

Fibonacci Inequality: skeleton

Show that $f(n) \leq 2^n$ for all $n \geq 0$ by induction.

Define $P(n)$ to be " $f(n) \leq 2^n$ ". We show $P(n)$ is true for all $n \geq 0$ by induction on n .

Base Cases: ($n = 0$): $f(0) = 1 \leq 1 = 2^0$.

($n = 1$): $f(1) = 1 \leq 2 = 2^1$.

Inductive Hypothesis: Suppose $P(0) \wedge P(1) \wedge \dots \wedge P(k)$ for an arbitrary $k \geq 1$.

Inductive step:

Target: $P(k + 1)$. i.e. $f(k + 1) \leq 2^{k+1}$

$$f(0) = 1; \quad f(1) = 1$$

$$f(n) = f(n-1) + f(n-2) \text{ for all } n \in \mathbb{N}, n \geq 2.$$

Induction: Hats! (goal)

You have n people in a line ($n \geq 2$). Each of them wears either a **purple hat** or a **gold hat**. The person at the front of the line wears a purple hat. The person at the back of the line wears a gold hat.

Show that for every arrangement of the line satisfying the rule above, there is a person with a purple hat next to someone with a gold hat.

Yes this is kinda obvious. I promise this is good induction practice.

Yes you could argue this by contradiction. I promise this is good induction practice.

What is $P(n)$?

String proof: skeleton

Let $P(y)$ be " $\text{len}(x \cdot y) = \text{len}(x) + \text{len}(y)$ for all $x \in \Sigma^*$."

We prove $P(y)$ for all $x \in \Sigma^*$ by structural induction.

Base Case:

Inductive Hypothesis

Inductive Step:

We conclude that $P(y)$ holds for all string y by the principle of induction. Unwrapping the definition of P , we get $\forall x \forall y \in \Sigma^* \text{len}(xy) = \text{len}(x) + \text{len}(y)$, as required.

String proof: Inductive Step Setup

Let $P(y)$ be " $\text{len}(x \cdot y) = \text{len}(x) + \text{len}(y)$ for all $x \in \Sigma^*$."

We prove $P(y)$ for all $x \in \Sigma^*$ by structural induction.

Base Case: Let x be an arbitrary string, $\text{len}(x \cdot \epsilon) = \text{len}(x)$
 $= \text{len}(x) + 0 = \text{len}(x) + \text{len}(\epsilon)$

Inductive Hypothesis: Suppose $P(w)$ for an arbitrary string w .

Inductive Step: Let $y = wa$ for an arbitrary $a \in \Sigma$. We show $P(y)$. Let x be an arbitrary string.

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Therefore, $\text{len}(xy) = \text{len}(x) + \text{len}(y)$, as required.

We conclude that $P(y)$ holds for all string y by the principle of induction. Unwrapping the definition of P , we get $\forall x \forall y \in \Sigma^* \text{len}(xy) = \text{len}(x) + \text{len}(y)$, as required.