# CSE 311 Section 6

Induction

#### Administrivia

#### **Announcements & Reminders**

- HW4
  - Grades out now
  - If you think something was graded incorrectly, submit a regrade request!
- HW5
  - due yesterday (10/30 @ 11:59 pm)
- HW6
  - o due 11/06 @ 11:59 pm on Gradescope
- Midterm is COMING!!!
  - Wednesday 11/13 @ 6-7:30 pm in PAA
  - More information coming soon!

# Induction

#### (Weak) Induction Template

Let P(n) be "(whatever you're trying to prove)".

We show P(n) holds for all n by induction on n.

Base Case: Show P(b) is true.

<u>Inductive Hypothesis:</u> Suppose P(k) holds for an arbitrary  $k \geq b$ .

<u>Inductive Step:</u> Show P(k + 1) (i.e. get  $P(k) \rightarrow P(k + 1)$ )

### (Weak) Induction Template

Let P(n) be "(whatever you're trying to prove)". We show P(n) holds **for all** n by induction on n. Note: often you will condition n here, like "all natural numbers n" or " $n \ge 0$ "

Base Case: Show P(b) is true.

<u>Inductive Hypothesis:</u> Suppose P(k) holds for an arbitrary  $k \geq b$ .

<u>Inductive Step:</u> Show P(k + 1) (i.e. get  $P(k) \rightarrow P(k + 1)$ )

<u>Conclusion</u>: Therefore, P(n) holds for all n by the principle of induction.

Match the earlier condition on *n* in your conclusion!

- a) Show using induction that  $0 + 1 + 2 + \dots + n = \frac{n(n+1)}{2}$  for all  $n \in \mathbb{N}$ .
- b) Define the triangle numbers as  $\triangle_n = 1 + 2 + \dots + n$ , where  $n \in \mathbb{N}$ . In part (a) we showed  $\triangle_n = \frac{n(n+1)}{2}$ . Prove the following equality for all  $n \in \mathbb{N}$ :  $0^3 + 1^3 + \dots + n^3 = \triangle_n^2$

Lets walk through part (a) together.

We can "fill in" our induction template to construct our proof by induction.

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

Let P(n) be "". We show P(n) holds for (some) n by induction on n.

Base Case: P(b):

Inductive Hypothesis: Suppose P(k) holds for an arbitrary  $k \ge b$ .

Inductive Step: Goal: Show P(k + 1):

<u>Conclusion:</u> Therefore, P(n) holds for (some) n by the principle of induction.

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

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Let P(n) be " $0 + 1 + 2 + \cdots + n = \frac{n(n+1)}{2}$ ". We show P(n) holds for all  $n \in \mathbb{N}$  by induction on n.

Base Case: P(0):  $0 + \cdots = 0 = \frac{0(0+1)}{2}$  so the base case holds.

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Inductive Step: Goal: Show  $P(k + 1): 0 + 1 + \dots + k + (k + 1) = \frac{(k+1)(k+2)}{2}$ 

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

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Inductive Step: Goal: Show  $P(k + 1): 0 + 1 + \dots + k + (k + 1) = \frac{(k+1)(k+2)}{2}$ 

$$0 + 1 + \cdots + k + (k + 1) = \cdots$$

...

$$=\frac{(k+1)(k+2)}{2}$$

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

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$$0 + 1 + \dots + k + (k + 1) = (0 + 1 + \dots + k) + (k + 1)$$

...

$$=\frac{(k+1)(k+2)}{2}$$

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

Let 
$$P(n)$$
 be " $0+1+2+\cdots+n=\frac{n(n+1)}{2}$ ". We show  $P(n)$  holds for all  $n\in\mathbb{N}$  by induction on  $n$ .

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Inductive Step: Goal: Show 
$$P(k + 1): 0 + 1 + \dots + k + (k + 1) = \frac{(k+1)(k+2)}{2}$$

$$0+1+\cdots+k+(k+1)=(0+1+\cdots+k)+(k+1)$$
 
$$=\frac{k(k+1)}{2}+(k+1)$$
 by I.H.

...

$$=\frac{(k+1)(k+2)}{2}$$
?

Show using induction that  $0+1+2+\cdots+n=\frac{n(n+1)}{2}$  for all  $n\in\mathbb{N}$ .

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Inductive Step: Goal: Show 
$$P(k + 1): 0 + 1 + \dots + k + (k + 1) = \frac{(k+1)(k+2)}{2}$$

$$0+1+\dots+k+(k+1) = (0+1+\dots+k)+(k+1)$$

$$= \frac{k(k+1)}{2}+(k+1)$$

$$= \frac{k(k+1)}{2}+\frac{2(k+1)}{2}$$
by I.H.

$$=\frac{(k+1)(k+2)}{2}$$

# Show using induction that $0 + 1 + 2 + \dots + n = \frac{n(n+1)}{2}$

for all  $n \in \mathbb{N}$ .

### Problem 1 – Induction with Equality

Let P(n) be " $0+1+2+\cdots+n=\frac{n(n+1)}{2}$ ". We show P(n) holds for all  $n\in\mathbb{N}$  by induction on n.

Base Case: P(0):  $0 + \cdots = 0 = \frac{0(0+1)}{2}$  so the base case holds.

Inductive Hypothesis: Suppose P(k) holds for an arbitrary  $k \ge 0$ , i.e.  $0 + 1 + 2 + \cdots + k = \frac{k(k+1)}{2}$ 

<u>Inductive Step:</u> Goal: Show  $P(k + 1): 0 + 1 + \dots + k + (k + 1) = \frac{(k+1)(k+2)}{2}$ 

$$0 + 1 + \dots + k + (k + 1) = (0 + 1 + \dots + k) + (k + 1)$$

$$= \frac{k(k+1)}{2} + (k + 1)$$

$$= \frac{k(k+1)}{2} + \frac{2(k+1)}{2}$$

$$= \frac{k(k+1) + 2(k+1)}{2}$$

$$= \frac{(k+1)(k+2)}{2}$$
?

# Show using induction that $0 + 1 + 2 + \dots + n = \frac{n(n+1)}{2}$

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### Problem 1 – Induction with Equality

Let P(n) be " $0+1+2+\cdots+n=\frac{n(n+1)}{2}$ ". We show P(n) holds for all  $n\in\mathbb{N}$  by induction on n.

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Inductive Step: Goal: Show  $P(k+1): 0+1+\cdots+k+(k+1)=\frac{(k+1)(k+2)}{2}$ 

$$0+1+\dots+k+(k+1) = (0+1+\dots+k)+(k+1)$$

$$= \frac{k(k+1)}{2}+(k+1)$$

$$= \frac{k(k+1)}{2}+\frac{2(k+1)}{2}$$

$$= \frac{k(k+1)+2(k+1)}{2}$$

$$= \frac{(k+1)(k+2)}{2}$$
factoring

factoring out (k+1)

- a) Show using induction that  $0 + 1 + 2 + \dots + n = \frac{n(n+1)}{2}$  for all  $n \in \mathbb{N}$ .
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$$0^3 + 1^3 + \dots + n^3 = \triangle_n^2$$

 $\triangle_n = 1 + 2 + \dots + n, \ n \in \mathbb{N}.$   $\triangle_n = \frac{n(n+1)}{2}. \text{ Prove for all } n \in \mathbb{N}:$   $0^3 + 1^3 + \dots + n^3 = \triangle_n^2$ 

Let P(n) be "". We show P(n) holds for (some) n by induction on n.

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<u>Conclusion:</u> Therefore, P(n) holds for (some) n by the principle of induction.

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Let P(n) be " $0^3 + 1^3 + \dots + n^3 = (0 + 1 + \dots + n)^2$ ". We show P(n) holds for all  $n \in \mathbb{N}$  by induction on n.

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Inductive Hypothesis: Suppose P(k) holds for an arbitrary  $k \ge 0$ . i.e.  $0^3 + 1^3 + \dots + k^3 + = (0 + 1 + \dots + k)^2$ 

Inductive Step: Goal: Show P(k + 1):

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$$\triangle_n = \frac{n(n+1)}{2}. \text{ Prove for all } n \in \mathbb{N}:$$

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<u>Inductive Hypothesis:</u> Suppose P(k) holds for an arbitrary  $k \ge 0$ . i.e.  $0^3 + 1^3 + \dots + k^3 + = (0 + 1 + \dots + k)^2$ 

Inductive Step: Goal: Show P(k+1):  $0^3 + 1^3 + \dots + k^3 + (k+1)^3 = (0+1+\dots+k+(k+1))^2$ 

$$\triangle_n = 1 + 2 + \dots + n, \ n \in \mathbb{N}.$$

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Inductive Step: Goal: Show P(k+1):  $0^3 + 1^3 + \dots + k^3 + (k+1)^3 = (0+1+\dots+k+(k+1))^2$ 

$$0^3 + 1^3 + \dots + k^3 + (k+1)^3 = \dots$$

 $= (0 + 1 + \dots + k + (k + 1))^2$ 

 $\triangle_n = 1 + 2 + \cdots + n, n \in \mathbb{N}.$ 

 $\triangle_n = \frac{n(n+1)}{2}$ . Prove for all  $n \in \mathbb{N}$ :  $0^3 + 1^3 + \cdots + n^3 = \triangle_n^2$ 

by (a)

Let P(n) be " $0^3 + 1^3 + \cdots + n^3 = (0 + 1 + \cdots + n)^2$ ". We show P(n) holds for all  $n \in \mathbb{N}$  by induction on n.

Base Case: P(0):  $0^3 = 0 = (0)^2$  so the base case holds.

i.e.  $0^3 + 1^3 + \dots + k^3 + = (0 + 1 + \dots + k)^2$ Inductive Hypothesis: Suppose P(k) holds for an arbitrary  $k \geq 0$ .

Inductive Step: Goal: Show P(k+1):  $0^3 + 1^3 + \dots + k^3 + (k+1)^3 = (0+1+\dots+k+(k+1))^2$  $0^3 + 1^3 + \dots + k^3 + (k+1)^3 = (0+1+\dots+k)^2 + (k+1)^3$ by I.H.

$$f(k) = (k + 1)^3 = (k + 1)^3 + (k + 1)^3$$

$$f(k) = (k + 1)^2 (k^2 + (k + 1))$$

$$f(k) = (k + 1)^2 (k^2 + (k + 1))$$

$$f(k) = (k + 1)^2 (k^2 + 4k + 4)$$

$$f(k) = (k + 1)^2 (k^2 + 4k + 4)$$

$$= (k+1)^2 \left(\frac{k+1}{4}\right)$$

$$= (k+1)^2 \left(\frac{(k+2)^2}{4}\right)$$

$$= \left(\frac{(k+1)(k+2)}{2}\right)^2$$
factor numerator

 $= (0 + 1 + \dots + k + (k + 1))^2$ 

# **Strong Induction**

### Why Strong Induction?

In **weak induction**, the inductive hypothesis only assumes that P(k) is true and uses that in the inductive step to prove the implication  $P(k) \rightarrow P(k+1)$ .

In **strong induction**, the inductive hypothesis assumes the predicate holds for every step from the base case(s) up to P(k). This usually looks something like  $P(b_1) \wedge P(b_2) \wedge \cdots \wedge P(k)$ . Then it uses this stronger inductive hypothesis in the inductive step to prove the implication  $P(b_1) \wedge \cdots \wedge P(k) \rightarrow P(k+1)$ .

Strong induction is necessary when we have multiple base cases, or when we need to go back to a smaller number than k in our inductive step.

## **Strong Induction Template**

Let P(n) be "(whatever you're trying to prove)". We show P(n) holds for all  $n \ge b_{min}$  by induction on n.

Base Case: Show  $P(b_{min})$ ,  $P(b_{min+1})$ , ...,  $P(b_{max})$  are all true.

Inductive Hypothesis: Suppose  $P(b_{min}) \land \cdots \land P(k)$  hold for an arbitrary  $k \ge b_{max}$ .

<u>Inductive Step:</u> Show P(k+1) (i.e. get  $P(b_{min}) \land \cdots \land P(k) \rightarrow P(k+1)$ )

Xavier Cantelli owns some rabbits. The number of rabbits he has in any given year is described by the function f:

$$f(0) = 0$$
  
 $f(1) = 1$   
 $f(n) = 2f(n-1) - f(n-2)$  for  $n \ge 2$ 

Determine, with proof, the number, f(n), of rabbits that Cantelli owns in year n. That is, construct a formula for f(n) and prove its correctness.

First, let's construct a formula for f(n). How many rabbits does he have each year? Let's do some calculations, and see if we can find a pattern. Then, we'll use induction to prove the pattern holds for all n!

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 $f(3) = 2f(3-1) - f(3-2) = 2f(2) - f(1) = 2(2) - 1 = 4 - 1 = 3$ 

Xavier Cantelli owns some rabbits. The number of rabbits he has in any given year is described by the function f:

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Xavier Cantelli owns some rabbits. The number of rabbits he has in any given year is described by the function f:

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Determine, with proof, the number, f(n), of rabbits that Cantelli owns in year n. That is, construct a formula for f(n) and prove its correctness.

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It seems like we have a pattern here!

$$f(n) = n$$

But we don't want to have to check for EVERY n, so let's see if we can prove it with induction instead!

What kind of induction should we use?

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Strong induction!

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Strong induction!

#### Two big clues:

- Multiple base cases in the formula: f(0) = 0 and f(1) = 1
- Recursively defined step of formula goes back further than just n:
  - ∘ f(n) based on both f(n-1) and f(n-2)
  - o for P(n) to be true, both P(n-1) and P(n-2) must be true

Let P(n) be "(whatever you're trying to prove)". We show P(n) holds for all  $n \ge b_{min}$  by induction on n.

Base Case: Show  $P(b_{min})$ ,  $P(b_{min+1})$ , ...,  $P(b_{max})$  are all true.

Inductive Hypothesis: Suppose  $P(b_{min}) \land \cdots \land P(k)$  hold for an arbitrary  $k \ge b_{max}$ .

<u>Inductive Step:</u> Show P(k + 1) (i.e. get  $P(b_{min}) \land \cdots \land P(k) \rightarrow P(k + 1)$ )

<u>Conclusion</u>: Therefore, P(n) holds for all  $n \ge b_{min}$  by the principle of induction.

Fill in the strong induction template to prove the claim!

Let P(n) be "".

We show P(n) holds ...

**Base Cases:** 

**Inductive Hypothesis:** 

**Inductive Step:** 

<u>Conclusion:</u> Therefore, P(n) holds for all ... by the principle of induction.

Let P(n) be "f(n) = n".

We show P(n) holds for all  $n \ge 0$  by induction on n.

**Base Cases:** 

**Inductive Hypothesis:** 

**Inductive Step:** 

```
Let P(n) be "f(n) = n".
We show P(n) holds for all n \ge 0 by induction on n.
Base Cases: (n = 0, n = 1): f(0) = 0 and f(1) = 1 by definition of f.
Inductive Hypothesis:
```

**Inductive Step:** 

```
Let P(n) be "f(n) = n".
```

We show P(n) holds for all  $n \ge 0$  by induction on n.

Base Cases: (n = 0, n = 1): f(0) = 0 and f(1) = 1 by definition of f.

Inductive Hypothesis: Suppose  $P(0) \wedge P(1) \wedge \cdots \wedge P(k)$  hold for an arbitrary all  $k \geq 1$ .

#### **Inductive Step:**

```
Let P(n) be "f(n) = n". We show P(n) holds for all n \ge 0 by induction on n. Base Cases: (n = 0, n = 1): f(0) = 0 and f(1) = 1 by definition of f. Inductive Hypothesis: Suppose P(0) \land P(1) \land \cdots \land P(k) hold for an arbitrary all k \ge 1. i.e. f(k) = k, f(k - 1) = k - 1, f(k - 2) = k - 2, etc. Inductive Step:
```

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<u>Inductive Hypothesis:</u> Suppose P(0) \wedge P(1) \wedge \cdots \wedge P(k) hold for an arbitrary all k \geq 1.
i.e. f(k) = k, f(k-1) = k-1, f(k-2) = k-2, etc.
<u>Inductive Step:</u> Goal: Show P(k+1): f(k+1) = k+1
f(k+1) = ...
          = k + 1
```

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<u>Inductive Step:</u> Goal: Show P(k+1): f(k+1) = k+1
f(k+1) = 2f(k) - f(k-1)
                                           definition of f
         = k + 1
```

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<u>Inductive Step:</u> Goal: Show P(k+1): f(k+1) = k+1
f(k+1) = 2f(k) - f(k-1)
                                           definition of f
         = 2(k) - (k-1)
                                           by I.H.
         = k + 1
```

# That's All, Folks!

Thanks for coming to section this week!

Any questions?