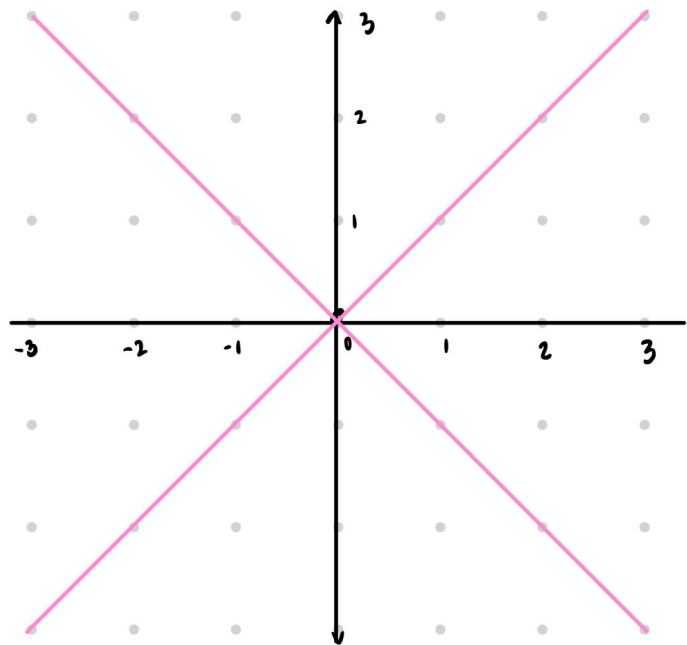
Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .



If all points on the line of $y = x$ are in the relation then the relation is reflexive

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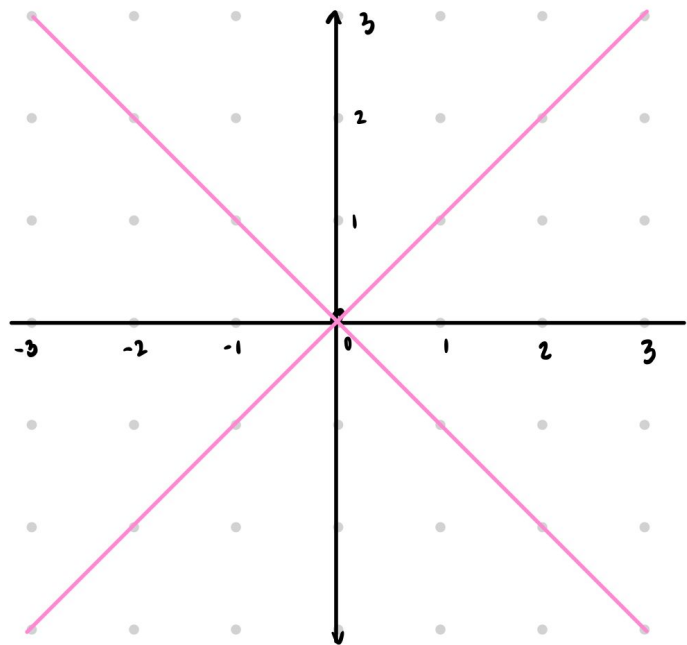


If all points on the line of $y = x$ are in the relation then the relation is reflexive

The relation is reflexive!

Task 4b

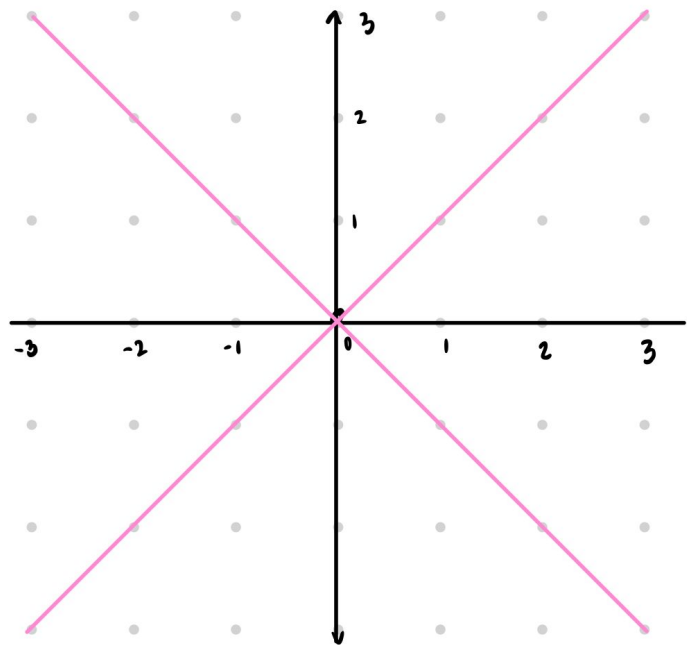
Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .



If all points that are reflected across $y = x$ are also in the relation, then the relation is symmetric

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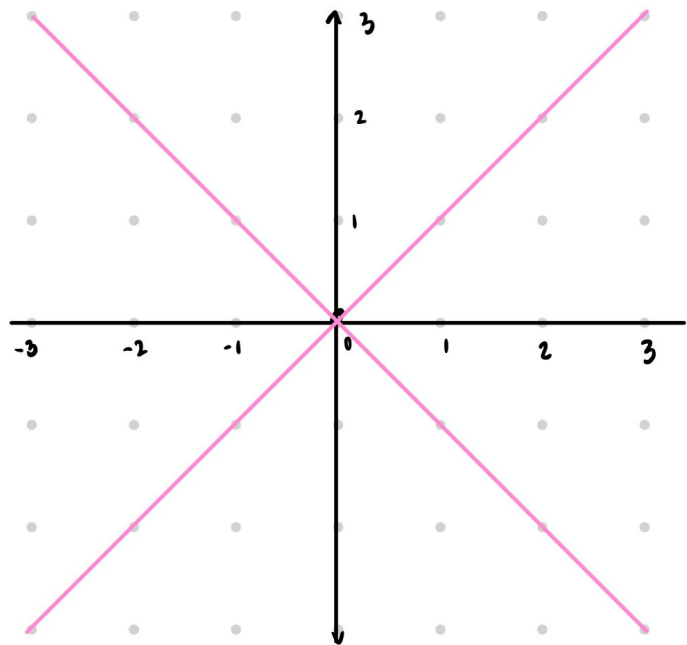


If all points that are reflected across $y = x$ are also in the relation, then the relation is symmetric

The relation is symmetric!

Task 4b

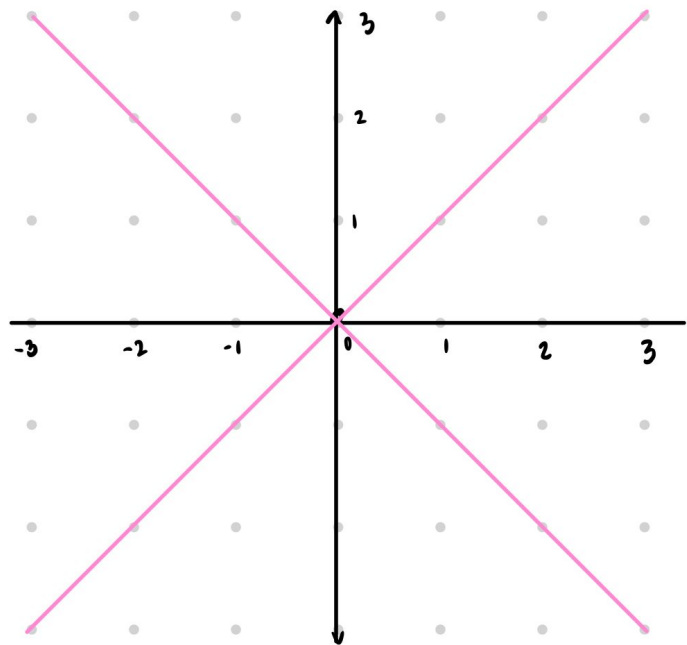
Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .



If all points that are reflected across $y = x$ are not in the relation, then the relation is antisymmetric

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If all points that are reflected across $y = x$ are not in the relation, then the relation is antisymmetric

The relation is not antisymmetric!

Task 4b

Let $R = \{(x, y) : x^2 = y^2\}$ on \mathbb{R} .

reflexive, symmetric, not antisymmetric (counterexample: $(-2, 2) \in R$ and $(2, -2) \in R$ but $2 \neq -2$), transitive

Proving Relations!



Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

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Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

R is reflexive means $\forall a \in A, (a, a) \in R$

We are trying to prove $\forall x, x \in R \rightarrow x \in R \circ R$

It can be helpful to convert definitions to formal logic so you know precisely what you are proving and where to get started.

Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

Let x be an arbitrary object.

Since x was arbitrary, we have proven that R is a subset of $R \circ R$.

Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

Let x be an arbitrary object.

Suppose that $x \in R$.

Since x was arbitrary, we have proven that R is a subset of $R \circ R$.

Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

Let x be an arbitrary object.

Suppose that $x \in R$.

Since $x \in R$, $x = (a, b)$ for some a and b in A .

Since x was arbitrary, we have proven that R is a subset of $R \circ R$.

Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

Let x be an arbitrary object.

Suppose that $x \in R$.

Since $x \in R$, $x = (a, b)$ for some a and b in A .

We want to show that $(a, b) \in R \circ R$. So we need to show that there is some c such that (a, c) is in R and (c, b) in R .

Thus, $x \in R \circ R$, by the definition of composition.

Since x was arbitrary, we have proven that R is a subset of $R \circ R$.

Task 6

Let R be a relation on a set A . Given that R is reflexive, it follows that $R \subseteq R \circ R$.

Let x be an arbitrary object.

Suppose that $x \in R$.

Since $x \in R$, $x = (a, b)$ for some a and b in A .

Since R is reflexive, we know that $(b, b) \in R$ as well. This shows that there is some $c \in A$ (namely, $c = b$) such that $(a, c) \in R$ and $(c, b) \in R$.

Thus, $x \in R \circ R$, by the definition of composition.

Since x was arbitrary, we have proven that R is a subset of $R \circ R$.

That's All, Folks!

Thanks for coming to section this week!
Any questions?