

## Structural Induction Template

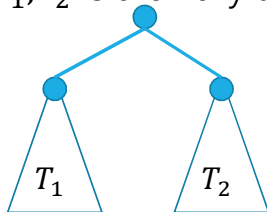
1. Define  $P()$  Show that  $P(x)$  holds for all  $x \in S$ . State your proof is by structural induction.
2. Base Case: Show  $P(x)$  for all base cases  $x$  in  $S$ .
3. Inductive Hypothesis: Suppose  $P(x)$  for all  $x$  listed as in  $S$  in the recursive rules.
4. Inductive Step: Show  $P()$  holds for the "new element" given.  
You will need a separate step for every rule.
5. Therefore  $P(x)$  holds for all  $x \in S$  by the principle of induction.

## Binary Trees

Basis: A single node is a rooted binary tree.

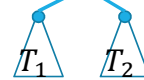


Recursive Step: If  $T_1$  and  $T_2$  are rooted binary trees with roots  $r_1$  and  $r_2$ , then a tree rooted at a new node, with children  $r_1, r_2$  is a binary tree.



$$\text{size}(\bullet) = 1$$

$$\text{size}(\text{tree}) =$$



$$\text{size}(T_1) + \text{size}(T_2) + 1$$

$$\text{height}(\bullet) = 0$$

$$\text{height}(\text{tree}) =$$



$$1 + \max(\text{height}(T_1), \text{height}(T_2))$$

## Regular Expressions

### Basis:

$\varepsilon$  is a regular expression. The empty string itself matches the pattern (and nothing else does).

$\emptyset$  is a regular expression. No strings match this pattern.

$a$  is a regular expression, for any  $a \in \Sigma$  (i.e. any character). The character itself matching this pattern.

### Recursive

If  $A, B$  are regular expressions then  $(A \cup B)$  is a regular expression matched by any string that matches  $A$  or that matches  $B$  [or both].

If  $A, B$  are regular expressions then  $AB$  is a regular expression. matched by any string  $x$  such that  $x = yz$ ,  $y$  matches  $A$  and  $z$  matches  $B$ .

If  $A$  is a regular expression, then  $A^*$  is a regular expression. matched by any string that can be divided into 0 or more strings that match  $A$ .

## More Examples

$(0^*1^*)^*$

$0^*1^*$

$(0 \cup 1)^*(00 \cup 11)^*(0 \cup 1)^*$

$(00 \cup 11)^*$